Comparison of Online Exploration and Coverage Algorithms in Continuous Space

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Abstract: We propose a framework to compare algorithms for multi agent exploration in an unknown continuous 2d environment. To analyze trade offs we compare algorithms with varying robot hardware requirements. We showcase our approach on Random Ballistic Walk (RBW), frontier-based exploration (The Next Frontier, TNF), Spiraling and Selective Backtracking (SSB), and Local Voronoi Decomposition (LVD). Algorithms that operate in a discrete grid-based space, such as LVD and SSB, are mapped to a continuous space for comparison with other algorithms. To our knowledge, no other extensive comparison of these exploration algorithms operating under the same testing environment has been conducted. The algorithms are tested in a custom 2D physics-driven simulation (Multi Agent Exploration Simulator, MAES), with two types of maps, namely the Cave map (C-Map) and the Building map (B-Map). The performance of each algorithm is evaluated in terms of coverage and exploration of the best in terms of coverage in all tested scenarios. TNF performed the best in terms of exploration, especially on bigger maps. RBW achieved good results in terms of both coverage and exploration in C-Maps, but not in B-Maps. LVD performed similarly to RBW in C-Maps, but better in the B-Maps.

1 INTRODUCTION

Collective terrain exploration and coverage using swarm robotics has many applications such as search & rescue and surveillance, and in fact it keeps a prominent role among the use cases for multi agent algorithms (Schranz et al., 2020; Brambilla et al., 2013).

The problem of terrain coverage by multiple agents comprises Offline Terrain Coverage and Online Terrain Coverage (OTC), where *online* refers to the terrain being unknown beforehand and thus explored while the algorithm is executed (Agmon et al., 2008). In this paper we focus on OTC, and we argue that many great algorithms have been proposed for terrain coverage using swarm robotics, however they are hardly comparable since they model their agents in very different manners; moreover, we raise the issue of the limited realism introduced by the requirements of some of the algorithms in terms of robot capabilities, such as (i) the environment being con-

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sidered a grid-like structure that can be maneuvered by moving from cell to cell (Cheraghi et al., 2020; Albani et al., 2019) where (ii) the robots cannot suffer from colliding with each other; (iii) communication between robots has unlimited range and is not blocked by walls (Kambayashi et al., 2009; Oikawa et al., 2015); (iv) robots know each other's location, not impeded by walls (Kegeleirs et al., 2021), and are provided with real-time distributed Simultaneous Localization And Mapping (SLAM) (Gonzalez and Gerlein, 2009a).

We propose a framework, and its accompanying tool Multi Agent Exploration Simulator (MAES), for taking algorithm simulations one step closer to the real world by representing their movement space as a continuous 2D plane, rather than a grid of cells (see Figure 1), and for mapping the different approaches on a common ground in terms of hardware capabilities of the robots.

To showcase the approach, we first of all (Section 2) provide information regarding selected algorithms, namely Spiraling and Selective Backtracking (SSB) from (Gautam et al., 2018), Local Voronoi Decomposition for Multi-Agent Task allocation (LVD) from (Fu et al., 2009), The Next Frontier (TNF) from (Co-

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Figure 1: Left: An agent moving in continuous space in MAES. Right: a traditional grid based simulation.

lares and Chaimowicz, 2016), and Random Ballistic Walk (RBW) from (Kegeleirs et al., 2019), the latter being used as a baseline for the comparisons.

Section 3 delves into the mapping of the algorithms onto a continuous space, and other design decisions needed to map the algorithms on a common comparison framework.

Section 4 describes the MAES tool¹, which was developed to demonstrate our approach. It is based on the Unity physical engine (Unity Technologies, 2021), and allows to simulate the algorithms on building maps (B-Map) featuring hallways and rooms, and cave maps (C-Map) full of irregular shapes. The tool allows for visual inspection for easier debugging and comparison of implemented algorithms, and draws comparisons in terms of coverage and exploration, which we define in Section 2.1.

Results of simulations performed on both B-map and C-map are reported in Section 5, and Section 6 draws conclusions and proposes future work, both in terms of extension of the MAES simulator, and regarding the application of our framework to more scenarios.

Summarizing, the main contributions of this paper include:

- A mapping SSB, LVD and TNF to continuous space and on common hardware capabilities (Section 3)
- MAES, a simulation tool for testing algorithms in continuous space and on common hardware capabilities. (Section 4)
- A comparison of the performance of the mapped algorithms in a realistic simulator. Since the algorithms have different hardware requirements, this also provides insight into the benefit of extra hardware capabilities. (Section 5)

2 BACKGROUND INFORMATION

This section prepares the scene for the rest of the paper, by defining a few key concepts formally, and describing and comparing the algorithms we focus on.

2.1 Definitions

Definition 1. A grid space G is a movement space where for any given position $\mathbf{x} \in G$, $\mathbf{x} \in (\mathbb{Z} \times \mathbb{Z})$.

Movement in a grid space consists of discrete steps from cell to cell. A step can only move an agent from its current cell to one of its immediate neighbours. There are four or eight possible movement directions, depending on whether diagonal steps are allowed.

Definition 2. A Continuous space C is a movement space where for any given position $\mathbf{x} \in C$, $\mathbf{x} \in (\mathbb{R} \times \mathbb{R})$

In a Continupus space, if an agent moves a distance of d in any direction between time t_1 and t_2 , it will passy by uncountably infinitely many possible positions **x** in that time range.

Definition 3. An agent's physical body has an area that covers (or 'shadows') some movement space, and **Coverage** describes the amount of accumulated area that an agent's body-area has physically covered during the process of moving through a movement space.

Definition 4. An agent has a field of view (FOV) which covers (or "shadows") a circular area of the movement space. **Exploration** describes the amount of accumulated open area an agent's FOV has covered during the process of moving through a movement space.

Definition 5. Simultaneous localization and mapping (SLAM) is constructing and updating a map of explored area in a movement space, while simultaneously keeping track of an agent's position within it.

Definition 6. Environment tagging refers to leaving information at some position in the environment, that can be used for indirect communication with other agents.

2.2 Algorithms

2.2.1 Random Ballistic Walk

Random walk describes a category of approaches where agents randomly alternate between rotating and moving straight ahead. Random Walks require that the agents be able to detect collision, move straight ahead and rotate in place. Some variations also require the ability to estimate how far they have travelled since the previous rotation.

¹A video demonstration of MAES can be found at https://youtu.be/lgUNrTfJW5g. MAES is open source and the source code can be found at https://github.com/MalteZA/MAES.

In (Kegeleirs et al., 2019), several variations on random walks are compared including Brownian motion, Lévy Walk, and RBW. In RBW, the straight movement continues until the agent collides with either an obstacle or another agent. According to (Kegeleirs et al., 2019) RBW produce the best results for coverage and mapping of unknown environments, and thus RBW was selected as the baseline for the comparison with more sophisticated algorithms.

2.2.2 Spiraling and Selective Backtracking

In the SSB (Gautam et al., 2018) algorithm, the agents perform coverage by traversing sections of the map following inward spiral patterns.

The SSB algorithm is designed for a grid space, and movement is modelled as discrete steps between cells. SSB also assumes that agents can communicate globally. Finally, the algorithm design is rooted in the assumption that agents are able to accurately build a map of their environment and continually synchronize this map across all agents. These two requirements reduces the usability in real world scenarios.

In the SSB algorithm agents alternate between two states: spiraling and backtracking. In the spiraling state, agents follow a simple set of rules to traverse an area in the pattern of an inward spiral. While spiraling, the agents make note of neighbouring unexplored cells as potential back-tracking points. Once spiraling is completed the agent broadcasts its collected backtracking points to all other agents. Then, the agent determines which backtracking point to choose as its target by performing an auction where agents bid on available backtracking points. Each agent calculates a bid based on its distance from the backtracking point. If an agent is in spiraling mode during an auction, its bid values will also depend on the estimated number of steps required to finish the ongoing spiral. If the agents wins a backtracking point, it will travel to that point and then begin the spiraling phase again. If the agent does not win any backtracking points it will still travel to the closest candidate to avoid being idle.

2.2.3 Local Voronoi Decomposition

The idea behind LVD (Fu et al., 2009) is to achieve task allocation using only indirect communication in the form of the position of visible agents and obstacles.

Each agent computes Voronoi regions from the areas of the map within line of sight (usually 360 degrees around the agent) at a given time, using other agents' location and walls/obstacles within line of sight. The visibility range is not mentioned in the original paper, but the illustrations in (Fu et al., 2009) seem to indicate that it is infinite, so we will in this paper assume infinite visibility range for LVD.

In LVD, an occlusion point is defined as a corner or obstacle where the line of sight is broken. All LVD agents employ local SLAM to detect areas that they have themselves previously covered, as well as previously covered occlusion points.

Pseudo code of the LVD algorithm can be seen in Listings 1 and 2. The agents see the map as a grid of tiles, which are *explored* when an agent has visited it². A tile can be within view but not explored yet.

Listing 1: LVD Algorithm 1 - Divide and Conquer.

11	Divide part
Foi	each cell in view
	${\tt if}$ the cell is closer to agent than to other agents
	mark the cell as within region
	end if
enc	for
//	Conquer part
Ĺ£	there is just one unexplored cell within region
	move to it
el:	se if there are more than one unexplored cell
	move in an ordered list, e.g. [North, East, South, West]
1	se
	Enter Search Mode
enc	lif
	to 1
go	Listing 2: LVD Algorithm 2 - Search Mode.
go	Listing 2: LVD Algorithm 2 - Search Mode.
go //	Listing 2: LVD Algorithm 2 - Search Mode.
]0 /// if	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then
// //	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1
yo /// if	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 a if
go /// if enc if	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been
go 	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode
jo /// if if →	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been
jo /// if if →	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been \leftrightarrow visited
yo /// if enc if →	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been \leftrightarrow visited se if all occlusion points have already been visited
jo // if enc if →	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been \leftrightarrow visited se if all occlusion points have already been visited move to the least recently visited occlusion point
go /// if enc if ⇒ el: el:	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been ↔ visited se if all occlusion points have already been visited move to the least recently visited occlusion point se if there are no occlusion points or if the only occlusion
go go if if enc if ecl ecl if ecl if ecl ecl if if ecl if ecl if if ecl if if ecl if if if ecl if if if if if if if if	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been ↔ visited se if all occlusion points have already been visited move to the least recently visited occlusion point se if there are no occlusion points or if the only occlusion point has just been visited
<pre>go</pre>	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been ↔ visited se if all occlusion points have already been visited move to the least recently visited occlusion point se if there are no occlusion points or if the only occlusion point has just been visited move to the nearest Voronoi boundary which does not
go go if if enc if el : el : \leftrightarrow	Listing 2: LVD Algorithm 2 - Search Mode. Consider all occlusion points within view. at any time an unexplored cell appears within region then exit Search Mode and go back to Algorithm 1 if there is at least one occlusion point which has not been visited during the Search Mode move to the nearest occlusion point which has not been ↔ visited se if all occlusion points have already been visited move to the least recently visited occlusion point se if there are no occlusion points or if the only occlusion point has just been visited move to the nearest Voronoi boundary which does not coincide with any obstacle

In its *Divide and Conquer Mode*, the algorithm considers that agents sense other agents close by and then use this information to divide the currently visible tiles into Voronoi regions, i.e. always delegating a given tile to the agent closest to it. Each division is computed independently by the agent. If any tiles within a given agent's Voronoi region are unexplored, the agent will move to them in a configurable order, e.g. first north, then east, south and west. If no tile is unexplored, the agent enters the *Search Mode*.

In the *Search Mode* the agent considers occlusion points within line of sight, and it moves to an occlu-

²Thus, LVN definition of explored correspond to what this paper defines as covered.

sion point that is likely to lead to an area with unexplored tiles. If all occlusion points have been visited during the current phase of *Search Mode*, the agent moves to the least recently visited one. If no occlusion point is within line of sight, or they have all been visited recently, the agent moves to the nearest Voronoi boundary which does not coincide with an obstacle, in order to search for new unexplored tiles, while also possibly "pushing" other agents further out and away from the already explored area. If at any time during the *Search Mode* an unexplored tile comes into line of sight, the agent exits *Search Mode* and returns to the *Divide and Conquer Mode*.

2.2.4 The Next Frontier

TNF (Colares and Chaimowicz, 2016) is a frontierbased exploration algorithm that works by detecting and moving to boundaries between explored and unexplored areas (frontiers). TNF prioritizes frontiers through a combination of factors that are combined in a utility function

$$U(f) = Inf(f) + Dist(f) - Coord(f)$$
(1)

where f is a frontier, Inf(f) is an Information Factor (prefer higher concentrations of cells of interest), Dist(f) is a Distance Factor (prefer frontiers at a certain configurable distance away), and Coord(f) is a Coordination Factor (dismiss frontiers that are close to other agents). TNF agents can detect and map out the local environment through SLAM, and can merge maps with other agents when they are within communication distance. The TNF paper does not mention any specific range requirements in regard to communication and area detection.

The Information Factor is used to prioritize cells that are likely to contribute more valuable information by being explored. An uncertain cell sitting in a large concentration of other uncertain cells in a frontier should be more valuable to explore than a single uncertain cell in a completely discovered area. A cells occupancy-value is seen as a number between 0 and 1, where 0 is a free cell, 0.5 an unexplored cell and 1 an occupied cell, and the cell's contribution to the Information Factor is modeled as a non-normalized Gaussian function centered on 0.5. The Information Factor of a frontier f can then be calculated as

$$\operatorname{Inf}(f) = \sum_{c \in f} \left(\mathbf{G}(c) + \sum_{n_i \in N} \mathbf{G}(n_i) \right)$$
(2)

where G is the Gaussian function mentioned previously, N is all neighbours of cell c, and n_i is the *i*th neighbour of c.

The Distance Factor is based on the idea that an agent might prefer to prioritize frontiers at a certain

distance over others. The distance of a frontier is calculated using what is referred to as a "wavefront", starting from the position of the agent going towards the frontier. This distance is then normalized to a value (a real number) between [0..1], and the final Distance Factor is calculated as

$$\text{Dist}(f) = \text{wavefront}(f)^{\alpha - 1} \times (1 - \text{wavefront}(f))^{\beta - 1}$$
 (3)

where α and β are the configuration parameters used to control at which distance a frontier is favored. The TNF authors found the most promising results using the configurations ($\alpha = 3, \beta = 9$), ($\alpha = 5, \beta = 9$), or ($\alpha = 8, \beta = 8$).

The Coordination Factor is the final part of the Utility Function, where an agent will prefer to explore frontiers that are relatively far away from neighbouring agents. Once again, a Wavefront Function is used, but this time the starting position is of a neighbouring agent, and this Coordication Factor is calculated for each neighbour the agent can see at the time of calculation.

2.2.5 Requirements Summary

Table 1 lists the agent capability requirements for each of the algorithms. The Environment Tagging and SLAM criteria for LVD are parenthesized because the functionality required by LVD can be achieved using either one, as described in subsection 2.2.3.

Table 1: Requirements for each algorithm.

	Collision Det.	Object Det.	Env. tagging	SLAM	Distri. SLAM	Local Com.	Global Com.
RBW	1						
LVD	1	1	(\checkmark)	(\checkmark)			
TNF	\checkmark	\checkmark		\checkmark	1	1	
SSB	1	1		1	1	1	1

2.3 Related Work

The simulation of multiple agent coverage algorithms in realistic simulation settings is a relatively unexplored research area. Only few papers can be found in the bibliography on this topic.

In (Gautam et al., 2018) a comparison of coverage performance is provided for existing algorithms with reduced communication range. This paper compares SSB, Multiple-Depth-First-Search (MDFS), Brick & Mortar (BNM), Boustrophedon and backtracking mechanism (BOB) and Backtracking Spiral Approach - Cooperative Multi Robot (BSA-CM). SSB achieved the best results. The comparison is, however, performed in a grid base simulation using only two different maps - a basic and a Cluttered map.

In (Gautam et al., 2021), SSB, BOB, and BSA-CM are compared with different communication range restrictions and number of agents. The algorithms are compared in terms of coverage and redundant coverage using hardware FireBird V robots on a small plane. SSB and BSA-CM performs similarly and better than BOB in terms of avoiding redundant coverage.

3 MAPPING TO CONTINUOUS SPACE

LVD and SSB operate in a grid space, where the cell size is roughly equal to the size of the agent, and this section describes a mapping from continuous space to grid space for them. Moreover, other issue can appear when mapping the algorithm to a more realistic framework, such as by considering that agent can actually collide with each other. Finally, TNF leaves some design parameters to be defined , and this section takes care of this detail too.

Both LVD and SSB use SLAM, and the mapping can be implemented by overlaying a grid on top of their environment map in the Continuous Space. Each tile in the overlay is considered either open or solid, with the entire tile considered solid if any part of a solid object is detected within the bounds of the tile. For SSB, this grid overlay and accompanying tile statuses are synchronized along with the map itself.

The above approach limits maneuverability, as otherwise traversable areas may be marked as solid, depending on the alignment of the grid. An illustration of this problem can be seen in Figure 2 where (a) shows the environment map constructed by an agent. The agent could comfortably fit through both the west and east openings. However, as illustrated in (b), the alignment of the grid can change perceived traversability of the openings, if the tiles are too big.

To guarantee that the grid will always provide an opening through a passage, the passage must be at least as wide two tiles, as illustrated in (c). Similarly, to guarantee open tiles in diagonal passages, the passages must be at least twice as wide as the diagonal of a tile. This approach also reduces the area that will be covered, as some free area will be perceived as solid by the agent, thus not needed to be covered.



Figure 2: Illustration of grid overlay. (a) An agent (red) in a Continuous Space with obstacles (black). (b) A grid overlay with large tiles. Red tiles are perceived as solid and grey tiles as traversable. (c) A more fine grained grid overlay.

3.1 Selective Spiraling and Backtracking

SSB is built upon an algorithm called BSA-CM (Gonzalez and Gerlein, 2009b). Apart from the aforementioned grid mapping approach, BSA-CM assumes collision avoidance to be trivial, as all agents can move at the same time in discrete steps between cells. In our simulation, on the other hand, coordination and timing issues arise, as an agent may partially occupy up to four cells at the same time. To avoid collisions, we introduce a tile reservation system, where agents will secure a reservation for a tile before moving into it. Reservation requests are broadcast to nearby agents, and after a time delay, the reservation is assumed to be accepted. To avoid delaying movement unnecessarily while waiting for reservations, agents in our implementation will attempt to predict their path, and reserve several tiles in advance.

Multiple agents may attempt to reserve the same tile at the same time, in which case the conflict is resolved by accepting the reservation from the agent with the highest id.

3.2 Local Voronoi Decomposition

The mapping of LVD to continuous space leads to a series of challenges, which are not present in the original paper since it uses Grid Space.

3.2.1 Empty Voronoi Regions

In a grid, an agent's Voronoi region will never be empty in terms of tiles, since the tile on which the agent stands is always contained in the agent's Voronoi region. In continuous space, however, an agent can be positioned in-between other agent that are closer to the center of tiles and thus own them, see Figure 3. We solved this issue by having the center agent stay idle, since the surrounding agents will have their respective Voronoi regions reach out from each of the four corners and thus eventually move away and set the center agent free.



Figure 3: Illustration of an agent with an empty voronoi region.

3.2.2 Calculation Interval

In LVD the agents move in a grid using discrete movement, and recalculate the Voronoi regions after every discrete move. However, this is not possible in continuous space, as a change in position may be infinitesimally small. This begs the question of when to recalculate Voronoi regions and find the next task to complete. We decided that an agent would recalculate whenever it collides with anything, or if it reaches the tile it was heading for.

Additionally, LVD mandates, that an agent should exit *Search Mode* whenever an unexplored tile comes into line of sight. It would be unfeasible to look for unexplored tiles at every infinitesimally small step of movement. We decided to recalculate periodically at a given interval, even if they have not collided or reached their current target. The interval is shorter while in *Search Mode* to make it likely that the agent exits *Search Mode* as soon as an unexplored tile comes into line of sight.

3.2.3 Occlusion Points with Multiple Visits

In *Search Mode* the agent first explores all occlusion points within line of sight. In continuous space with a high resolution slam map it can, however, be difficult to keep track of which occlusion points have been visited. When visiting occlusion points, the agent heads



Figure 4: (a) An agent identifying the same occlusion point as two distinct points (b) Two points are considered the same point, and the agent will not revisit the occlusion point

for the closest occlusion point it sees, which can however result in multiple visits to the same occlusion points. Consider Figure 4(a), where the agent sees an occlusion point and thus moves to this point and around the obstacle. Now, however, the agent sees the same occlusion point but from a different perspective. Due to the shape of the obstacle the previous occlusion now appears to be a new occlusion point, since it is slightly offset to the right. This causes the agent to revisit an already visited occlusion point. In order to avoid this issue, we consider all occlusion points within some distance of each other to be the same (see Figure 4(b)).

3.3 The Next Frontier

TNF was originally designed to operate in a continuous space. As mentioned in Section 2.2.4, TNF sees the environment as a set of cells, where the boundaries between the fully explored cells and the unknown cells make up the frontiers to move to. However, a design decision is required regarding how to implement the components of the Utility Function.

In order to prioritize cells where an agent is the least certain of the cell's state (i.e. cells with an Information Factor contribution close to 0.5), the authors of (Colares and Chaimowicz, 2016) made use of a non-normalized Gaussian function centered around 0.5. However, the authors only mention the center of the function's peak, and do not specify its height nor its deviation. For our implementation, we assume a function peak of 5 and a deviation of 0.1, as this seems to produce a usable behavior.

With regards to the Distance Factor, the TNF authors mention the use of a wavefront function to measure the distance from an agents position to the cells of a frontier, without providing much details on how to build it. For our implementation, we assume euclidean distance (L^2 norm) to each of the cells, as this is consistent with how a wavefront is depicted in the paper. With regards to the normalization of the distances to the frontiers to values between [0, 1], we assume that a wavefront is normalised by identifying the most distant cell's distance D_{MAX}, and dividing the distance of every other cell by D_{MAX}. As mentioned in Section 2.2.4, the final Distance Factor value is configurable with the variables α and β . For this implementation, we chose to use one of the authors' suggested configuration of $(\alpha = 8, \beta = 8)$.

We implemented the *Coordination Factor* by summing up the results for each neighbouring agent. We acknowledge that other approach might be viable, for example dividing each wavefront distance by the number of known neighbours. Finally, to take into consideration collisions between two (or more) agents, we use a mitigation-procedure that implies that, after a collision, each agent looks in its immediate surroundings, and go to a location that is free and not occupied by another agent.

4 THE MAES TOOL

This section describes the MAES tool, which we developed for conducting the experiments shown in Section 5. MAES is a deterministic 2D discrete timestep physics-based simulation, visualized in 3D. The simulator uses the Unity Engine (Unity Technologies, 2021), for visualization and physics simulation. Physics are simulated at a rate of 100 ticks per simulated second. The *reaction time* of the agent, i.e. how often the agents apply their exploration algorithms to the current data, is 10 ticks.

MAES features a map generator than can generate two types of maps, i.e. the cave map type (C-Map) and the Building map type (B-Map). MAES provides an agent control interface that allows implementation of exploration algorithms. This interface provides access to movement controls, communication, object detection and SLAM. MAES allows many of its feature to be configured, and it takes parameters for agent constraints, for physics simulation, and for map generation. A list of all possible parameters for the simulator can be seen on https://github.com/MalteZA/MAES, where the code is made available to the public.

4.1 Environment Maps

In order to study useful settings, MAES generates random maps fitting the characteristics of B-maps and Cmaps.

A map consists of a number of tiles, where each tile consists of 8 triangles arranged, as shown in Figure 5(a). An agent can be up to as big as a tile. In Fig-



Figure 5: The tile structure of the map. (a) The arrangement of the of 8 triangles forming a single tile. (b) The size of an agent relative to tile. (c) Four tiles, where solid triangles form an obstacle.



Figure 6: An example of a 50x50 generated cave map.

ure 5(b) an agent with size 0.6 times a tile is shown. Finally, each triangle may be open or solid, where open tiles are freely traversable and solid triangles are impassable. Obstacles with complex shapes can be formed by chaining solid triangles as shown in 5(c).

4.1.1 Cave Map Type

The C-map type is relevant for example for the use cases of exploration and search & rescue in an irregularly shaped environment. The C-map generator starts by randomly distributing wall and room type cells in a grid of a given size. It then uses neighbour smoothing, based on rules similar to Conway's Game of Life, to create sections of walls and rooms, e.g. a tile with 4 or more solid neighbour tiles will turn solid and the other way around. Now the generator can interconnect all rooms (groups of open tiles) created while ensuring that every part of the cave is reachable from every other part. In order to make the cave more realistic and irregular we use the Marching Squares algorithm to round off edges. An example image of a C-map can be seen in Figure 6.

4.1.2 Building Map Type

The B-map type is relevant for simulating use cases such as floor cleaning and search & rescue in buildings. The generated maps are made to look like the floor plan of a building. The B-map generator creates a grid of a given size and marks each cell as being either wall, hallway or room. The generator starts by



Figure 7: An example of a 50x50 generated building map.



Figure 8: The 3D model of the MONA-inspired agent used in MAES.

connecting all the hallways, then it recursively splits the spaces between to hallways to create rooms, with room size defined at random. Finally all the rooms are connected to each other with doors in such a way that all rooms are reachable from the hallways. Figure 7 shows an example of a B-map.

4.2 Agents

Within MAES, an agent's capabilities can vary depending on the parameters that describe the simulation at hand, e.g. limited vision, broadcasting range etc. An image of the 3D model of an agent, inspired by the MONA robots (Arvin et al., 2019), can be seen in Figure 8.

4.2.1 Movement

An agent is able to rotate in place and moving straight ahead, and the agent is not able to rotate while moving forwards. Movement is simulated through the Unity 2D Physics Engine by applying force at the position of each wheel. This simulation accounts for inertia, drag, and collisions with obstacles and other agents. The agents of the MAES simulator can reach a top speed of 3 tiles / second (or 10 logic ticks). The tile size can vary depending on the scale of the map. We decided that the agent takes about 30 physics ticks to reach its top speed. Drag is a function of speed, which in combination with inertia results in non-constant acceleration, leading to an agent reaching half its top speed (about 1.5 tiles / second) after just 4 ticks.

4.2.2 Sensors and Communication

Agents are able to sense other agents at a given distance, which is a simulation parameter. In order to accommodate variety of scenarios with differing hardware capabilities, the signal sensing other agents can also be configured to be blocked by walls, i.e. requiring line of sight. Agents detect collisions with walls and with other agents, and can detect a nearby wall and the angle to said wall. This could for example be achieved using a LIDAR scanner in the real world. Agents can communicate through broadcasting, and both communication range and the capabilities to pass walls is defined via simulation parameters. Line of sight is determined using ray tracing. Beyond line of sight and maximum range, no other signal loss is simulated.

Finally, agents can drop tags on the ground to deposit information in the environment and communicate indirectly with other agents, as required for example by LVD (Section 2.2.3). Tags can only be dropped at an agent's current position, but data can be written to and read from at a configurable distance.

4.2.3 Simulated SLAM

If enabled, the agent can provide an environment map generated via SLAM to the algorithm being simulated. SLAM is simulated by performing a series of ray traces from the position of the agent, and measuring the distance that the rays traveled before colliding. This emulates the behavior of technologies such as LIDAR scanners.

The agent continuously constructs a "SLAM map" using the ray tracing information when it becomes available. If any object is detected within the region of a tile, then that entire tile is marked as solid. If a ray trace is sent in the direction of a tile, and no object is found, the tile is assumed to be open, unless previous traces indicate that it is solid. The agent has access to an approximation of its location within the SLAM map. The simulation can be configured to automatically synchronize SLAM maps of agents that are within communication range of each other.

4.2.4 Interfaces

MAES is intended to support the implementation of many different algorithms. For this reason we expose an interface for the algorithms to control the agents and provide access to sensor information. The interface is created to allow for implementing all of the hardware requirements mentioned in Table 1, e.g. SLAM map, environment tagging, etc.

4.3 Debugging Features

As MAES should function as a testbed for many different algorithms we include a wide variety of debugging tools. A menu is included for controlling the camera view over the simulation, as well as changing the simulation speed. Additionally, agents can be individually selected, which makes the camera follow the agent as well as reveal debugging information in a side bar regarding the selected agent.

Furthermore, slam maps, communication, and environment tagging can be visualized. When a simulation is running the surface of the map is highlighted in green if any agent at any point has explored it. If an agent is selected, the surface reveals in blue the tiles included in the slam map for said agent. The slam map can also include sections of the map revealed by other agents, if slam synchronization is enabled and the two agents have been within communication distance of each other. Environment tagging is visualised using colored boxes on the ground where an agent has tagged the environment.

5 EXPERIMENTS AND RESULTS

For quantifying performance we use two metrics: coverage and exploration, explained in Definition 3 and Definition 4 respectively. The coverage metric is important for purposes like robotic lawn mowers and cleaning robots such as autonomous vacuum cleaners, where the robots need to efficiently and precisely cover areas. The exploration metric is important for agents when exploring an area to generate a navigable map of the environment, or in search & rescue scenarios where it is vital to quickly explore large areas.

As described in Section 4.1 a MAES map consists of tiles, where each tile consists of 8 triangles. Exploration is measured in triangles, and a triangle is considered to be explored when an agent's simulated Lidar trace has intersected with the triangle. Coverage is measured in whole tiles, and a tile is considered covered once the center point of an agent has entered the bounds of the tile.

For both B-map and C-map, in order to test the scalability of the algorithms, we simulated different map sizes, namely 50x50, 100x100 and 200x200. We used an agent size of 0.6 tiles. We run each algorithm 20 times for 60 simulated minutes, each with a different random seed, leading to different maps. If an algorithm reaches a coverage of 99.5% (or TNF reaches an exploration of 99.5%), the simulation ends before the 60 minute mark. Each algorithm has its own agent constraints as showed on Table 1. For example, SSB assumes global communication, while LVD does not need direct communication at all.

5.1 Building Scenario

According to Figure 9, LVD leads in exploration for the first few minutes after which SSB takes the lead

Average coverage/exploration for 50x50 building maps



Figure 9: Comparison of algorithm performance in terms of exploration and coverage for the building type map of size 50x50. Average results of 20 random seeds.

in the C-map with a size of 50x50. In terms of coverage, SSB performs the best and has the map fully covered on average after about 5 minutes. RBW and LVD achieve similar coverage with a slight advantage to LVD.

Figure 10 shows the results for the 200x200 map can be seen. Here TNF achieves the best result in terms of early exploration. SSB, however, eclipses the exploration performance of TNF after about 30 minutes. LVD and RBW are significantly behind on exploration and neither manages to finish within 60 minutes. LVD achieves a slightly better result than RBW in terms of exploration. SSB achieves the fastest coverage and finishes on average after about 37 minutes. RBW and LVD reaches about 56% and 64% coverage respectively after 60 minutes.

We omit the results for 100x100 B-maps, since they were intermediate between the 50x50 and the 200x200 B-maps.



Figure 10: Comparison of algorithm performance in terms of exploration and coverage for the building type map of size 200x200. Average results of 20 random seeds.

Average coverage/exploration for 50x50 cave maps



Figure 11: Comparison of algorithm performance in terms of exploration and coverage for the cave type map of size 50x50. Average results of 20 random seeds.

5.2 Cave Scenario

In Figure 11 it can be seen, that LVD explores the fastest, followed by RBW and then TNF and SSB. In terms of coverage SSB is slightly faster with RBW and LVD both finishing last after about 6.5 minutes.

The results for the 200x200 C-map, seen on Figure 12, has TNF exploring the map the fastest, followed by SSB. LVD and RBW achieves similar exploration rates. SSB covers the larger map significantly faster than the other algorithms. Neither RBW nor LVD finishes the coverage within the 60 minutes of the simulation. RBW and LVW both manage about 83% coverage.

We omit the results for 100×100 C-maps, since they were intermediate between the 50×50 and the 200×200 C-maps.

Average coverage/exploration for 200x200 cave maps

100 80 <u>60</u> 80



Figure 12: Comparison of algorithm performance in terms of exploration and coverage for the cave type map of size 200x200. Average results of 20 random seeds.

Minutes

6 CONCLUSIONS AND FUTURE WORK

The paper proposes a continuous space mapping for the SSB, LVD and TNF algorithms, and presents the open source simulation tool MAES, which can be used for developing and evaluating exploration algorithms in a continuous 2D space. The mapped algorithms were implemented and compared statistically through repeated runs with randomly generated maps. RBW was implemented as a baseline for comparison with other the algorithms.

In terms of coverage, SSB performed significantly better than the other algorithms in both B-map and Cmap types. This performance comes at the cost of reduced realism, as SSB assumes global communication and distributed SLAM capabilities.

TNF shows good results for exploration, exceeding other algorithms in large maps for both B-map and C-map types. TNF does, however, also require distributed SLAM to operate.

LVD achieves slightly better results in terms of coverage and exploration in most maps than RBW while only using strictly local information.

Future work will consider to implement more state-of-the-art algorithms, and propose algorithms tailored on the realistic framework we propose. Moreover, there are aspects of the algorithms that can be further studied such as scalability with respect to the number of agents, and the effect of lower visibility and transmission ranges. Finally, the MAES tool can benefit from more development, such as to allow for batch (no UI) and distributed simulations for enhanced performance.

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