

Bi-Objective CSO for Big Data Scientific Workflows Scheduling in the Cloud: Case of LIGO Workflow

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Abstract: Scientific workflows are used to model scalable, portable, and reproducible big data analyses and scientific experiments with low development costs. To optimize their performances and ensure data resources efficiency, scientific workflows handling big volumes of data need to be executed on scalable distributed environments like the Cloud infrastructure services. The problem of scheduling such workflows is known as an NP-complete problem. It aims to find optimal mapping task-to-resource and data-to-storage resources in order to meet end user's quality of service objectives, especially minimizing the overall makespan or the financial cost of the workflow. In this paper, we formulate the problem of scheduling big data scientific workflows as bi-objective optimization problem that aims to minimize both the makespan and the cost of the workflow. The formulated problem is then resolved using our proposed Bi-Objective Cat Swarm Optimization algorithm (BiO-CSO) which is an extension of the bio-inspired algorithm CSO. The extension consists of adapting the algorithm to solve multi-objective discrete optimization problems. Our application case is the LIGO Inspiral workflow which is a CPU and Data intensive workflow used to generate and analyze gravitational waveforms from data collected during the coalescing of compact binary systems. The performance of the proposed method is then compared to that of the multi-objective Particle Swarm Optimization (PSO) proven to be effective for scientific workflows scheduling. The experimental results show that our algorithm BiO-CSO performs better than the multi-objective PSO since it provides more and better final scheduling solutions.

1 INTRODUCTION

Scientific applications handling big volumes of data usually consist of a huge number of computational tasks with data dependencies between them, which are often referred to as scientific workflows. They are frequently used to model complex phenomena, to analyze instrumental data, to tie together information from distributed sources, and to pursue other scientific endeavors (Deelman, 2010). Data intensive scientific workflows need from few hours to many days to be processed in a local execution environment. In this paper, we consider the scientific application of Laser Interferometer Gravitational Wave Observatory (LIGO) ¹ which is a network of gravitational-wave detectors, with observatories in Livingston, LA and Hanford, WA. The observatories' mission is to detect and measure gravitational waves predicted by general relativity Einstein's theory of gravity in which grav-

ity is described as due to the curvature of the fabric of time and space. The data collected by the different LIGO observatories is about Tera-octets per hour. To analyze this huge amount of data, scientists modeled their workloads using scientific workflows. The LIGO workflow may contain up to 100 sub-workflows and 200,000 jobs (Juve et al., 2013) which requires many hours to be performed in a Grid computing environment. To obtain simulation results within acceptable timeframes, large scientific workloads are executed on distributed computing infrastructures such as grids and clouds (Brewer, 2000). The Cloud Computing environment offers on-demand access to a pool of computing, storage and network resources through Internet that allows hosting and executing complex applications such as scientific workflows. As a subscription based computing service, it provides a convenient platform for scientific workflows due to features like application scalability, heterogeneous resources, dynamic resource provisioning, and pay-as-you-go cost model (Poola

¹<https://www.ligo.caltech.edu/>

et al., 2014). Scheduling data intensive scientific workflows on the Cloud environment is usually an NP-complete problem (Guo et al., 2012). It aims to find convenient resources for executing the workflow tasks and storing its data in order to meet the user's quality of service objectives such as minimizing the makespan or the cost of the workflow. In the Cloud Computing environment, the overall objective of the task scheduling strategy is to guarantee the service-level agreements (SLAs) of allocated resources to make cost-efficient decisions for workload placement (Bousselmi et al., 2016b). In real-life situations, there are three main QoS constraints that need to be considered for efficient utilization of resources: (1) minimizing the execution cost of resources, (2) minimizing the execution time of workloads, (3) reducing the energy consumption and at the same time meeting the cloud workload deadline (Singh and Chana, 2015). In our context, we focus on the objectives of reducing both the cost and the makespan of the workflow since they are the most important quality metrics for the workflow's end users (Bousselmi et al., 2016b). These metrics are conflicted because high-performance Cloud resources with high storage capacities are often more expensive than others.

The scheduling problem is more challenging for data intensive scientific workflows as they produce a huge amount of data that should be communicated to other tasks in order to perform their execution. Since data storage and data transfer through the Internet is paying, the cost and time needed for the execution of the overall workflow can be consequent. For instance, both Amazon S3 and EBS use fixed monthly charges for the storage of data, and variable usage charges for accessing the data. The fixed charges are 0.15\$ per GB-month for S3, and 0.10\$ per GB-month for EBS (Berriman et al., 2010). The variable charges are 0.01\$ per 1000 PUT operations and 0.01\$ per 10000 GET operations for S3, and 0.10\$ per million I/O operations for EBS (Berriman et al., 2010).

This work is an extension of an earlier conference publication (Bousselmi et al., 2016b). In (Bousselmi et al., 2016b), the scheduling problem of scientific workflows was formulated as a mono-objective optimization problem and we used typical workflows for the experiments. In this paper, our objective is to propose a scheduling solution for data intensive scientific workflows that considers the time and cost of data storage and transfer. The proposed extension consists on formulating our scheduling problem as a multi-objective optimization problem dedicated to Big Data scientific workflows. Indeed, we model

our scheduling problem as a bi-objective optimization problem that aims to reduce the overall cost and the response time of the workflow. As an optimization technique, we propose the extension of the cat swarm optimization algorithm (CSO) algorithm to allow the resolution of a multi-objective optimization problem. We chose the CSO algorithm since it has proven its efficiency compared to other evolutionary algorithms such as the particle swarm optimization algorithm especially in solving common NP-hard optimization problems (Chu et al., 2007) Thus, the principal contribution of our paper is the proposition of a new multi-objective optimization solution that allows the scheduling of Big Data scientific workflows based on the quality metrics of time and cost.

The remainder of this paper is organised as follows. In section 2, we summarize some works that were interested in the scheduling of scientific workflows in the cloud computing environment. In section 3, we present our problem formulation as well as the formulation of the scientific workflows and the execution environment. Section 4 describes the CSO algorithm and its components. In section 5, we detail the proposed algorithm for the scheduling of scientific workflows in the cloud. We describe our application case, namely, the LIGO workflow in section 6. Then, we report and discuss our experimental results before concluding in the final section.

2 RELATED WORKS

With the development of cloud technology and extensive deployment of cloud platform, the problem of workflow scheduling in the cloud becomes an important research topic. The challenges of the problem lie in: NP-hard nature of the task-resource mapping with diverse Quality of Service (QoS) requirements (Wu et al., 2015). Several approaches were proposed recently to resolve this problem with considering multiple QoS metrics such as the makespan, cost and energy consumption, others focused only on a single-objective problem by aggregating all the objectives in one analytical function such as in (Bousselmi et al., 2016b). (Durillo and Prodan, 2014) proposed Multi-objective Heterogeneous Earliest Finish Time (MOHEFT) which gives better trade-off solutions for makespan and financial cost when the results are compared with SPEA2* Algorithms (Yu et al., 2007). A multi-objective heuristic algorithm, Min-min based time and cost trade-off (MTCT), was proposed by Xu et al. in (Xu et al., 2016). As scheduling approach, the Balanced and file Reuse-Replication Scheduling (BaRRS) algorithm, was proposed to se-

¹<https://pegasus.isi.edu/application-showcase/ligo/>

lect the optimal solution based on makespan and cost (Casas et al., 2017).

To reduce the cost and time of data transfer, some scheduling strategies prefer employing the same cloud provider for the data storage and VMs, since the data transfer cost between inner services is free (Beriman et al., 2010). However, executing too many workloads on a single cloud will cause workloads to interfere with each other and result in degraded and unpredictable performance which, in turn, discourages the users (Dillon et al., 2010). Only few works were interested in the problem of reducing the time and cost of data transfer while executing data intensive workflows in multi-cloud environment. In (Yao et al., 2017), the authors designed a multi-swarm multi-objective optimization algorithm (MSMOOA) for workflow scheduling in multi-cloud that optimizes three objectives (makespan, cost and energy consumption) with taking into consideration the structure characteristic of cloud data center but they considered only computational-intensive workflows. Wangsom and al. proposed in (Wangsom et al., 2019) a multi-objective peer-to-peer clustering algorithm for scheduling workflows where cost, makespan and data movement are considered, but this work lacks details in the reported results.

3 EXECUTION ENVIRONMENT AND PROBLEM FORMULATION

In this section, we explain how we formulate our multi-objective combinatorial optimization problem for scheduling big data scientific workflows in the cloud. We start by defining the model used for the presentation of big data scientific workflows. Then, a model of our execution environment, namely the cloud computing, is proposed. Finally, we present our problem formulation.

3.1 Big Data Scientific Workflows Modeling

Scientific workflows are usually modeled using Directed Acyclic Graphs (DAG). Since our objective is to handle big data workflows, we model a workflow as a DAG graph where tasks are represented by nodes and dependencies between them are represented by links. Dependencies typically represent the dataflow in the workflow, where the output data produced by one task is used as input of one or multiple other tasks. A single workflow task T_i can be modeled using the

triplet $T_i = (D, Dt_{in}, Dt_{out})$ with:

- $Dt_{in} = \{Dt_{in_j} | j \in [0..(i-1)]\}$ represents the amounts of data to be processed by T_i with Dt_{in_j} is the dataset from the task j .
- $D = \{D_k | k = 1..r\}$ is the vector of demands of the T_i in terms of resource capacities with D_k is the minimum capacity required of the resource type k to process the task and r the total number of resource types.
- $Dt_{out} = \{Dt_{out} | l = 1, \dots, s\}$ represents the number of datasets generated after the completion of T_i .

3.2 Modeling the Cloud Computing Environment

The Cloud Computing environment can be modeled as a set of datacenters. Each datacenter is composed of a set of cloud clusters where each cluster is composed of a set of physical machines, network devices and storage devices. The physical machines, that we note PM , can be used for both computing and storage. PM may be a dedicated computing machine, a set of virtual machines or a shared machine. In this scope, we consider only physical machines that are composed of a set of virtual machines such as $PM = \{VM_i | i = 1, \dots, m\}$. The communication link between two physical machines $BW(PM_i, PM_j)$ is defined as the bandwidth demand or the traffic load between them. In addition, we define the quality of service metrics used to evaluate the performances of a virtual machine VM_i by the vector $Q_i = (T_{exe}, C, A, R)$ where:

- $T_{exe} = \sum_{j=1}^k Dt_{in_j} T_u$ is the execution time of a single task assessed as the sum of times needed to process each input data file Dt_{in_j} . T_u is assessed in million instructions per second (MIPS).
- The cost of processing a task on a VM_i denoted $C = T_{exe} * C_u$, where C_u represents the cost per time unit incurred by using VM_i .
- A is the availability rate of VM_i .
- R is the reliability value of VM_i .

In this paper, we will consider only the first two quality metrics, namely the execution time and the cost of VM_i .

3.3 Problem Formulation

Dispersion, heterogeneity and uncertainty of cloud resources bring challenges to resource allocation, which cannot be satisfied with traditional resource allocation

policies. Resource scheduling aims to allocate appropriate resources at the right time to the right workloads, so that applications can utilize the resources effectively which lead to maximization of scaling advantages (Singh and Chana, 2015). In this paper, our objective is to propose a new multi-objective scheduling approach for big data scientific workflows in the Cloud. The proposed approach considers two QoS metrics that can represent, according to our understanding, the most important QoS metrics for end users, namely: the makespan and the cost.

The execution of a big data workflow in the cloud implies executing each task using the allocated virtual machine with respect to the data-flow dependencies among different tasks. At the end of a task execution, generated data is transferred to one or several following tasks to perform their execution. These data are often massive and require additional time and cost to be transferred to other geographically distributed virtual machines. Some final or intermediary generated data have to be stored also in dedicated storage services to be analyzed later (Yuan et al., 2012).

We denote X_k a realizable workflow schedule where each task is assigned to a specific VM. Then, the overall makespan of the schedule X_k can be assessed by the equation 1 as the sum of the data processing and data transfer times between dependent tasks, such as:

$$Makespan(X_k) = T_{(data-processing)}(X_k) + T_{(data-transfer)}(X_k) \quad (1)$$

We argue that the time needed for data storage in dedicated storage services (SS) can't be assessed in the makespan since these data could remain in SS services even after the workflow execution end for analyze requirements. Likewise, the overall cost of a schedule solution X_k is assessed as the sum of the data processing cost on the allocated VMs, the overall cost of data transfer between tasks and data storage on SS services as follows:

$$Cost(X_k) = C_{(data-processing)}(X_k) + C_{(data-transfer)}(X_k) + C_{(data-storage)}(X_k) \quad (2)$$

Like in (Bousselmi et al., 2016a), to assess the global quality metrics values of the workflow during its execution, such as its makespan and its cost, we use aggregation functions of different considered quality metrics as shown in the table 1 with N is the total number of tasks of the workflow, T_{exe_i} , C_i , $DT_{(i,j)}$, $d_{(i,i+1)}$, T_{rem} are respectively the data processing time of a task i on a specific VM, its execution cost, the data-flow transmitted between two VMs i and j corresponding to two consecutive tasks, the throughput of network bandwidth between two VMs i and j and

the time during which data is remaining stored on a specific SS. In the case of parallel tasks of the workflow, the operators of the aggregation function for the execution time and data transfer time metrics are replaced by the maximum operator.

To sum up, our scheduling objective can be formulated as a multi-objective optimization problem that aims to reduce both the makespan and the cost of the workflow using the following linear representation:

$$\begin{cases} \min & makespan(X_k) \\ \min & cost(X_k) \end{cases} \quad (3)$$

The objective functions of our optimization problem are reported by the formulas 1 and 2. In the next session, we will introduce the CSO algorithm used for the resolution of our formulated problem.

4 CAT SWARM OPTIMIZATION (CSO) ALGORITHM

Biologically inspired computing has been given importance for its immense parallelism and simplicity in computation. In recent times, quite many biologically motivated algorithms have been invented and are being used for handling many complex problems of the real world (Das et al., 2008). Using a biologically inspired algorithm for the scheduling of scientific workflows in the cloud became an attractive research track due to complexity of the problem of finding the best match of task-resource pair based on the user's QoS requirements.

Recently, a family of biologically inspired algorithms known as Swarm Intelligence has attracted the attention of researchers working on bio-informatics related problems all over the world (Das et al., 2008). Algorithms belonging to this field are motivated by the collective behavior of a group of social insects (like bees, termites and wasps). For instance, the Cat Swarm Optimization algorithm (CSO) aims to study the behavior of cats to resolve complex combinatory optimization problems. According to (Bahrami et al., 2018), cats have high alertness and curiosity about their surroundings and moving objects in their environment despite spending most of their time in resting. This behavior helps cats in finding preys and hunting them down. Compared to the time dedicated to their resting, they spend too little time on chasing preys to conserve their energy. The first optimization algorithm based of the study of cat swarms was proposed in 2007 by (Chu et al., 2007). The steps of the CSO algorithms are as follows. First, the CSO algorithm creates an initial population of N cats and

Table 1: Aggregation functions for QoS metrics.

| Quality metric | Data processing time | Data Transfer time | Data processing cost | Data storage cost | Data transfer cost |
|------------------|------------------------------------|---|----------------------|--|--|
| Sequential Tasks | $\sum_{i=1}^N T_{exe_i}$ | $\sum_{i=1}^N \frac{DT_{i,i+1}}{d_{i,i+1}}$ | $\sum_{i=1}^N C_i$ | $\sum_{i=1}^N T_{rem_i} * C_i$ | $\sum_{i=1}^n \frac{DT_{i,i+1}}{d_{i,i+1}}$ |
| Parallel Tasks | $\max_{1 \leq i \leq n} T_{exe_i}$ | $\max_{1 \leq i \leq n} \frac{DT_{i,i+1}}{d_{i,i+1}}$ | $\sum_{i=1}^N C_i$ | $\max_{1 \leq i \leq n} T_{rem_i} * C_i$ | $\max_{1 \leq i \leq n} \frac{DT_{i,i+1} * C_{ij}}{d_{i,i+1}}$ |

defines their characteristics. Each cat represents a solution of the formulated optimization problem and its characteristics represent the parameters of this solution, namely its fitness value, its mode (in "seeking" or "tracing" mode), its position and its velocity. Second, cats are organized randomly into two groups, one group of cats is in "Seeking mode" and the other in "Tracing mode". All cats are evaluated according to the fitness function and the cat with the best fitness value is retained. After that, if the cat is in "Seeking mode", then the process of "Seeking mode" is applied, otherwise the process of "Tracing mode" is applied. Finally, if the stop condition has not been reached yet, then the algorithm restarts from the second step. The optimum solution is represented by the cat having the best fitness value at the end of the algorithm. The details of the "Seeking mode" and "Tracing mode" procedures could be found in (Chu et al., 2007). Four essential parameters are to be defined and applied in the CSO algorithm, namely:

- SMP (Seeking Memory Pool): the search memory pool that defines the size of the seeking memory of each cat.
- SRD (Seeking Range of the selected Dimension): the mutative ratio for the selected dimensions.
- CDC (Counts of Dimension to Change): the number of elements of the dimension to be changed.
- SPC (Self Position Consideration): a Boolean number indicating whether the current position is to be considered as a candidate position to be moved on or not.

5 PROPOSED BI-OBJECTIVE CSO (BiO-CSO) ALGORITHM

In this paper, we formulated our scheduling problem as a multi-objective optimization problem that aims to reduce both the overall makespan and cost of the workflow execution in the cloud. To resolve this latter, we applied the discrete version of the CSO algorithm proposed in (Bouzidi and Riffi, 2013) since our

Table 2: Example of scheduling solution representation.

| Workflow's task | Task 1 | Task 2 | Task 3 | Task 4 | Task 5 | Task 6 | Task 7 | Task 8 |
|-----------------|--------|--------|--------|--------|--------|--------|--------|--------|
| VM identifier | 4 | 1 | 2 | 3 | 4 | 3 | 2 | 5 |

objective function uses discrete integers to model a task-resource affectation.

We start by representing the scheduling solution for the CSO algorithm. Then, we will detail the steps of the proposed BiO-CSO algorithms and its tracing and seeking mode procedures.

5.1 Mapping between Problem Formulation and CSO

In our context, a scheduling solution is a set of N affectations $\langle task, VM \rangle$ for all the tasks of the workflow to the available cloud VMs that can perform these tasks. A realizable scheduling solution should respect the order of execution of the tasks so that two parallel tasks can't be assigned to the same VM. We suppose that the number of available VMs is less than or equal to the number of tasks. The table 2 shows an example of realizable scheduling solution of a workflow sample with 8 tasks. We suppose that for each task, there are 5 candidates VMs that can perform it. We suppose that the tasks 1,2, 3 and 4 can be performed in parallel and then the tasks 5,6,7,8 are in sequence.

As our scheduling problem is multi-objective, we consider that the best scheduling solution is the one having the best value of makespan and cost at the same time than all the other solutions. All non-dominated scheduling solutions are retained and considered as best solutions. To find these solutions we should first map our problem formulation with the formalism of the CSO algorithm. In fact, in the CSO algorithm, each cat represents a solution for the considered optimization problem. The overall objective of the CSO is to find the cat having the optimal solution from all the population of cats, i.e. the cat with best fitness value. Then, the mapping between our problem formulation and the CSO algorithm can be summarized in the following table 3.

Table 3: Mapping between our problem formulation and CSO.

| Optimization problem formulation | CSO algorithm |
|--|----------------------------|
| A scheduling solution | A cat |
| The number of workflow tasks | A cat's position dimension |
| $\langle tasks, VM \rangle$ affectations | A cat's position vector |
| The objective function | The fitness value of a cat |

5.2 BiO-CSO Algorithm Details

The CSO algorithm was designed for mono-objective optimization problems (Chu et al., 2007; Bouzidi and Riffi, 2013). We extended this algorithm to consider multi-objective optimization problems. First, we define a vector of best solutions to store a fixed number of non-dominated solutions along the evolution process. Then, for each cat, we replace its fitness value by a structure that stores the two values of our objective functions (the makespan and the cost of the scheduling solution described by this cat). Also, we implement additional functions to compare the new generated solutions in order to update the vector of best solutions. The pseudo code of the proposed algorithm is given by Algorithm 1 detailed hereafter.

Algorithm 1: BiO-CSO.

Input: MaxLoopIteration :maximum number of iterations

Variables:
 CATS: population of cats (scheduling solutions)
 BestParetoSet: vector of stored non-dominated solutions

- 1: Initialize the cats' population CATS
- 2: Initiate the position and velocity for each cat
- 3: Pick randomly a mode for each cat (CatModes)
- 4: **for all** iteration $I < MaxLoopIteration$ **do**
- 5: **for all** cat in CATS **do**
- 6: fitnessValue = EvaluateFitness(cat) // *evaluate the makespan and cost*
- 7: Apply BiO-Replacement(fitnessValue, BestParetoSet) // *Apply cat mode procedure*
- 8: **if** catModes (cat) == SeekingMode **then**
- 9: Apply seekingModeProcedure(cat,CATS)
- 10: **else**
- 11: Apply TracingModeProcedure(cat,CATS)
- 12: **end if**
- 13: Re-pick randomly a mode for each cat CatModes
- 14: **end for**
- 15: **end for**

BiO-CSO works as follows. First, the algorithm creates N cats and initialize their properties (step 1 and 2), each cat has an M dimensional position vector that represent a realizable scheduling solution with M is the total number of tasks per workflow. The values of the position vector represent the identifiers of VMs affected respectively to each task of the workflow. The fitness value of a cat is the Bi-value vector that represents the makespan and the cost of the represented scheduling solution. Second, BiO-CSO randomly chooses a sequence of cats using the MR parameter value, set them in the 'Tracing mode' and set the rest of cats in 'Seeking mode' (step 3). Third, it evaluates the fitness values (step 6) of each cat (the makespan and cost of the scheduling solution represented by a cat) and keep the cat with the best fitness values. The retained solutions may be several as our problem is multi-objective. Then, we keep all the non-dominated solution in a vector representing the best solutions (step 7). The update of this vector is described below (algo 2). If the cat is in 'Seeking mode', then the 'Seeking mode' procedure is applied as described in the algorithm 3, otherwise, apply the procedure of 'Tracing mode' as described in the algorithm 4. Forth, the algorithm re-selects a sequence of cats, sets them to tracing mode and puts the rest of the cats in 'Seeking mode'. Finally, if the termination condition is reached (the maximum number of iterations is reached or a desirable best solution is found), then terminate the algorithm. If not, return to step 4.

Bi-Objective Replacement: As for most of the multi-objective optimization algorithms (Reyes-Sierra et al., 2006), the proposed BiO-CSO is Pareto-based. Its purpose is to generate a Pareto solution set with optimized makespan and cost. The algorithm maintains a list of best Pareto solutions that consists of a subset of non-dominated designs found so far. Since the number of best solutions can quickly grow very large, the size of the Pareto list is limited. Each cat is inserted into the Pareto set only if it is non dominated as described in algorithm 2.

Tracing Mode and Seeking Mode Procedures: The Seeking mode corresponds to the resting state of the cats while being in continuous research of a better position. In this mode, the BiO-CSO algorithm performs multiple copies of the cat's position (the scheduling solution), then, modifies this solution using the parameter CDC for all the position copies (CDC defines the number of values of the position to be changed). After that, it calculates the distance between the new position's fitness value and the old one. According to the distance value, the algorithm chooses the new position to assign to the cat according to its probability value. In tracing mode, cat tries to trace targets by

updating their velocity and positions. The details of these procedures are given in the algorithms 3 and 4.

Algorithm 2: BiO-Replacement.

Inputs:

BestParetoSet: vector of stored non-dominated solutions

X: a cat (a solution)

- 1: **if** X is dominated by an other cat Y in BestParetoSet **then**
 - 2: Replace X parameters by Y parameters
 - 3: **else**
 - 4: **if** X is non-dominated solution and BestParetoSet is not full **then**
 - 5: Add X to BestParetoSet
 - 6: **else**
 - 7: Find a solution from BestParetoSet to be replaced by the new one X
 - 8: **end if**
 - 9: **end if**
-

Algorithm 3: Seeking Mode Procedure.

Inputs:

cat: a solution

SMP: number of cat copies

- 1: Perform SMP cat position copies in PositionCopies
 - 2: **for all** position in PositionCopies **do**
 - 3: modify cat position
 - 4: **end for**
 - 5: **for all** cat corresponding to a position in PositionCopies **do**
 - 6: Compute cat fitness
 - 7: Calculate the distance of the new solution to previous one
 - 8: Move cat to the new picked position with a given probability
 - 9: **end for**
-

Algorithm 4: Tracing Mode Procedure.

Inputs: CATS: a set of solutions

solutions

- 1: **for all** cat in CATS **do**
 - 2: Update cat velocity
 - 3: Update cat position
 - 4: **end for**
-

6 BIG DATA SCIENTIFIC WORKFLOWS: CASE OF LIGO

The LIGO concept built upon early work by many scientists to test a component of Albert Einstein's theory of relativity, the existence of gravitational waves [10]. LIGO operates two gravitational wave observatories in the USA: the LIGO Livingston Observatory in Livingston, Louisiana, and the LIGO Hanford Observatory, Washington. These sites are separated by 3,002 kilometers (1,865 miles) straight line distance through the earth, but 3,030 kilometers (1,883 miles) over the surface [10]. Other similar observatories were built later in other countries in the world. The workflow of LIGO Inspiral Analysis is one of the LIGO project defined workflows using scientific workflows to detect and analyze gravitational waves produced by various events in the universe. The LIGO Inspiral Analysis workflow [22] is a CPU and Data intensive workflow used to analyze the data obtained from the coalescing of compact binary systems such as binary neutron stars and black holes. The time-frequency data from each of the three LIGO detectors is split into smaller blocks for analysis. For each block, the workflow generates a subset of waveforms belonging to the parameter space and computes the matched filter output. If a true inspiral has been detected, a trigger is generated that can be checked with triggers for the other detectors. A simplified representation of the workflow is shown in figure 1 with sub-workflows outlined below. This workflow may contain up to 100 sub-workflows and 200,000 jobs (Juve et al., 2013). The execution of a sample instance of the LIGO Inspiral workflow with 3981 tasks on a local machine runs for 1 day and 15 hours and uses more than 1 Go of input data and more than 2 Go of output data [9].

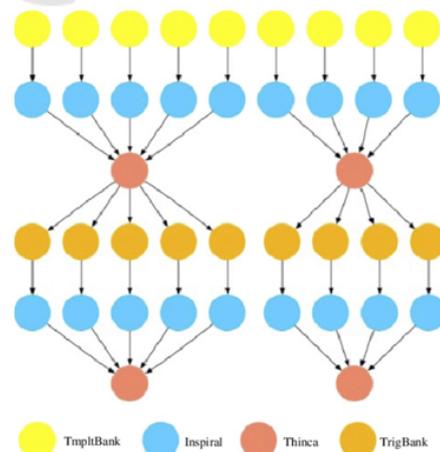


Figure 1: Structure of the LIGO Inspiral Workflow.

In order to establish a confident detection or measurement, a large amount of auxiliary data will be acquired (including data from seismometers, microphones, etc.) and analyzed (for example, to eliminate noise) along with the strain signal that measures the passage of gravitational waves (Deelman et al., 2002). The raw data collected during experiments is a collection of continuous time series at various sample rates. The amount of data that will be acquired and cataloged each year is on the order of tens to hundreds of terabytes (Deelman et al., 2002). These data are treated in real time by the LIGO workflow, so, several instances of this workflow should run simultaneously to achieve the objectives of the project. Consequently, the execution environment should provide efficient tools and enough resources for the execution and the data storage of the LIGO workflow. Therefore we have chosen this workflow as an application case for our proposed algorithm BiO-CSO for its excessive use of computing and storage resources.

7 RESULTS AND DISCUSSION

In this section, we present the execution environment for our proposed BiO-CSO algorithm and the used parameters for our algorithm and the PSO algorithm. Then, we will expose our experimentation results and discuss the performances of our proposition.

7.1 Execution Environment and Algorithm Parameters

As a simulation environment, we used the WorkflowSim (Chen and Deelman, 2012) on which we implemented our proposed algorithm BiO-CSO. WorkflowSim extends the CloudSim (Calheiros et al., 2011) simulation toolkit by introducing the support of workflow preparation and execution with an implementation of a stack of workflow parser, workflow engine and job scheduler. It supports a multi-layered model of failures and delays occurring in the various levels of the workflow management systems. CloudSim is an extensible simulation toolkit that enables modeling and simulation of Cloud computing systems and application provisioning environments (Calheiros et al., 2011). We tested the LIGO Inspiral workflow with different scales using respectively 30, 100 and 1000 tasks. To compare our algorithm performances, we implemented also the Multi-objective version of the Particle swarm optimization (PSO) algorithm. PSO is an optimization algorithm

which is based on swarm intelligence (Tripathi et al., 2007). It is inspired by the flocking behavior of birds, which is very simple to simulate and has been found to be quite efficient in handling the single objective optimization problems. The simplicity and efficiency of PSO motivated researchers to apply it to the MOO problems since 2002 (Calheiros et al., 2011). For the MOPSO algorithm, we consider $c1=c2=2$, and the inertia weight 'w' varying from 0.4 at the beginning to 0.9 at the end of execution of the algorithm. The table 4 illustrates the general parameters of the CSO algorithm. The initial population size of cats is set to 32.

Table 4: Generic parameters of CSO algorithm.

| Parameter | Value |
|-----------|-------|
| SMP | 5 |
| SRD | 20 |
| CDC | 80% |
| MR | 10% |
| c1 | 2.05 |
| r1 | [0,1] |

7.2 Experimentation Results

The experiments were repeated 20 times by varying the value of MP parameter of CSO for each workflow instance. Average values of the output metrics are reported in this section. This manipulation allowed to generate more values of best solutions for our considered objective functions. Thus, all non-dominated solutions obtained over each run are recorded in a general Pareto set. Figure 2 shows the obtained Pareto set for LIGO workflow with 30, 100 and 1000 tasks respectively for the proposed BiO-CSO algorithm and the MOPSO one. In the figure 2, it is clear that the makespan and deployment cost functions are strongly correlated with a negative slope. With respect to the Pareto front, BiO-CSO demonstrates the ability to generate better diversity of solutions with a uniform distribution than MOPSO. In fact, the CSO algorithm was designed to avoid the problem of premature convergence thanks to the weighted position update strategy.

It can also be seen that the Pareto fronts obtained using BiO-CSO are superior for all the different scales of LIGO workflow. In figure 2, the non-dominated solutions generated by the BiO-CSO algorithm are better spread over the solutions space than those of the MOPSO algorithm which offers more choices to the workflow end user to make cost-efficient decisions.

In figure 3, we can see that the BiO-CSO algorithm achieves better average values for both Makespan and Cost for all the LIGO workflow instances. For the Makespan, the BiO-CSO allows

¹<https://pegasus.isi.edu/portfolio/ligo/>

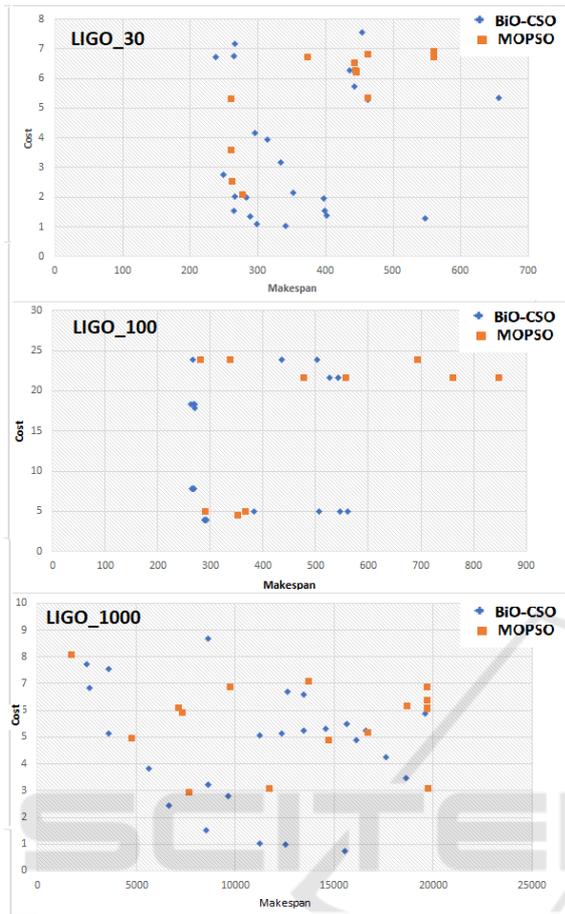


Figure 2: Non-Dominated solution sets generated by BiO-CSO and MOPSO for LIGO workflow.

to gain an improvement of 16,03%, 47,21%, and 11,23% compared to the MOPSO algorithm for the different scaled of LIGO workflow. The cost of BiO-CSO was better by 18,41%, 27,75%, and 34,82% over the MOPSO algorithm for the LIGO workflow with 30, 100 and 1000 tasks respectively.

Table 5: Average values for computation and data transfer time.

| Workflow | Data transfer time | | Computation time | |
|-----------|--------------------|----------|------------------|---------|
| | BiO-CSO | MOPSO | BiO-CSO | MOPSO |
| LIGO_30 | 261,91 | 295,05 | 96,87 | 109,13 |
| LIGO_100 | 274,39 | 382,42 | 81,96 | 114,23 |
| LIGO_1000 | 9255,56 | 11022,28 | 1895,72 | 2257,58 |

The table 5 illustrates the average values for computation and data transfer time obtained by BiO-CSO and MOPSO for the LIGO workflow with 30, 100 and 1000 tasks. It shows that our algorithm BiO-CSO achieves better average values than MOPSO for all the LIGO workflow scales with a gain ranging from 11% to 28%. For example, for the LIGO_1000, the BiO-CSO data transfer time is 16% better the the



Figure 3: Cost and makespan average values obtained by BiO-CSO and MOPSO for LIGO workflow.

MOPSO one. It can be seen also that the data transfer time represents roughly 75% of the workflow execution time due to the huge amount of data treated.

8 CONCLUSION

In this paper, we addressed the problem of Big Data Scientific workflow scheduling in the Cloud. We started by formulating the problem as a bi-objective optimization problem minimizing both the makespan and the cost of the workflow. We then proposed the BiO-CSO algorithm which is an extension of the CSO one to deal with bi-objective optimization. The objectives conflict is resolved using the Pareto analysis. The experiments were conducted on the LIGO workflow and discussed in comparison to the current algorithm frequently used for workflow scheduling e.g. PSO. From the experimental results we conclude that BiO-CSO is able to produce better schedules based on multiple objectives. Future work will investigate the extension of BiO-CSO to deal with further objectives such as minimizing energy consumption and its application on different types of Big Data scientific workflows.

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