

Towards Multi-objective Optimisation of Quantitative Goal Models using Constraint Programming

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Keywords: Multi-objective Optimisation, Goal-oriented Requirements Engineering, Search-based Software Engineering, Quantitative Reasoning, Pareto Front, Tool Support.

Abstract: Goal Model are widely used to capture system goals and refine them into operational requirements assigned to human, hardware or software. Such models support alternative goal refinements resulting in a potentially large design space to explore. A given design can be quantitatively evaluated in terms of its fulfilment of a set of non-functional requirements (e.g. reliability, performance) or business goals (e.g. costs, stakeholder satisfaction). Optimisation techniques can be used to explore the design space to determine an optimal design according to a single objective like the cost but also according to multi-objective techniques to propose a set of Pareto-optimal solutions in which a best solution can be selected. In this paper, we show how to translate a goal-oriented requirements model, expressed in the KAOS notation, into a constraint programming (CP) problem. The OsaR.CP engine is used to get, from all alternatives explored, either global or Pareto-optimal solutions. Our method is implemented as a tool plugin of a requirements engineering platform and is benchmarked on the classical meeting scheduler case study.

1 INTRODUCTION

Requirements Engineering (RE) is concerned with the elicitation, evaluation, specification, consolidation, and evolution of the objectives, functionalities, qualities, and constraints of a software-based system (van Lamsweerde, 2009). It is a crucial phase as many project failures can be traced back to flaws in requirements. In addition to dealing with completeness and consistency issues, RE also allows one to take various high-level design decisions related to several ways to reach the system goals using different combinations of software/hardware/human agents, generally also resulting in different trade-offs among Non-Functional Requirements (NFR), i.e. qualities like performance, usability or security. Such alternative designs can be precisely modelled using Goal-Oriented Requirements Engineering (GORE) notations like KAOS (Dardenne et al., 1993), i* (Yu and Mylopoulos, 1997) or GRL (ITU, 2012).

Complex systems can contain many alternatives with various constraints restricting how they can be combined and also complex impacts on different kinds of properties, especially NFR. The process to select an adequate design can actually be described

as an optimisation problem over the system design space (Mogk, 2014). More precisely, it is a multi-objective optimisation problem given the many, so different qualities of a system needs to simultaneously met. Translating such problems into a single optimisation problem is not recommended because it is difficult to assess the relative importance of those qualities and therefore to combine them into a unique objective function. An alternative approach is to compute a set of solutions known as Pareto-optimal from which trade-offs can be explicitly evaluated (Zhang et al., 2008). Search-based software engineering (SBSE) (Harman, 2007) is a specialised field of software engineering concerned with the application of optimisation tools to this kind of problems. A famous one is the Next Release Problem which is known to be NP-Hard (Bagnall et al., 2001).

The aim of this paper is to explore how to use SBSE to tackle the problem of alternative selections in a multi-objective context by computing the Pareto front containing all solutions explicitly. To support this work, the KAOS notations have been used. However, the same concepts and methods can be applied to the other GORE notations mentioned above. On the SBSE side, we will consider the use of constraint pro-

gramming, more specifically using the Open Source Oscala library which provides, among others, powerful support to multi-objective problems (Oscala Team, 2012). In order to illustrate our approach, we have focused on a well-known RE case study: the meeting scheduler (van Lamsweerde et al., 1995) which has already been used for multi-objective reasoning (Nguyen et al., 2018).

Our paper is structured as follows. First, Section 1 reviews existing work in this area. Then Section 2 presents the goal modelling notation used here and illustrates it on our meeting scheduler case study. Section 3 details our mapping of the variability part of the model into a CP model using Oscala.CP. Section 4 exploits the model to perform single and multi-objective searches and illustrate the result on our case study. Finally, Section 5 concludes and identifies some future work.

2 LITERATURE OVERVIEW

This section reviews different reported works investigating the use of CP to reason about GORE models.

Looking at other standard goal modelling notations, GRL support quantitative reasoning, especially through Key Performance Indicators (KPI) (Barone et al., 2011). A mapping on constraint programming was also developed using the JaCoP library (Luo and Amyot, 2011). However the work is more directed towards the definition of a better declarative semantics of GRL. The exploitation seems limited to the propagation on trees without alternatives designs and does not investigate the possibility to explore the design space as we do.

Related to i^* , an extensive multi-objective framework is available (Nguyen et al., 2018). It formalises the notion of Constraint Goal Models (CGM) to supports different kinds of requirements (nice-to-have, preferences, optimisation, constraints...). It exploits automated reasoning SMT-based solvers to support sound and complete reasoning on such goal models. An Eclipse-based CGM tool is available. However, its optimisation capabilities are limited to a lexicographic approach, i.e. with priority ranking among goals. A special attention is devoted to ensuring the reasoning scales on large models derived from the meeting scheduler also used in our work.

Another tool called WebREd-Tool was developed for the i^* method (Aguilar et al., 2011; Calderon et al., 2012). It is a set of Eclipse plugins assisting the designer in the early phases of a Web application development process. It helps the designer to compare different configurations of functional requirements,

while balancing and optimising non-functional requirements. The algorithm relies on Pareto efficiency and provides both tabular and radar based visualisations. The algorithm itself does not seem efficient and scalable as it explicitly stores all possible configurations. However, the authors performed a controlled experiment that confirm requirement engineers may perform better at the task of selecting FRs while optimising and balancing NFRs using the proposed approach and visualisation (Zubcoff et al., 2018).

Feature models are quite close to goal models in terms of structure, attributes, constraints. Complete mapping of feature models over constraint logic programming have been defined, e.g. (Karataş et al., 2010) and enable to perform operations such as propagation, emptiness check or enumeration of products. However, they do not relate to requirements and do not answer the problem we tackle here.

Finally, another KAOS approach has been developed to deal with uncertainty in goals (Heaven and Letier, 2011). It relies on a simulation approach using a genetic algorithm which also results in a multi-objective analysis. This approach requires a more complex modelling in which probability distributions must be specified and is more computation intensive. It can be seen as complementary to the present work and can be used to introduce a more realistic goal modelling.

3 GOAL MODELLING OF THE MEETING SCHEDULER

Goals prescribe, at different levels of abstraction, key properties the considered system should achieve. In this paper, the focus is only put on modelling goals although it is usually combined with other models to structure the domain and agent interactions.

KAOS uses several abstraction levels to express goals starting from high-level strategic goals. In the meeting scheduler case study (van Lamsweerde et al., 1995) used as running example, it can be a goal like “Achieve[EfficientSchedulingOfMeetings]” as depicted in Figure 1 as a light blue parallelograms. In KAOS, high-level goals progressively refined into more concrete and operational ones through refinement relationships. A refinement relationship links a parent goal to several subgoals with different fulfilment conditions using either *AND-refinement* (all subgoals need to be satisfied) or *OR-refinement* (a single subgoal is enough, i.e. possible alternatives). The “WHY” and “HOW” questions can be used to conveniently navigate from subgoals to parent goals (why) and from parent goals to subgoals (how).

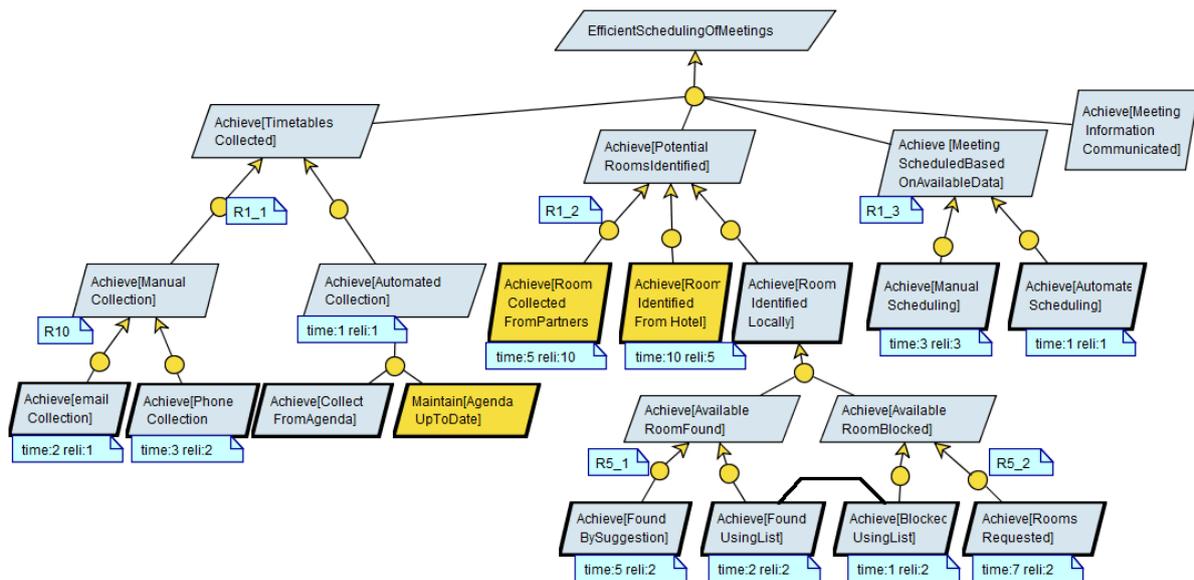


Figure 1: Goal graph for the Meeting Scheduler.

In our case study modelled in Figure 1, the top refinement decomposes the goal “EfficientSchedulingOfMeetings” into the following goals: “Achieve[TimetablesCollected]”, “Achieve [PotentialRoomsIdentified]”, “Achieve [MeetingScheduledBasedOnAvailableData]” and “Achieve[Meeting InformationCommunicated]”. This is an *AND-refinement* of the top level goal: the set of subgoals forms a chain of milestones that yields our top level goal. Such refinements are spotted with yellow circular connectors in our model.

As depicted in Figure 1, the first level of refinement for our case study is composed of the goals: “Achieve[TimetablesCollected]”, “Achieve [PotentialRoomsIdentified]”, “Achieve [MeetingScheduledBasedOnAvailableData]” and “Achieve[Meeting InformationCommunicated]”. Those goals form a *AND-refinement* of the higher level goal: the set of subgoals forms a chain of milestones that yield our top level goal. Such a refinement can be spotted by a yellow circular connectors in our model.

Our model contains many *OR-refinements* that can be detected as yellow circular connectors departing from the same parent goal. For example, the goal “Achieve[TimetablesCollected]” can be satisfied either by “Achieve[ManualCollection]” or “Achieve[AutomatedCollection]”. Goals may also have more than two alternatives. For example, the goal “Achieve[PotentialRoomsIdentified]” has three alternatives. Some alternatives may be restricted to specific choice done in another one. For example, the alternatives “Achieve[FoundUsingList]” and “Achieve[BlockedUsingList]” need to be selected to-

gether.

In this paper, we will assume only one alternative can be selected in a configuration and will discuss this again towards the end of our paper. A design configuration can then be represented as a vector $D = \{R_1, \dots, R_N\}$ where R_i is an integer in the range $0..N_i$ that represents the selected alternative. 0 means no alternative selected and is only relevant for a sub-alternative that is not in the scope (i.e. subtree) of a selected alternative, e.g. “Achieve[EmailCollection]” and “Achieve[PhoneCollection]” is not relevant when selecting “Achieve[AutomatedCollection]”. The identification number of each alternative is captured in the model. They are graphically represented in Figure 1 from left to right. For example, R1_1 has two alternatives: #1 is “Achieve[ManualCollection]” and #2 is “Achieve[AutomatedCollection]”.

The functional goal decomposition stops when reaching a goal controllable by an agent, i.e. answering the “WHO” question about responsibility assignment. These goals are either requirements (depicted like goals but with a thicker border) on the software or expectations on the behaviour of agents in the environment (depicted as requirements but with a yellow background).

The satisfaction of goals can be assessed by their contribution to different several qualities usually modelled as non-functional requirements (NFR) using fulfilling standard classifications such as (ISO, 2011). Each functional requirement can contribute to one or several NFR. This should be preferably measured in a well-defined domain units, either generic

(e.g. reliability of a goal can be measured as the ratio of number of failures vs. total number of occurrences) or problem-specific (e.g. usability can be assessed as the amount of manual work that needs to be done). NFR satisfaction on higher goals can usually be assessed as a combination of the contribution of lower level nodes (eventually possibly leaf requirements) using propagation rules. For an *OR-refinement*, the maximal value is selected as we look for the optimal (as maximum). For an *AND-refinement*, it will depend on the type of quality assessed:

- many properties related to resources usage can be additive on such refinements, e.g. our usability metric measured in hours.
- reliability is not additive but multiplicative, especially when considering chains of tasks as in our model structure
- user satisfaction will generally be reflected by the minimal level of satisfaction of a goal refinement.

A more complete framework to reason about the partial satisfaction of goals was also developed by (Letier and van Lamsweerde, 2004).

Figure 1 shows two quality attributes: time and reliability for leaf-level goals. From there we can infer higher level values using the above rules. For example, at top level the total cost will be the sum of the cost of R1_1, R1_2 and R1_3.

4 TRANSLATING GOAL VARIANTS INTO A CP MODEL

Translating the KAOS model into an OscaR.CP model requires three main steps

- extracting all alternatives from the model
- generating specific constraints across alternatives
- for each target quality: producing an objective function

Note the target CP model is written in the same language as OscaR, i.e. the Scala programming language which is a very expressive general-purpose programming language providing support for functional programming (Odersky, 2004).

4.1 Extracting Alternatives

Alternatives are extracted by querying our model repository. In our experimental setting, we are using the Objectiver tool which can be queried using OQL (Respect-IT, 2005). We also use some post-processing to only keep goals with more than

one *AND-refinement*, that is, goals with an *OR-refinement*. Such queries can easily be adapted to other contexts. The result is shown in Listing 1. Each *OR-refinement* is given a unique identifier R1_1, etc.), is labelled with the parent goal, and decorated with the size of the *AND-refinement* set composing the *OR-Refinement* (named in short in the following, the size of the *OR-Refinement*).

Listing 1: Listing of alternatives with fan-out.

| | | |
|------|---|---|
| R1.1 | Achieve[TimetablesCollected] | 2 |
| R1.2 | Achieve[PotentialRoomsIdentified] | 3 |
| R1.3 | Achieve[MeetingScheduledBasedOnAvailableD.] | 2 |
| R5.1 | Achieve[AvailableRoomFound] | 2 |
| R5.2 | Achieve[AvailableRoomBlocked] | 2 |
| R10 | Achieve[ManualCollection] | 2 |

In the following, each *OR-refinement* in the goal graph will have a state. The values of the states are:

- *Disabled*: the *OR-refinement* has been discarded and none of its alternatives can be selected in the design
- *Enabled*: the *OR-refinement* is not discarded and one of its alternatives may be selected in the design (but none is currently selected)
- *Selected*: the *OR-Refinement* is enabled and one alternative is selected in the design

Based on this, for each alternative, a variable is declared, with the identifier of an *OR-Refinement* as a name and with a domain ranging from 0 to the size of the *OR-refinement* as computed in the previous step. A special kind of variable called *CPIntVar* capturing the specified range is used. A strictly positive value means the *OR-refinement* is enabled and the value tells the alternative that is selected. A “0”-value means the alternative is disabled.

4.2 Generating Specific Constraints

In what follows, we will call *sub – alternative*, an alternative in the subtree of another alternative.

The following specific constraints express that at the end of the computation a complete and consistent design has been built:

- *completeness*: all top level alternatives *OR-refinements* are enabled selected (i.e. all associated variables $?!=0$)
- *consistency#1*: if an alternative is not selected, all its sub- alternatives should must be disabled (by recursion, it is enough to disable the closest sub-alternatives using refinement links)

- *consistency#2*: if an alternative is selected, its sub-alternatives should be enabled (again the same recursion remark holds)
- *domain specific*: e.g. if some alternatives need to be selected together or if two alternatives are mutually exclusive.

The computation of those constraints is quite easy to translate using Boolean equality/equivalence `?===`, difference `?!=`, implication `==>`, conjunction `&&` and disjunction `||` operators as depicted in Listing 2.

Listing 2: CP model for goal variants and constraints.

```

val R1_1 = CPIntVar(0 to 2)
val R1_2 = CPIntVar(0 to 3)
val R1_3 = CPIntVar(0 to 2)
val R5_1 = CPIntVar(0 to 2)
val R5_2 = CPIntVar(0 to 2)
val R10 = CPIntVar(0 to 2)

// toplevel alternatives
add(R1_1 ?!= 0)
add(R1_2 ?!= 0)
add(R1_3 ?!= 0)

// enabling and disabling constraints
add((R1_2 ?!= 3)
    ==> ((R5_1 ?=== 0) && (R5_2 ?=== 0)))
add((R1_2 ?=== 3)
    ==> ((R5_1 ?!= 0) && (R5_2 ?!= 0)))

// domain specific constraints
add((R1_1 ?!= 1) ==> (R10 ?=== 0))
add((R1_1 ?=== 1) ==> (R10 ?!= 0))
    
```

Based on this set of constraints, the list of alternatives can be easily produced using the commands of Listing 3. It generates 24 different configurations.

Listing 3: Optimisation search.

```

search {
  binaryStatic(
    Seq(R1_1, R1_2, R1_3, R5_1, R5_2, R10))
} onSolution {
  println("R1_1:" + R1_1 + ... )
} start()
    
```

4.3 Objective Function

Objective functions are also encoded as *CPIntVar* and can be expressed using expressions involving our variables. The rules described in Section 2 are applied to produce evaluation at all levels enabling leaf values to propagate up to the top level based on the selected alternatives. The alternative selection operator is simply the truth value (`RX?===expected_value`)

which will evaluate to 0 when false and to 1 when true. Note that we consider only refinements without any shared goals which deserve a special attention and are discussed at the end of this section.

For example, our time quality can be encoded as shown in Listing 4.

Listing 4: Optimisation search.

```

def time_R1_1 : CPIntVar =
(R1_1 ?=== 1) * ((R10 ?=== 1) * 2 + (R10 ?=== 2) * 3)
+ (R1_1 ?=== 2) * 1

def time_R1_2 : CPIntVar =
(R1_2 ?=== 1) * 5 +
(R1_2 ?=== 2) * 10 +
(R1_2 ?=== 3) * ((R5_1 ?=== 1) * 5 + (R5_1 ?=== 2) * 2
+ (R5_2 ?=== 1) * 1 + (R5_2 ?=== 2) * 7)

def time_R1_3 : CPIntVar =
(R1_3 ?=== 1) * 3 + (R1_3 ?=== 2) * 1

val time = (time_R1_1 + time_R1_2 + time_R1_3)
    
```

At this point, it is easy to perform a single objective search, for example minimising the time can be done inserting the command `minimize(time)` before the search. It will yield the result of Listing 5

Listing 5: Optimisation search.

```

objective tightened to 10 lb:-1
R1_1:1 R1_2:1 R1_3:1 R5_1:0 R5_2:0 R10:1
objective tightened to 8 lb:-1
R1_1:1 R1_2:1 R1_3:2 R5_1:0 R5_2:0 R10:1
objective tightened to 6 lb:-1
R1_1:1 R1_2:3 R1_3:2 R5_1:2 R5_2:1 R10:1
objective tightened to 5 lb:-1
R1_1:2 R1_2:3 R1_3:2 R5_1:2 R5_2:1 R10:0
    
```

5 MULTI-OBJECTIVE OPTIMISATION

Performing a multi-objective search is also quite easy thanks to the multi-objective search function and Pareto visualisation contributed by (Hartert and Schaus, 2014). After defining a second objective function, the optimisation command before the search becomes `paretoMaximize(obj1, obj2)` as shown in Listing 6. Note that the first objective function was turned into a maximisation by computing `20 - time` (i.e. time resource left) just to ensure both functions need to be maximised.

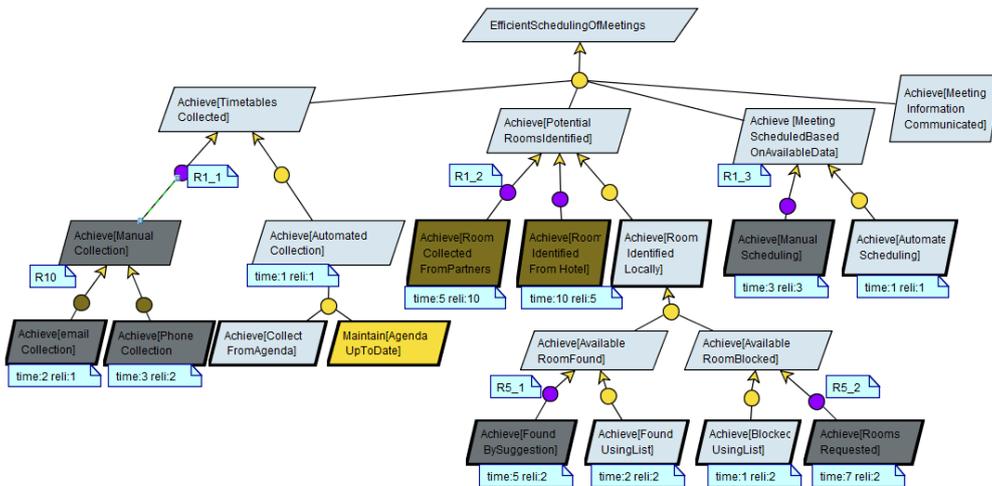


Figure 3: Visual representation of the selected design.

Listing 6: Multi-objective search (for costs see Listing 6).

```

def reli_R1_1 : CPIntVar =
(R1_1?===1)*((R10?===1)*1+(R10?===2)*2)+
(R1_1?===2)*1

def reli_R1_2 : CPIntVar =
(R1_2?===1)*10+
(R1_2?===2)*5+
(R1_2?===3)*2

def reli_R1_3 : CPIntVar =
(R1_3?===1)*3+(R1_3?===2)*1

val obj1=time*(-1)+20
val obj2=reli_R1_1*reli_R1_2*reli_R1_3

cp.paretoMaximize(obj1, obj2)

search {
  binaryStatic(
    Seq(R1_1, R1_2, R1_3, R5_1, R5_2, R10))
  } onSolution {
    paretoPlot.insert(obj1.value,
                      obj2.value)
  } start()
    
```

Figure 2 depicts both the complete search as a cloud of points and the resulting Pareto front which can also be presented in tabular form with more information enabling the selection (see Table 1). The reliability information has been normalised to a real value between 0 and 1 without impacting the solutions but to better match conventional representation for this attribute. The plot and the table can be used to guide the analyst in selecting a right solution, e.g. by ruling out too costly alternatives, unreliable ones or making sure a specific *OR-refinement* is selected.

Table 1: Tabular view of Pareto solutions.

| R1_1 | R1_2 | R1_3 | R5_1 | R5_2 | R10 | Cost | Reliability |
|------|------|------|------|------|-----|------|-------------|
| 1 | 1 | 1 | 0 | 0 | 2 | 9 | 0,95 |
| 2 | 1 | 1 | 0 | 0 | 0 | 11 | 0,8 |
| 2 | 1 | 2 | 0 | 0 | 0 | 13 | 0,7 |
| 2 | 3 | 2 | 2 | 1 | 0 | 15 | 0,66 |

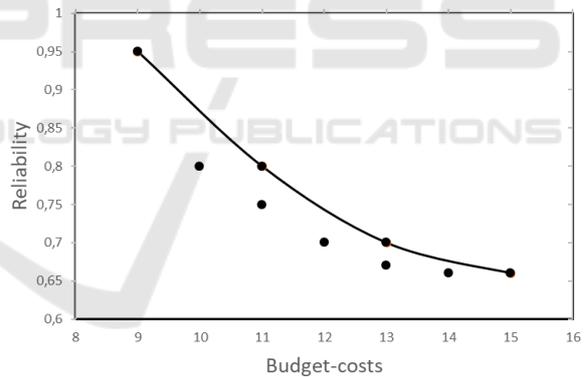


Figure 2: Pareto front.

The tool also provides a filtering mechanism to grey out the part of the model not covered by the selected alternative. Figure 3 highlights the instance selected by the time minimisation reported earlier.

5.1 Discussion

Shared goals require some extra processing in order to avoid counting them twice. Actually, a shared goal may have a contribution of its own to be counted once and a contribution for being involved in one or more refinements. Making sure the goal is counted once can be encoded through selection checking operators referring to the other alternatives involving it.

About the readability of the constraints, although the generated code is meant to be executed by an optimisation engine, the selection operators could be stated in more user-friendly form, e.g. using `isSelected(alt)` rather than `"alt?!" == 0`. Operators for checking other states can also be introduced.

The current mapping is limited to the selection of a single alternative which is usual for goal refinements but not for the dual of goals, i.e. obstacles. In this case, making sure an obstacle is managed may require to select multiple *OR-refinements*. The current mapping cannot cope with this but could be adapted, e.g. an array of *CPBoolVar*, and specific constraints for imposing *no*, *one* or *some* alternatives. This results in the generation of far larger design spaces which also raises the need for better guidance.

6 CONCLUSION & NEXT STEPS

In this paper, we developed an approach to deal with the exploration of the design space at requirements engineering time. As multiple qualities need to be assessed simultaneously, a multi-objective approach computing a Pareto front was used to provide the analyst with a candidate set of solutions which can then be reviewed more carefully. Our translation from the KAOS notation to the Oscar.CP was tested on the meeting scheduler benchmark.

Our work is still quite in early phase. Our next steps are to cover the different extensions identified such as better language primitives, dealing with shared goals and allowing the selection of multiple alternatives. We also plan to improve usability and to validate on larger models and with external analysts.

ACKNOWLEDGEMENTS

Thanks to Respect-IT for giving access to the Objectiver tool and to the Oscar team for its Open Source optimisation library.

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