Reconfiguration Penalty Calculation for Cross-cloud Application Adaptations

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Abstract: Cloud's indisputable value for SMEs and enterprises has led to its wide adoption for benefiting from its costeffective and on-demand service provisioning. Furthermore, novel systems emerge for aiding the cross-cloud application deployments that can further reduce costs and increase agility in the everyday business operations. In such dynamic environments, adequate reconfiguration support is always needed to cope with the fluctuating and diverse workloads. This paper focuses on one of the critical aspects of optimal decision making when adapting the cross-cloud applications, by considering time-related penalties. It also contributes a set of recent measurements that highlight virtualized resources startup times across different public and private clouds.

1 INTRODUCTION

In cloud computing besides the on-demand provisioning of resources, users are enabled with features that allow the seamless adaptation of the allocated resources, used for hosting applications according to the constantly fluctuating workload needs. This is achieved by either scaling in or scaling out the infrastructure in times of lower or higher demand, respectively. This ability to dynamically acquire or release computing resources according to user demand is defined in the computer science as elasticity (Verma et al., 2011). Providing infrastructural resources i.e. Virtual Machines (VMs), becomes very important when these resources can be ready in time to be used according to the users' expectations.

Nowadays, modern data-intensive applications increasingly rely on more than one cloud vendors, a fact that makes elasticity even more challenging as a feature (Horn et al., 2019). In order to decide in each situation, based on a given application topology and fluctuating workload, several aspects of the reconfiguration costs should be considered (such as time cost, data lifecycle cost etc.). In this work, we present how the time dimension of this cost can be considered, based on the VM startup times and the application component deployment times. This cost is evaluated as a part of a utility function that can reveal whether a certain reconfiguration action is optimal for the current cross-cloud application. Specifically, an algorithm and a software tool are presented, in section 3, for calculating the reconfiguration cost of each alternative topology that should be examined towards a cloud application reconfiguration. In section 4, we highlight the importance of the penalty calculator for the reconfiguration decision making by using an illustrative example. Through a set of related startup measurements of virtualised resources in prominent vendors, we reveal important findings about VM provisioning. Last, we conclude this work and discuss next steps in section 5.

2 RELATED WORK

In this section, we discuss some of the studies performed that focus on the reconfiguration costs in the cloud and the multi-cloud environment. Such a work (Mao and Humphrey, 2012) provides a systematic study on the cloud VM startup times across three cloud providers (i.e. Amazon EC2, Windows Azure and Rackspace). In this study, measurements were reported, while an analysis of the relationship among the VM startup time and different factors, is given for comparing the three cloud

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providers. These factors include the size of OS instance image, the instance type of the VM, the number of instances concurrently deployed and the time within the day that the reconfiguration/startup is performed. Although this study is valuable the measurements have been performed back in 2012 and they need to be updated, while there is a lack of exploitation of such data in terms of reconfiguration decision making. In another work (Salfner et al., 2011) the authors analyse the VM live migration downtime during the reconfiguration process using different cloud resources. The results from the analysis after various experiments, showed that the total migration time as well as the downtime of the services running on the migrated VMs are mainly affected by the memory usage of the VMs used. But besides the significant findings, there is no described method on how to take into consideration this time cost in the reconfiguration process of the cloud infrastructure in order to minimize its impact. What is more the multi-cloud case of reconfiguration is not examined at all. In a different approach, the authors (Yusoh and Tang, 2012) propose a penalty-based Grouping Genetic Algorithm for deploying various Software as a Service (SaaS) composite components clustered on VMs in different clouds. Their main objective was to minimize the resources used by the application and at the same time maintain an adequate quality of service (QoS), respecting any constraints defined. Based on the experimental results, their proposed algorithm always produces a feasible and cost-effective solution with a quite long computation time though. In addition, no action is taken in this study to incorporate in this penalty calculation the time dimension for provisioning VMs, as a crucial aspect of the reconfiguration process and the availability of cloud applications.

Considering time aspects for the reconfiguration penalty in the multi-cloud environments, it is also noteworthy to examine cases were resources should be used for which no prior data is available (e.g. a custom VM for which no previous measurements are available). In such cases several approaches exist that are valuable. Uyanik and Guler (Uyanik and Guler, 2013) analyse in their study whether or not the five independent variables in the standard model were significantly predictive of the KPSS score (Kokoszka and Young, 2015), the dependent variable, based on ANOVA statistics (Rutherford, 2001). Their primary objective was to exemplify the multiple linear regression analysis with its stages. The assumptions, necessary for this analysis, were examined and the

¹ http://camel-dsl.org/

regression analysis was performed using related data that were satisfying the assumptions. The standard model's prediction degree of the dependent model was R=0.932, while the variance of the dependent variable was R2=0.87. The model seems to predict appropriately the dependent variable, but it is not so accurate as the ordinary least squares (OLS) Multiple Linear Regression algorithm (Rutherford, 2001). Specifically, in the case of OLS algorithm a greater than 95% value of *R2* is achieved which means that the proportion of the variance in the dependent variable that is predicted from the independent variables is greater than 95%. OLS regression algorithm is one of the major techniques used to analyse data and specifically to model a single response variable which has been recorded on at least an interval. For the above reasons the specified OLS method is used for the Penalty Calculator Algorithm described in the section 3.2.

3 PENALTY CALCULATOR

In this paper, an innovative platform which is called Melodic is used as an automatic DevOps for managing the life cycle of cross-cloud applications (Horn et al., 2019), (Horn and Skrzypek, 2018). The Melodic platform is built around a micro-services architecture, able to manage container-based applications and support some of the most prominet big data frameworks. The main idea of Melodic is based on models@run.time and states that the application architecture, its components and the data to be processed can all be described using a Domain Specific Language (DSL). The application description includes the goals of the efficient deployment (e.g. reduce cost), complies with the given deployment constraints (e.g. use data centres located in various locations), and registers the current state of the application topology, through monitoring, in order to optimize the deployment of each application component.

The Melodic platform-as-a-service (PaaS) is conceptually divided into three main parts: i) the Melodic interfaces to the end users; ii) the Upperware; and iii) the Executionware. The first part comprises tools and interfaces used to model users' applications and datasets along with interacting with the PaaS platform. Moreover, the PaaS is using modelling interfaces that are established through the CAMEL¹ modelling language, which provides a rich set of DSLs with modelling artefacts, spanning both the design and the runtime of a cloud application as well as data modelling traits. The second part (Upperware) is responsible to calculate the optimal application component deployments and the appropriate data placements on dynamically acquired cross-cloud resources. The optimal configuration of the cross-cloud application topology refers to a utility function evaluation. The utility function can be defined as the function, introduced as a measure of fulfillment for applying reconfiguration for crosscloud applications. This utility function requires the use of a Penalty Calculator (used as a library) which focuses on the reconfiguration time cost. In this paper the time reconfiguration cost is mainly examined, while there can also be other parameters to consider such as the cost of transferring data. An important part of this evaluation includes the VMs startup times along with the expected deployment times of specific application components that are to be reconfigured. The Penalty Calculator provides normalized output values between 0 and 1, where 0 indicates the lowest possible penalty which indicates the most desired solution and *1* indicates the highest possible penalty which is the less desired solution.

For example, a utility function can be defined as follows:

$$UtilityFunction = \frac{1}{C_{Solution} + C_{\text{Reconfiguration}}}$$
(1)

Where the $C_{Solution}$ is a function of the number of resources used for deployment. This implies the satisfaction of certain goals (e.g. minimize the deployment cost, minimize response time etc.) expressed as a mathematical function:

$$C_{Solution} = f(no_of_resources)$$
(2)

While the $C_{Reconfiguration}$ is a function of the result of the Penalty Calculator. This result represents the value given by the Min-Max normalization method, applied over the Ordinary Least-Squares Regression (OLS) algorithm result described in paragraph 3.2:

$$C_{Reconfiguration} = f(penalty_value)$$
 (3)

The third part of Melodic includes the Executionware which executes the actual cloud application deployments and reconfigurations by directly invoking the cloud providers APIs.

3.1 Approach

Penalty Calculator's objective is to calculate a normalized reconfiguration penalty value by

comparing the current and the new candidate configuration, coming from a constraing programming solver component of the Upperware. Therefore, the system examines a sequence of candidate configurations under specific constraints and optimization goals (e.g. reduce cost and increase service response time) that will serve according to the desired QoS the incoming workload. The Penalty Calculator affects the decision on accepting and deploying a new candidate cross-cloud application topology based on its' function value. The smaller the penalty function is, the better is for the candidate solution as it implies a smaller time for materializing the proposed reconfiguration.

The Penalty Calculator is a part of the Melodic Upperware and it is used as a library by the Utility Generator, a component that calculates a single value for each candidate solution, according to a utility function that expresses the overall goals of the application. The Penalty Calculator receives from the Utility Generator, XMI files describing the collections of configuration elements for the current and the new proposed configuration (OS, hardware and location related information of the virtualised resources to host certain application components).In order to use a high-performance, distributed memory caching system intended to speed up the penalty function calculations, the VM startup time data are stored in memory, using the Memcache² solution. The categorization of various VM startup times include multiple variables for resources such as the RAM, CPU cores, Disk, VM types names etc. Regarding the component deployment times, these are persisted and retrieved from a time-series database. The various application components that are deployed in cross-clouds are constantly measured with respect to the deployment time needed and based on the virtualized resource used. By using a timeseries database for this purpose, it allows for a quick retrieval of the average deployment times. In this work InfluxDB³ was used.

3.2 The Penalty Calculator Algorithm

Based on the feed from the Utility Generator in Melodic, the Penalty Calculator algorithm is applied for comparing the old and the new proposed (candidate) solution, issuing a penalty value, thus affecting the decision on whether or not a specific new solution should be deployed. This algorithm uses measured VM startup times and measured component deployment times (their average values) for

² https://memcached.org/

³ https://www.influxdata.com/

calculating the time-related cost for changing from the current to a new application topology. If there are no component deployment times from past measurements, the algorithm takes into consideration only the VM startup times. In case of new custom VMs are to be provisioned, the Ordinary Least Squares Regression Algorithm (Hutcheson, 2011) is used to estimate the expected startup time by exploiting the measurements of the available predefined cloud providers' flavours. The general flow of the Penalty Calculator is given in Figure 2.



Figure 2: Penalty Calculator's Flowchart.

It is very important to note that since the VM startup times is not a constant property of the VM and of each cloud provider, but depends on the current state, load and configuration of each given cloud infrastructure in conjunction with the chosen VM, the used startup times values in the algorithm are real ones and are updated and fetched in real-time from a time-series database where these are stored.

Regarding the OLS algorithm, a single response variable is used to model the VM startup time which has been recorded for a specified range of values. The specific technique is applied to multiple variables that have been appropriately coded (i.e. RAM usage, CPU core number and Disk usage). It's purpose is to calculate startup times regarding (custom) VM flavours for which we do not have mesurements from previous deployments. The general format of the OLS model includes the relationship between a continuous response variable *Y* and some continuous variables *X* by using a line of best-fit, where *Y* is predicted at least to some extend by variables *X*:

$$Y = a + b_1 * X_1 + b_2 * X_2 + b_3 * X_3$$
(4)

In equation (4), α indicates the value of Y when all values of the explanatory variables X are equal to zero. Each parameter b indicates the average change in Y that is associated with a unit change in X, whilst controlling the other explanatory variables in the model. The Min-Max normalization method is used as a last step in the Penalty Calculator by considering the average values of all the VMs (to be used) startup times of new configuration plus the average value of the component deployment times.

4 AN ILLUSTRATIVE EXAMPLE

We note that the Penalty Calculator, presented in this paper has been tested and evaluated in several realapplication scenarios. In this section, we present one of them as an illustrative example for highlighting the value of such an approach. We refer to a traffic simulation application which is used by the company CE-Traffic for the analysis of traffic and mobilityrelated data as a basis for optimization and planning in major European cities. The initial deployment consists of five main components instances (also seen in Figure 3): i) traffic evaluation component (single instance); ii) simulation manager (single instance); and iii) simulation workers (three instances). The traffic evaluator component is responsible for the traffic analysis and sends to the simulation manager information about the need of executing a simulation. On the other hand, simulation workers are components responsible for evaluating traffic simulation settings received from the simulation manager.



Figure 3: Model of CET Traffic Simulation App.

We consider the following constraints and requirements described in the data farming application CAMEL model:

- Single instance of the traffic evaluation component
- Single instance of the simulation manager
- Between 0 and 10 instances of workers
- At least 2 CPU Cores per worker
- At least 2GB of RAM per worker

As expressed in the CAMEL model of the application, reconfiguration and later horizontal scaling of simulation worker instances is supposed to happen within a limit of 1 to 100 instances. To trigger this reconfiguration, the simulation manager collects several metrics:

- TotalCores the total number of cores available in workers
- RawExecutionTime the time of performing a single task (running a single simulation) by a worker
- SimulationLeftNumber the number of tasks (simulations) which should be still performed
- RemainingSimulationTimeMetric the remaining time in which the data farming experiment should be finished

Values of these metrics are computed and updated by the simulation manager which sends them to PaaS platform described in the introduction of section 3. In this PaaS platform we have implemented a distributed complex event processing system that is able to process incoming monitoring data in hierarchical way (Stefanidis et al., 2018). Based on this processing the system is able to detect at the appropriate time when a new reconfiguration should be initiated to cope with the detected current workload of the application. The specified system receives values of metrics and checks whether the data farming experiment is expected τo be finished on time. Finally the 'MinimumCores' composite metric is calculated in order to help for the reconfiguration. In our case example, 2 new workers are added in order to finish the traffic simulation that described before. When the simulations are finished the 'MinimumCores' composite metric is equal to zero and in the next reconfiguration workers are being removed.

Although the above system works fine in the majority of the cases, there are edge cases where a reconfiguration might start (based on the scalability rules) although the remaining simulation time is quite small. In fact this means that we might observe a behaviour where our system starts a reconfiguration cycle which until it is fully implemented, the application simulation will have been completed. Therefore the consideration of the time that is needed for any reconfiguration and as a consequence the time penalty that our component calculates, is a critical factor to be considered.

Such cases are resolved succesfully by using a Penalty Calculator component that receives two configurations schemas that are provided to it as input. The new configuration schema presents new elements (i.e. a new predefined VM flavour) and some custom VM flavours, not predefined in the used cloud providers (i.e. t1.microcustom). Specifically, the predefined VM types in this example come from 2 cloud vendors: Amazon EC2 and Openstack. By using the normalized value that it is produced from Penalty Calculator and considered in the Utility Function (UF) the previous described unecessary reconfigurations are avoided. Zero is the most desired output of Penalty Function and if the output is closer to that value, it implies a smaller time for materializing the proposed reconfiguration. On the other hand, if the output of Penalty Calculator is closer to one, then this is not desired and affects negatively the UF for a new reconfiguration. In this way reconfigurations that impose delays unacceptable according to the current application context are avoided.

4.1 Experiment Measurement Results and Analysis

Openstack Flavours	VCPUs	RAM (in MB)
m1.small2	2	1024
m1.medium2	4	4096
m1.large2	8	8192
m1.xlarge	8	16384

Table 1: Openstack Flavours Used.

Table 2: Amazon EC2 Flavours Used.

EC2 Flavours	VCPUs	RAM (in MB)
t2.micro	1	1024
t2.small	1	2048
t2.medium	2	4096
t2.large	2	8192
t2.xlarge	4	16384
t2.2xlarge	8	32768

Considering the importance of the VM startup times in cloud application reconfigurations, we conducted a performance study that is presented in this section. Similar to this work (Mao and Humphrey, 2012), we conducted new measurements across one private and one public cloud provider, specifically: i) an internal testbed offered by the university of ULM in Germany that corresponds to an Openstack installation; and ii) Amazon AWS. A number of different regions from the public providers and several VM types were used in this analysis, which focused on the VM startup times. More than 2500 measurements were conducted that involved the provisioning of different VM flavours, hosted in different data centre locations and with an increasing number of VMs instantiated simultaneously (Table 1, Table 2). For all of these VMs, the same Ubuntu images were used.

To describe the lifecycle of the cloud VM instances, cloud providers use a set of status tags to indicate the states of the provisioned VM instances. To make the definition of startup time consistent across the cloud providers that were used in our measurements, we ignore the status tags and considered as VM startup time the duration from the time of issuing a VM provision request to the time that the acquired instances can be logged in remotely.



Figure 4: Average Startup times by Openstack VM flavour (including standard deviations).

The first set of measurements across the cloud providers focused on the relationship between the The second set of measurements was focused VM startup time and the VM flavour used. Each set of measurements included for each specific VM flavour the provisioning of 1-20 instances either sequentially or in parallel (by incrementally increasing the VMs requested simultaneously). The threshold of 20 instances per set of measurements was imposed by the API limitations of the providers.



Figure 5: Average Startup times by Amazon EC2 VM flavour (including standard deviations).

According to the outcome of these measurements which can be found in Figures 4-5 the VM startup time is longer for the private cloud provider than the public one. Specifically, the Openstack VMs are provisioned with an average startup time from 81 to 123 seconds depending on the VM flavour, while the rest startup times are found from 49 to 66 seconds for EC2 Cloud VMs. This is quite expected if we consider the wide range of resources that is employed by big vendors. With respect to the variance of the conducted measurements, we found that the standard deviation in Openstack VMs' startup time is also significantly higher than those of the public provider. This implies a much more unstable environment in the case of the private provider both in terms of infrastructural resources and scheduling mechanism.



Figure 6: Average Startup times by Amazon EC2 Availability Zones (including standard deviations).

around the different data centre locations offered by the AWS public cloud providers and how this may impact the startup time of VMs. In Figure 6, we present the findings of our measurements. It is important to note that we do not find a significant fluctuation of the VM startup times as the requests for VMs provisioning change among regions and availability zones. The 55 seconds was the average startup time even for VMs provisioned in US locations. A slight improvement by 5 seconds was observed in all VMs provisioned from the data centre located in Paris, while the standard deviation of these measurements didn't exceed the 15 seconds.

In the last set of measurements, we tried to examine the impact in the VM startup times as we increased the number of VMs that were requested simultaneously, reaching up to 20 VMs in parallel (which is the threshold set by the Cloud providers). The results are presented in Figures 7-8. In Openstack VMs, we detected, as expected a much higher fluctuation in the VM startup times, which is gradually reduced as the requested VMs increase. In addition, we found significant fluctuations among the same number of instances startup that reached even the amount of 89 seconds when 7 VMs requested in parallel, a fact that reveals unstable behaviour in case of the private cloud provider. In the case of the public provider, we witnessed a much more balanced behaviour with minor fluctuations in the startup time. Specifically, we observed average startup times between 45 (for 10 instances) and 53 seconds (for 11 instances) as different simultaneous VMs startup requests were submitted. This is quite reasonable as the scheduling is done online and there are always enough spare resources to directly schedule the considerably small amount of resources that we were requesting. We also note that in the previous similar work (Mao and Humphrey, 2012), the authors have measured in 2012 an average startup time in Amazon EC2 VMs that of 100 seconds while in our recent measurements we witnessed 48% shorter times. This fact affirms the significant investments in infrastructure that public cloud providers have made over the last years.



Figure 7: Average Startup times in Openstack by the number of concurrent instances (including standard deviations).



Figure 8: Average Startup times in Amazon EC2 by the number of concurrent instances (including standard deviations).

4.2 Penalty Calculator Results

According to the measured values from the previous paragraph 4.1, we present the VM startup times stored in the system (Memcached memory) in order to be used by the proposed Penalty Calculator for the needs of our example: t2.micro-56 sec, t2.small-58 sec, t2.medium-66 sec, t2.large-52 sec, t2.xlarge-50 sec, t2.2xlarge-49 sec, m1.tiny-55 sec, m1.small-80 sec, m1.medium-120 sec, m1.large-90 sec, m1.xlarge-93 sec.

A table is used with the specific values on RAM, CPU cores, and Disk for each type described in Table 3. This table is also stored in Memcached for fast retrieval.

VM startup time (sec)	Number of cores for vCPU	RAM (GB)	Disk (GB)
56	1	0.6	0.5
58	1	1.7	160
66	4	7.5	850
52	8	15	1690
50	7	17.1	420
49	5	2	350
55	1	0.5	0.5
80	_1	2.048	10
120	2	4.096	10
90	4	8.192	20
93	- 8 - 1	16.384	40

Table 3: VM Startup Times mapped to resources.

The values of Table 3 are used to train the Ordinary Least-Squares Regression algorithm which is used to help in the prediction of the unknown VM startup times. By that way, the weights of the OLS algorithm are adapted. The new custom VM type that is used in this case is the *t1.microcustom* with a predicted startup time of 57 sec.

Moreover, the component deployment times have to be considered in the penalty calculator as explained in section 3. The measured component deployment times that have been stored in the InfluxDB are: *Traffic evaluation compontent* - 372.7659902248333 sec, *Simulation manager component* -383.61119407688045 sec and *Simulation Worker* (per each of the 3 instances) - 323.87364700952725 sec. By using the above VM startup times and the component deployment times the following regression parameters of the equation (4) are produced:

A=96.69038582442504	_
B1=-8.070707346640273	
B2=1.7404837523622727	
B3=7.407279675477281E-4	
With a r-Squared parameter: 0.9894791420723722	

Based on these results, this algorithm is quite accurate and depends on the value of the 3 explanatory variables to 98.95% and 1.7% to the constant value of *a*. This is used in order to give an accurate prediction for any custom VM type that may be used as part of a new configuration in the new Cloud infrastructure. Last, by using the Min-Max normalization method, the system calculates a Penalty value which is the *normalized average value* of the VM startup time and component deployment time and equals to 0.52197146827194. Based on this value, the Utility Generator component is able to decide the most appropriate configurations.

5 CONCLUSIONS

In this paper we focused on one of the critical aspects for optimal decision making, with respect to reconfiguration, in the dynamic environment of crosscloud applications. Specifically, we presented a system for calculating time-related penalties when comparing candidate new solutions that adapt a current application topology which is unable to serve an incoming workload spike. The algorithm implemented considers both VM startup times, across different providers and application component deployment times for calculating a normalized penalty value. This paper also discussed a set of recent measurements that highlight virtualization resources startup times across different public and private providers.

The next steps of this work include the extension of the VMs startup time measurements across more providers, regions using additional VM flavours. Moreover, this work will continue with the consideration of data management and migration related times for considering the complete lifecycle management when calculating reconfiguration (timerelated) penalties.

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