Stability Analysis for Adaptive Behavior (Position Paper)

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Abstract: One of the biggest challenges related to the research and development of autonomous systems is to prove the correctness of their autonomy. Nowadays, autonomous and adaptive systems are the roadmap to AI and the verification of such systems needs to set boundaries that will provide the highest possible guarantees that AI will be safe and sound, so trust can be established in its innocuous operation. In this paper, the authors draw upon their work on integrating stabilization science as part of a mechanism for verification of adaptive behavior. Stability analysis is studied to find an approach that helps to determine stable states of the behavior of an autonomous system. These states are further analyzed to determine behavior trajectories and equilibrium orbits. KnowLang, a formal method for knowledge representation and reasoning of adaptive systems, is used as a platform for stability analysis of autonomous systems.

1 INTRODUCTION

In today's technologies, the terms autonomous and adaptive started to progressively underline the new trends of research and development of softwareintensive systems. Autonomous systems, such as automatic lawn mowers, smart home equipment, driverless train systems, or autonomous cars, perform their tasks without human intervention. Eventually, the most challenging question which comes up when following the life cycle of the terms "autonomy" and "adaptation" is the potential to construct a system that behaves and operates similarly to, or even better than, a human being. Obviously, this is the roadmap to AI, and proving the correctness of autonomous and adaptive systems becomes extremely important. The point is that the verification of an autonomous system needs to set the boundaries of such AI and provide autonomic operations at least in a certain context with highest safety guarantees, and finally establish trust in its innocuous operation.

This paper draws upon the authors' work on a special approach to Adaptive Behavior Verification (ABV) (Vassev and Hinchey, 2014) where *stability analysis* is performed to identify unstable behavior

that will eventually require autonomy and adaptation to restore the stability of the system.

2 VERIFICATION OF ADAPTIVE BEHAVIOR

The ABV approach consists of the following parts (see Figure 1):

- a stability analysis capability that identifies instabilities given a system model and partitions the system model into stable and unstable component models;
- 2) a *state-space reduction capability* that prunes the state space of an unstable component model without loss of critical fidelity;
- high performance computing (HPC) simulations to explore component behavior over a wide range of an unstable component's reduced state space and produce a statistical verification for the component;
- 4) a *compositional verification capability* that integrates individual component verifications;

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Figure 1: Adaptive Behavior Verification (Vassev and Hinchey, 2014).

5) *operational monitors* to detect and take action to correct undesired unstable behavior of the system during operation.

This work is based on the principles of stabilization science (Arora, 2000). In our approach, stability analysis is performed over a model built with the KnowLang framework (Vassev and Hinchey, 2015a; Vassev and Hinchey, 2015b) (see Section 4) where the model is partitioned into system elements providing the overall behavior as a collection of stable (deterministic) and unstable (nondeterministic) components. The stable parts (or components) can then be verified with traditional verification methods, e.g., model checking. To verify the unstable parts we reduce their state space first and then use HPC simulation (see Figure 1) and testing. Finally, a compositional verification approach integrates the verification results of both stable and unstable parts by using special behavior invariants as approximation of safe states along with interaction invariants as interaction states.

3 STABILITY

In this work, the term "stable" informally means *resistant to change in behavior*. Stabilization science (Arora, 2000) provides a common approach to studying system stability where a system behavior is linearized around an operating point to determine a small-signal linearized model of that operating point. The stability of the system is then determined using *linear system stability analysis criterions* such as: Circle criterion (Shiriaev et al., 2003), Routh-Hurwitz criterion (Gantmacher, 1959), Nyquist stability criterion (Pippard, 1985), etc.

In this work, the *mathematical model for stability* means differential equations used for the simulation of a self-adaptive system. In general, we deal with a semi-discrete model, which is *discrete in space* and *continuous in time*. The mathematical model for stability is based on the *Lyapunov Stability Theory* (Lyapunov, 1892). Here, the stability of a system is

modeled for behaviors near to a *point of equilibrium* and where a behavior orbit around that point is defined stable if the forward orbit of any point is in a small enough neighborhood or it stays in a small (but perhaps, larger) neighborhood. In simple terms, if a behavior B_{x_e} determined by a sequence of actions that starts out near an equilibrium point x_e stays near x_e forever, then B_{x_e} is Lyapunov stable near x_e .

The theory behind the stability analysis is based on the *qualitative theory of differential equations* and dynamical systems, and deals with asymptotic properties of a system and the trajectories describing what happens with that system after a long period of time. A simple stable behavior B_{x_e} is exhibited by *equilibrium points* (or fixed points) and by *periodic orbits*.

We adapted the Lyapunov stability theory to address the following three questions:

- 1) Will a nearby behavior orbit indefinitely stay close to an *equilibrium orbit*?
- 2) Will a nearby behavior orbit converge to an *equilibrium orbit*?
- 3) Will a nearby behavior orbit depart away from an *equilibrium orbit*?

In the first case, the behavior is called *stable*, in the second case, it is called *asymptotically stable*, and in the third case, the behavior is said to be *unstable*.

An equilibrium solution f_e (where all the evaluated behaviors are stable) to an *autonomous* system is called:

- stable if for every small depart from the equilibrium orbit $\varepsilon > 0$, there exists a $\delta > 0$ such that every solution f(t) having initial conditions within distance δ , i.e., $||f(t_0) f_e|| < \delta$ of the equilibrium remains within distance ε , i.e., $||f(t) f_e|| < \varepsilon$ for all $t \ge t_0$;
- asymptotically stable if it is stable and, in addition, there exists $\delta_0 > 0$ such that whenever $\delta_0 > ||f(t_0) f_e||$ then $f(t) \rightarrow f_e$ as $t \rightarrow \infty$.

In this model stability means that the behavior trajectories do not change too much under small perturbations, e.g., changes in the environment. In this case, we can consider the behavior deterministic and verifiable with the methods of model checking.



Figure 2: Stable and Unstable Behavior.

In contrary, when a nearby behavior orbit is getting repelled from the given orbit of stability (stability point) (see Figure 2), we may consider such a behavior unstable and non-deterministic. Such a behavior cannot be verified with the traditional model checking, but through simulation and/or probabilistic model checking. In general, unstable behaviors may be perturbing in a trajectory asymptotically approaching the stable one or in a trajectory getting away from the stability point (see Figure 2). There may also be directions for which the behavior of the perturbed orbit is more complicated. e.g., neither converging nor escaping completely.

Various criteria have been developed to prove stability or instability of a behavior orbit (or trajectory). Under favorable circumstances, the question may be reduced to a well-studied problem involving *eigenvalues of matrices*. A more general method involves Lyapunov functions. In practice, any one of a number of different stability criteria are applied.

One of the key ideas in stability theory, and pursued in this approach, is that the qualitative behavior of a behavior orbit under perturbations can be analyzed using the linearization of the system near the equilibrium point. In particular, at each equilibrium of a self-adaptive system with an ndimensional phase space, there is a certain $n \times n$ matrix A whose eigenvalues characterize the behavior of the nearby points (Hartman-Grobman theorem) (Arrowsmith and Place, 1992). More precisely, if all eigenvalues are negative real numbers or complex numbers with negative real parts then the point is a stable equilibrium point (Lyapunov stability) and the nearby points converge to it at an exponential rate form a zone of asymptotical stability. If none of the eigenvalues are

purely imaginary (or zero) then the attracting and repelling directions are related to the eigenspaces of the matrix A with eigenvalues whose real part is negative and, respectively, positive, i.e., corresponding to unstable behavior.

4 KnowLang

KnowLang (Vassev and Hinchey, 2015a; Vassev and Hinchey, 2015b; Vassev and Hinchey, 2015c) is a framework for KR&R that aims at efficient and comprehensive knowledge structuring and awareness (Vassev and Hinchey, 2012) based on logical and statistical reasoning. Knowledge specified with KnowLang takes the form of a Knowledge Base (KB) that outlines a Knowledge Representation (KR) context. A key feature of KnowLang is a formal language with a multi-tier knowledge specification model (see Figure 3) allowing integration of ontologies together with rules and Bayesian networks (Neapolitan, 2013).



Figure 3: KnowLang Specification Model.

The language aims at efficient and comprehensive knowledge structuring and awareness. It helps us tackle (Vassev and Hinchey, 2015b): 1) explicit representation of domain concepts and relationships; 2) explicit representation of particular and general factual knowledge, in terms of predicates, names, connectives, quantifiers and identity; and 3) uncertain knowledge in which additive probabilities are used to represent degrees of belief. Other remarkable features are related to knowledge cleaning (allowing for efficient reasoning) and knowledge representation for autonomic behavior.

KnowLang employs special knowledge structures and a reasoning mechanism for modeling autonomic self-adaptive behavior (Vassev and Hinchey, 2015a; Vassev and Hinchey, 2015b). Such a behavior can be expressed via KnowLang policies, events, actions, situations and relations between policies and situations (see Definitions 1 through 10). Policies (Π) are at the core of autonomic behavior. A policy π has a goal (g), policy situations (Si_{π}) , policy-situation relations (R_{π}) , and policy conditions (N_{π}) mapped to policy actions (A_{π}) where the evaluation of N_{π} may eventually (with some degree of probability) imply the evaluation of actions (denoted $N_{\pi} \xrightarrow{[Z]} A_{\pi}$) (see Definition 6). A condition (n) is a Boolean expression over an ontology (see Definition 2), e.g., the occurrence of a certain event. Policy situations Si_{π} are situations (see Definition 7) that may trigger (or imply) a policy π , in compliance with the policy-situations relations R_{π} (denoted by $Si_{\pi} \xrightarrow{[R_{\pi}]} \pi$), thus implying the evaluation of the policy conditions N_{π} (denoted by $\pi \to N_{\pi}$)(see Definition 6). Therefore, the optional policy-situation relations (R_{π}) justify the relationships between a policy and the associated situations (see Definition 10).

$$\begin{split} \Pi &:= \{\pi_0, \pi_1, \dots, \pi_m\}, m \geq 0 \quad (policies) \quad (1) \\ A_{\pi} \subset A \quad (A_{\pi} - m) \\ policy actions; A - the set of all actions) \\ Si_{\pi} \subset Si \quad (Si_{\pi} - policy situations) \\ R_{\pi} \subset R \quad (R_{\pi} - policy - situation relations) \\ n &:= be(0) \quad (Boolean expression over ontology) \quad (2) \\ N_{\pi} &:= \{n_0, n_1, \dots, n_k\}, k \geq 0 \quad (policy conditions) \quad (3) \\ s &:= be(0) \quad (state) \quad (4) \\ g &:= \langle \Rightarrow s' \rangle | \langle s \Rightarrow s' \rangle \quad (goal) \quad (5) \\ \pi &:= \langle g, Si_{\pi}, [R_{\pi}], N_{\pi}, A_{\pi}, map(N_{\pi}, A_{\pi}, [Z]) > (policy) \quad (6) \\ N_{\pi} \stackrel{[Z]}{\to} A_{\pi} \quad (N_{\pi} \text{ implies the evaluation of actions } A_{\pi}) \\ Si_{\pi} \stackrel{[R_{\pi}]}{\to} \pi \to N_{\pi} \quad (Si_{\pi} \operatorname{trigger} \pi) \\ Si &:= \{si_0, si_1, \dots, si_n\}, n \geq 0 \quad (situation) \quad (8) \\ A_{si}^{\leftarrow} \subset A^* \quad (A_{si} - executed actions) \\ (A^{\leftarrow} - the set of all finite sequences with elements in A) \\ A_{si} \subset A \quad (A_{si} - posible actions) \\ E_{si}^{\leftarrow} \subset E^* \quad (E_{si}^{\leftarrow} - situation events) \\ (E^* - the set of all finite sequences with elements in E) \\ \end{split}$$

$$R := \{r_0, r_1, \dots, r_n\}, n \ge 0 \quad (relations) \quad (9)$$

$$r := < \pi, \ [rn], [Z], si > (relation) \quad (10)$$

$$si \in Si, \pi \in \Pi, si \stackrel{|Z|}{\to} \pi$$

Note that in order to allow for self-adaptive behavior, relations must be specified to connect policies with situations over an optional probability distribution (Z) where a policy might be related to multiple situations and vice versa. Probability distribution (Z) is provided to support probabilistic reasoning and to help the reasoner to choose the most probable situation-policy "pair". Thus, we may specify a few relations connecting a specific situation to different policies to be undertaken when the system is in that particular situation and the probability distribution over these relations (involving the same situation) should help the reasoner decide which policy to choose (denoted by $si \rightarrow \pi$) (see Definition 10). Hence, the presence of probabilistic beliefs (Z) in both mappings and policy relations justifies the probability of policy execution,

which may vary with time. Ideally, KnowLang policies are specified to handle specific situations, which may trigger the application of policies. A policy exhibits a behavior via actions generated in the environment or in the system itself. Specific conditions determine which specific actions (among the actions associated with that policy (see Definition 6) shall be executed. These conditions are often generic and may differ from the situations triggering the policy. Thus, the behavior not only depends on the specific situations a policy is specified to handle, but also depends on additional conditions. Such conditions might be organized in a way allowing for synchronization of different situations on the same policy. When a policy is applied, it checks what particular conditions N_{π} are met and performs the mapped actions A_{π} (map($N_{\pi}, A_{\pi}, [Z]$)) (see Definition 6). An optional probability distribution Z may additionally restrict the action execution. Although specified initially, the probability distribution at both mapping and relation levels is recomputed after the execution of any involved action. The recomputation is based on the consequences of the action execution, which allows for reinforcement learning.

States and goals drive the specification of any system modeled with KnowLang. States are used to specify goals (see Definition 5) and situations (see Definition 8), which on other side are used to specify policies (see Definition 6) often intended to provide self-adaptive behavior. The following is an example of a Boolean expression defining a state in KnowLang.

STATE ArrivedOnTime { carSafety.eCars.CONCEPT_TREES.Journey.STATES.Arrived AND (carSafety.eCars.CONCEPT_TREES.JourneyTime <=

carSafety.eCars.CONCEPT_TREES.Journey.PROPS.journeyTime)}

As shown, the logical expression above needs to be evaluated as true in order to consider the ArrivedOnTime state as active.

5 STABILITY POINTS WITH KnowLang

In this work, stability analysis works on states, goals, situations, and policies to determine behavior trajectories under perturbations in the execution environment. Here, the first task is to determine the *stable equilibrium points* (representing Lyapunov stability) and then the nearby points forming zones of *asymptotical stability* (see Section 3). The following elements are considered to be equilibrium points of stability in KnowLang specifications:

- 1) *departing states* (or expression of states) used in goal definitions;
- 2) *arriving states* (or expression of states) used in goal definitions (achieved goals);
- 3) policies handling situations to support goal achievement.

Note that policies are originally intended to handle situations that are considered to be critical for achieving a goal. Because these policies are part of the expected (deterministic) behavior, we consider as stable points both the execution of the policies' actions and handled situations. Here, the *asymptotical stability* Λ of an autonomous system modeled with KnowLang is a set of *areas of asymptotical stability* (see Definition 11). Here, each *area of asymptotical stability* is determined by zones defined by (see Definition 12):

- the goal through its goal states (zone ω_g see Definition 13);
- the goal-supporting policies through the execution of their actions (zones Θ_π see definitions 14 and 15);
- the critical situations (through their associated states) addressing that area of asymptotical stability, because they are deterministic (zones Θ_{si} see definitions 16 and 17).

$$\begin{split} &\Lambda := \{\lambda_0, \lambda_1, \dots, \lambda_m\}, m \geq 0 \quad (asymptotical stability) \quad (11) \\ &\lambda := < \varpi_g, \Theta_\pi, \Theta_{Si} > \quad (area \ of \ asymptotical \ stability) \quad (12) \\ &\varpi_g := < g(s \Longrightarrow s') > \quad (zone \ of \ goal \ stability) \quad (13) \\ &\Theta_\pi := \{\varpi_{0\pi}, \varpi_{1\pi}, \dots, \varpi_{m\pi}\}, m \geq 0 \quad (zone \ of \ policy \ stab) \quad (14) \\ &\varpi_\pi := < \Pi(g, A_\pi, Si) > \quad (zone \ of \ policy \ stability) \quad (15) \\ &\Theta_{Si} := \{\varpi_{0Si}, \varpi_{1Si}, \dots, \varpi_{mSi}\}, m \geq 0 \ (zone \ of \ stuat. \ stab) \quad (16) \\ &\varpi_{Si} := < Si(s) > \quad (zone \ of \ stuation \ stability) \quad (17) \end{split}$$

Here, the areas of asymptotical stability Λ are determined by the defined goals through their departing and arriving states. Note that in each area the zone of goal stability (see Definition 13) is determined by a specific goal g and the zones of policy stability Θ_{π} (see Definition 14) are basically formed by policies ϖ_{π} supporting that goal (see Definition 15). Finally, the zones of situation stability Θ_{si} are formed by situations associated with the policies forming the zones of policy stability. In this stability model, Λ defines a space of asymptotical stability and the continuous in time system is considered stable if at any time its behavior stays in that space or in a close proximity to that space. Here, if we consider that the system behavior *B* is determined by a sequence of actions *A* that starts out in the space of asymptotical stability Λ and stays in that space (or in a close proximity) forever, then B_{Λ} is Lyapunov stable near Λ .



Figure 4: KnowLang Asymptotical Stability.

Figure 4 demonstrate a Lyapunov stable behavior determined by a sequence of actions that form the *state trajectory* \widetilde{AB} . As shown, the exemplar space of asymptotical stability is formed by six areas of asymptotical stability (each driven by a distinct system goal g). Figure 5 depicts a possible area of asymptotical stability λ driven by a goal g that defines the zone of goal stability ϖ_g (a trajectory from a departing state to an arriving state), a zone of policy stability Θ_{π} presented as a set of policy trajectories each determined by a sequence of policy actions, and a zone of situation stability presented as a set of situation states.



Figure 5: KnowLang Area of Asymptotical Stability.

Note that when a *state trajectory* \overrightarrow{AB} goes through an area of asymptotical stability (see Figure 4), the presented behavior is considered stable if at any time it stays in one of the *zones of asymptotical stability* or stays in a close proximity to one of those zones. Here, the definitions for stability can be adapted to the KnowLang specifics as following:

An equilibrium solution f_{λ} (where all the evaluated behaviors are stable) to an *autonomous* system modeled with KnowLang is called:

- stable if for every small depart ε > 0 from the space of stability Λ, there exists a δ > 0 such that every solution f(t) having initial conditions within distance δ, i.e., ||f(t₀) f_λ|| < δ of the equilibrium remains within distance ε, i.e., ||f(t) f_λ_ρ|| < ε for all t ≥ t₀;
- asymptotically stable if it is stable and, in addition, there exists $\delta_0 > 0$ such that whenever $\delta_0 > ||f(t_0) - f_\lambda||$ then $f(t) \to f_\lambda$ as $t \to \infty$.

Here, one of the challenges in this solution that we still need to overcome is to determine the trajectory from a departing point to an arriving point (e.g., state transitions in goals). Further, we need to determine the borderline distances from stable states and deviations from stable trajectories ε , δ , and δ_0 . In this case, we need to break the state expressions down to atomic Boolean sub-expressions and determine sub-expressions that may vary in expression components without changing the overall expression evaluation.

Moreover, state-reduction techniques need to be developed to target the partitioning of the Boolean state expressions into sub-state expressions to determine parts that are irrelevant to the final result (TRUE or FALSE) and therefore can be excluded from the expression (replaced by TRUE or FALSE). Finally, a methodology for impact analysis to exclude low-impact partitions need to be developed, so the state space can be reduced.

6 CONCLUSION

An autonomous system is loaded with AI and operates in a potentially nondeterministic environment. Therefore, the verification of such systems needs to set boundaries that will provide the highest possible guarantees that the autonomy behavior will be safe and sound, so trust can be established in its innocuous operation. In this paper, we have presented our work on stability analysis for systems modeled with KnowLang. In this approach, a space of asymptotical stability is determined by analyzing the system goals along with the associated supportive policies and involved situations. The stable states and trajectories of action executions and state transitions form zones of asymptotical stability. Zones driven by a single system goal are grouped into areas of asymptotical stability and then these areas form a space of asymptotical stability. Here, if a system behavior determined by a sequence of actions that starts out in the space of asymptotical stability and stays in that space (or in a close proximity) forever, then that behavior is Lyapunov stable near that space of asymptotical stability.

Future work is considered with further development of this approach in terms of stability analysis automation, state trajectory definition and deviation borderlines. Moreover, state-reduction techniques need to be developed along with state impact analysis.

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