Implementation of Trajectory Planning for Automated Driving Systems using Constraint Logic Programming

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Abstract: Automated driving systems are a maturing technology that is considered to have a significant impact on mobility. Trajectory Planning is a safety-critical task that plays an important role in automated driving systems. In this paper, we present the implementation of a trajectory planning module called CLPT (CLP_TRAJECTORY_PLANNER) using constraint logic programming (CLP) and evaluate it in simulated traffic situations. CLP allows us to express the constraints of the problem of trajectory planning in a declarative way. The approach makes the code less complex and more readable for domain experts compared to code using an imperative programming language. Compared to approaches making use of neural networks to manage the complexity of the problem of trajectory planning, the results of CLPTP are more comprehensible and easier to verify. Thus CLPTP can be seen as a step towards solving the problem of trajectory planning with explainable artificial intelligence. An evaluation of the execution time performance of our implementation shows that further research is needed to apply the approach in real world vehicles.

1 INTRODUCTION

In the past few decades, the level of automation in road vehicles has been steadily increasing, resulting in a significant decrease of traffic accidents (ITF, 2017). Not only safety, but also accessibility, convenience and efficiency of automotive mobility are expected to be enhanced by automated driving systems (Paden et al., 2016; Arbib and Seba, 2017).

Although many different automated driving systems are known to exist in research and development projects, most of them share a similar functional architecture in which the driving task is split in three parts: Perception, decision and control, vehicle platform manipulation (Behere and Törngren, 2015). Decision and control is typically divided into three levels: route planning, behavioural decision making and trajectory planning (Paden et al., 2016).

It is the responsibility of the trajectory planning level to find a safe, comfortable and dynamically feasible trajectory for the vehicle in its dynamic environment. Due to space limitations, this paper focuses on trajectory planning. Other related problems such as object movement prediction or trajectory execution are not discussed in this paper. The problem of trajectory planning can be formulated as an optimization problem with constraints imposed by traffic rules, the dynamic environment, vehicle dynamics etc. Many algorithms have already been developed for solving this problem. Some of them represent vehicle states in the form of a graph and use search algorithms like A* (Hart et al., 1968) for finding a minimum cost path on the graph (Ziegler and Stiller, 2009; Gu and Dolan, 2012). Other solutions describe trajectories as polynomials that are optimized under consideration of a cost functional (Werling, 2011; McNaughton, 2011; Kelly and Nagy, 2016; Rathgeber, 2016). Aside from that, there are other implementations e.g. based on neural networks and reinforcement learning (Dubey et al., 2013; Zuo et al., 2014; Grigorescu et al., 2017). Refer to (Paden et al., 2016) for a detailed presentation of the trajectory planning problem and an extensive survey of motion planning techniques for self-driving vehicles.

Translating the aforementioned constraints of the problem of trajectory planning into software code can be a challenging task when using imperative programming languages because the specification of the constraints are often present as declaratively written requirements in natural language or in formal or semi-
formal notations. Moreover, operational details in imperative code make it hard to read for domain experts and more difficult to verify, which is an important aspect for safety-critical software. Approaches using neural networks overcome the challenging programming task by learning desired planning behaviour implicitly. However, verifying neural networks for safety-critical systems, where international standards such as ISO 26262 (ISO, 2011) must be applied, is still an open problem (Grigorescu et al., 2017).

In this paper, we present the design and implementation of a trajectory planning module called CLPTP (CLPTTRAJECTORYPLANNER). CLPTP uses constraint logic programming (CLP), a programming paradigm enhancing the declarative paradigm of logic programming by mechanisms for specifying and solving constraints. By using this programming paradigm with its high level of abstraction, program code is expected to be easier to understand and more concise than its imperative counterpart. The planning results are expected to be more comprehensible and easier to verify than the output of a neural network, thus supporting the trend for explainable artificial intelligence (Biran and Cotton, 2017) with a different approach than e.g. (Bojarski et al., 2017).

Logic programming has already been used in (Piaggio and Sgorbissa, 2000) for the symbolic component of an autonomous robot navigator. CLP has been successfully used to implement programs for solving general planning and optimization problems (Lever and Richards, 1994; Apt and Wallace, 2006). However, no attempt to apply CLP to the specific problem of trajectory planning seems to exist yet.

In Section 2, we present the background of trajectory planning as far as needed here. In Section 3, the design and implementation of CLPTP is described, Section 4 contains an evaluation of the system, and Section 5 concludes and points out further work.

2 TRAJECTORY PLANNING

Modeling Mobility of a Vehicle. Mobility of a vehicle can be modeled beginning with the notion of a vehicle configuration, representing its position in the world (Paden et al., 2016). In this paper, configuration is expressed as the planar coordinate of the vehicle center. Rather than directly using cartesian coordinates in some world coordinate system with a fixed base or the body frame of the vehicle, we use the dynamic Frenet frame (Bauchau, 2011) as the coordinate system for the vehicle configuration space. The Frenet frame is commonly used for trajectory planning implementations. As depicted in Figure 1, it is given by the tangential and normal vector at a certain point of some curve referred to as the centerline in the following (Werling et al., 2010). The centerline either represents an ideal path along a free lane or the result of a path planning algorithm for unstructured environments (Werling et al., 2010). Hence a position in the configuration space of the vehicle is expressed as a longitudinal offset \( s \) along the straightened centerline and a perpendicular lateral offset \( d \).

The Problem of Trajectory Planning. Trajectories prescribe the evolution of the configuration of a vehicle in time. So in theory, a trajectory is a continuous time-parameterized function \( \pi(t):[0,T] \rightarrow X \), where \( T \) is the planning horizon and \( X \) is the configuration space of the vehicle (Paden et al., 2016). The trajectory planning problem is to find a trajectory that starts at the initial vehicle configuration \( x_{init} \in X \) and satisfies given requirements or constraints.

Trajectory planning for automated vehicles has a high computational cost for most real world scenarios because of their complexity resulting from the number of dynamic and static objects. One possibility to reduce the computational cost is to split the planning into a general planning and a detailed planning (Rathgeber, 2016). In general planning, the whole available space is considered (e.g. multiple lanes on a highway) and a reduced target space is determined for the trajectory. In detailed planning, only this reduced space has to be considered for the optimization of the exact trajectory shape.

Because of space limitations and the goal of this paper to rather demonstrate the application of CLP to the problem of trajectory planning instead of developing a new algorithm, we only present the implementation of a general trajectory planning module. The presented approach is also applicable to detailed trajectory planning, but rather demanding on resources (cf. Section 4).

Our implementation strongly discretizes the con-
Table 1: Quantization of the 4-dimensional configuration space of the vehicle.

<table>
<thead>
<tr>
<th>Coordinate</th>
<th>Quantization step size</th>
</tr>
</thead>
<tbody>
<tr>
<td>s (longitudinal offset)</td>
<td>1 cm</td>
</tr>
<tr>
<td>d (lateral offset)</td>
<td>1 cm centerline</td>
</tr>
<tr>
<td>v (longitudinal velocity)</td>
<td>1 cm² s⁻¹</td>
</tr>
<tr>
<td>a (longitudinal acceleration)</td>
<td>20 cm³ s⁻¹</td>
</tr>
</tbody>
</table>

The configuration space of the vehicle in lateral direction so that only the centerline or lines parallel to the centerline (e.g., on a multi-lane road) are considered as valid configurations. This approximation is sufficient for general trajectory planning and increases execution performance. Time is discretized in a way that vehicle configurations can be planned at arbitrary discrete time instances \( t_k = k \cdot \Delta t \) with \( \Delta t = 1 \) s and \( k \in \mathbb{N} \).

As planned vehicle configurations are supposed to be eventually passed to a vehicle controller for execution, they do not only contain position coordinates in our implementation, but also values for velocity and acceleration in longitudinal direction. This means that the configuration space is 4-dimensional. As a simplification for ease of demonstration, we assume centerline changes to be completed within one configuration transition, so lateral velocity, steering etc. must not be considered. Furthermore, both centerlines are assumed to be blocked on the occasion of a (lane) change.

For efficiency reasons, the implementation presented here uses CLP over finite domains (CLP(FD)) instead of e.g., CLP over reals (CLP(R)). Thus, constraint variables can only be bound to integer values, so the vehicle configuration space must be quantized considering the conflicting goals of high accuracy and small configuration space. Table 1 summarizes our choice of configuration space quantization.

3 DESIGN AND IMPLEMENTATION OF CLPTP

System Overview. With logic programming, elegant programs solving complex problems can be implemented by applying the generate-and-test technique. Such a program consists of one (often very simple) predicate generating possible solutions and another predicate checking the solution for correctness. Thereby the search for a solution is accomplished by the interpreter’s backtracking mechanism rather than by an explicitly stated (and often difficult to find) algorithm. It is easy to realize that this approach is infeasible for problems with big solution spaces in practice. With CLP, the generate-and-test technique can be improved to a constrain-and-generate pattern, where the solution space is narrowed down by constraints, thereby increasing the efficiency of solution search or optimization, but still separating concerns very well. Therefore, the general structure of our implementation is based on this programming technique.

All code presented in this paper has been developed using SICStus Prolog (Carlsson and Mildner, 2012). SICStus Prolog supports CLP(FD). Whenever we mention predefined predicates or interpreter-specific details, we refer to SICStus Prolog. With respect to our approach, other Prolog variants supporting CLP(FD) like SWI-Prolog (Wieliemaker et al., 2012) or GNU-Prolog (Diaz, 2001; Diaz et al., 2012) are very similar in most aspects.

Before starting to actually define predicates that describe the problem of trajectory planning, we need to define how to represent trajectories, vehicle configurations and environmental elements such as objects.

A trajectory is represented by a list of vehicle configurations at discrete time instances. Vehicle configurations are represented by program terms \( \text{conf}(T, D, S, V, A) \), where \( T \) denotes the discrete time instance this configuration belongs to and \( D, S, V \) and \( A \) are the four dimensions of the quantized configuration space as outlined in Table 1. General vehicle parameters, quantization factors etc. are expected to be supplied by the predicate \( \text{param}/2 \) that implements some kind of key-value-store. Parameters could also be time- or position-parameterized in principle. Lateral boundaries of the configuration space (i.e., road boundaries in a particular planning situation) are modeled by two predicates \( \text{left Boundary}/2 \) and \( \text{right Boundary}/2 \) dependent on the longitudinal offset \( s \). Recognized objects are modeled in the clause database as facts \( \text{object}(ID, T, D, SR, SF, V) \), where \( ID \) uniquely identifies an object, \( T \) is the time the recognition is valid for, \( D \) denotes the lateral coordinate (i.e., the centerline) of the object, \( SR/SF \) are the longitudinal offsets of the rear/front object boundaries, respectively, and \( V \) denotes the longitudinal velocity of the recognized object. Speed limits are modeled similar to objects by a predicate \( \text{speed limit}(S, V) \). \( S \) is the longitudinal offset from where on the limit with value \( V \) is valid.

Now we can describe the two main predicates our implementation consists of. They are given in Listing 1. \( \text{optimal trajectory}/3 \) is the high-level predicate expected to be called by the user. It relates an initial vehicle configuration and a list of discrete time instances to an optimized trajectory. It is based on the constrain-and-generate technique explained above. Two helper predicates \( \text{trajectory}/4 \) and \( \text{variable order}/2 \) are used to, firstly, prescribe a feasible trajectory and its costs for optimization and,
secondly, order the constraint variables of a list of vehicle configurations (i.e. a trajectory) in a specific way that can be optimized for an efficient labeling strategy. These helper predicates can be implemented with fairly simple Prolog code and are therefore not outlined in detail here. transition/3 constrains the variables of the second of two consecutive vehicle configurations and relates costs to this transition in the configuration space. It makes use of four helper predicates vehicle_motion/2, collision_avoidance/3, speed_limit_compliance/3 and costs/3 that encapsulate the implementation of various requirements or constraints of trajectory planning. The implementation of these predicates is explained in detail hereafter.

Listing 1: Predicates implementing general trajectory planning using CLP(FD). The predicates findall/3, setof/3, minimize/2 and labeling/2 are part of the modules aggregate clpfd provided by SICStus Prolog.

```prolog
optimal_trajectory (InitialConf, TimeInstances, Trajectory) :-
  trajectory (TimeInstances, InitialConf, Trajectory, Costs),
  sum (Costs, #=, TotalCosts),
  variable_order (Trajectory, Vars),
  minimize (labeling ([leftmost, bisect], Vars), TotalCosts).

transition (Conf, NextConf, Cost) :-
  NextConf = conf (Te, De, Se, Ve, Ae),
  param (Vmax, Vmax), param (Vmin, Vmin),
  Ve in Vmin .. Vmax,
  param (Amax, Amax), param (Amin, Amin),
  Ae in Amin .. Amax,
  right_boundary (Se, Dmin),
  De in Dmin .. Dmax,
  vehicle_motion (Conf, NextConf),
  findall (ObjD, ObjRear, ObjFront, ObjV),
  object (_, Te, ObjD, ObjRear, ObjFront, ObjV),
  objects ),
  collision_avoidance (Conf, NextConf, Objects),
  setof ([SLimit, Limit], speed_limit (SLimit, Limit), Vlimits),
  speed_limit_compliance (Vlimits, Se, Ve),
  costs (Conf, NextConf, Cost).
```

Vehicle Motion. The most basic and important relation between two vehicle configurations is given by the motion of the vehicle. As we are not implementing detailed but general trajectory planning, acceleration of a configuration is considered to have the meaning of a constant average acceleration for the time period of the transition leading to the configuration. Considering the already mentioned simplifications we can thus characterize vehicle motion by the basic equations of motion for point masses with constant acceleration $a$:

$$
\begin{align*}
\dot{s}_r(t) &= 0 \\
\dot{s}_r(t) &= a \\
\ddot{s}_r(t) &= at
\end{align*}
$$

Listing 2 outlines a translation of these equations into code using CLP(FD). Note that these constraints already take into account the quantization presented in Table 1.

Listing 2: Implementation of constraints among the variables of two consecutive vehicle configurations resulting from the motion equations of a very simple vehicle model.

```prolog
vehicle_motion (conf (Te, De, Se, Ve, Ae)),
conf (Te, De, Se, Ve, Ae) :-
  param (jmax, Jmax), param (jmin, Jmin),
  param (aquant, AQuant),
  DT #= Te - T0,
  Ae #< A0 + (Jmax / AQuant) * DT,
  Ae #> A0 + (Jmin / AQuant) * DT,
  Ve #= V0 + Ae * AQuant * DT,
  De #= D0 + (Ve * (AQuant / 2) * DT) * DT.
```

Collision Avoidance. After the application of basic motion constraints, the configuration space of the to be planned configuration can be narrowed down further by imposing constraints on its position coordinates. These constraints are most likely caused by obstacles or objects such as other traffic participants and therefore implement collision avoidance. An example implementation assuming recognized objects being modeled as described in Section 3 is given in Listing 3. For all recognized objects, the constraints outlined in Listing 3 apply a static safety distance and a safety time resulting in a dynamic velocity-dependent safety distance to the coordinates of the vehicle configuration. The value for the static safety distance must take the vehicle dimensions into account.

Note that time discretization as explained above can cause collisions with small objects in this implementation. Therefore time instances must be chosen carefully. This problem could also be solved by adding an additional constraint regarding the coordinates covered by the transition in the configuration space.

1Labeling is the process of assigning values to constraint variables that satisfy all present constraints. The interpreter’s predefined predicate labeling/2 accepts parameters to configure the labeling strategy (Carlsson, 2009).
Listing 3: Implementation of constraints for the variables of two consecutive vehicle configurations that ensure keeping a static and dynamic safety distance to other objects.

collision_avoidance(_ , _, []). collision_avoidance (conf{To, D0, S0, V0, A0 }, conf{Te, De, Se, Ve, Ae}, [ObjD, ObjRear, ObjFront, ObjV]) :-
  param(safety_distance, Ss), param(safety_time, Ts),
  Smi = ObjRear - Ss - Ts * Ve,
  Smax = ObjFront + Ss + Ts * ObjV,
  (De = ObjD) #=> ((Se < Smin) \
                   { Se > Smax}),
  ((De = ObjD) #\ {De = D0}) #=> ((
    S0 < Smin) \
    {S0 > Smax}),
  collision_avoidance (conf{To, D0, S0, V0, A0}, conf{Te, De, Se, Ve, Ae}, Objs).

Speed Limits. By considering motion constraints and collision avoidance, important requirements for trajectory planning are satisfied. However, other constraints have to be taken into account as well, e.g. the planned trajectory has to be compliant with traffic rules. Listing 4 shows a possible implementation of speed limit constraints as an example on how to introduce traffic rules using CLP.

Listing 4: Constraints for the vehicle configuration variables S (longitudinal coordinate) and V (longitudinal velocity) implementing compliance with a list of speed limits. A speed limit is expected to be represented by a coordinate from where on the limit is valid and the value of the limit.

  (S => SLimit) #=> (V <= VLimit).
  speed_limit_compliance([(S1, VLimit1), (S2, VLimit2) | Limits], S, V) :-
  (S => S1 \ S < S2) #=> (V < VLimit1),
  speed_limit_compliance([(S2, VLimit2) | Limits], S, V).

Costs. Until now, we have specified and constrained the relation between a number of discrete vehicle configurations (represented by constraint variables over finite domains) representing a planned trajectory (the expected output of our general trajectory planning module). We could now use the implemented program to retrieve a feasible trajectory, i.e. bindings for all the constraint variables, by using a predicate such as labeling/2 provided by SICStus Prolog.

However, the goal of trajectory planning is to find an optimal trajectory, not just any feasible trajectory. Therefore we relate each vehicle configuration transition to some costs. These costs are combined to total costs which are used as the goal in the predicate minimize/2, as it can be seen in Listing 1. On the execution of that predicate, the constraint variables representing vehicle configurations are unified with values corresponding to the optimal solution for the queried trajectory with respect to the specified costs. In Listing 5, the implementation of a predicate prescribing the relation between costs and two consecutive discrete vehicle configurations is given. This implementation considers the following costs each at which is assigned a weight:

- **C1**: difference between planned velocity and desired velocity. Desired velocity is assumed to be the maximal velocity here but can of course be any value e.g. set by the vehicle user.
- **C2**: acceleration
- **C3**: change in acceleration (jerk)
- **C4**: lane change
- **C5**: deviation from a measure representing a kind of strategic lane advice determined by the upper layers of decision and control. It is assumed that this measure is represented by a value between 0 and 100 available via the predicate advice/3 for every position that is relevant for the current planning cycle.

These costs and the weights assigned to them in Listing 5 serve as demonstrating examples in this paper. In practice, they should be determined by theoretical and empirical analysis.

Listing 5: Implementation of the predicate costs/3, giving the costs of a transition from one vehicle configuration (indexed 0) to another vehicle configuration (indexed e).

costs (conf{T0, D0, S0, V0, A0}, conf{Te, De, Se, Ve, Ae}, Costs) :-
  advice(De, Se, Advice),
  param(vmax, VDesired),
  C1 = 50 * abs(Ve - VDesired),
  C2 = 100 * abs(Ae),
  C3 = 100 * abs(Ae - A0),
  C4 = 10 * abs(De - D0),
  C5 = 30 * (100 - Advice),
  sum([C1, C2, C3, C4, C5], #=, Costs).

4 EVALUATION

Functional Evaluation in Simulation. In this section, trajectories returned by the presented general trajectory planning module for two simulated traffic situations are visualized and discussed.

The first situation is visualized in Figure 2. Here, the simulated environment is a road section with four lanes that are separated by dashed lane markings.
Hence, lane changes are allowed. Two speed limits exist at $s_r = 0 \text{ m}$ and $s_r = 50 \text{ m}$, respectively. Furthermore, there are four dynamic objects in addition to the ego vehicle moving with different velocities. A simple object prediction assuming objects to continue on their current lane with their current velocity is applied for ease of demonstration. Predicted object positions are visualized for the discrete time instances $2\text{s}$, $4\text{s}$, $6\text{s}$ and $8\text{s}$, for which planning is executed. At $t = 0\text{s}$, the vehicle is in initial configuration. Values for strategic lane advice and important parameters such as $v_{\text{max}}$ etc. are given in Figure 2 as well.

It can be seen that many requirements mentioned in the previous section are fulfilled: Predictions of other objects are considered. Strategic advice is taken into account by a lane change. In general, velocity is planned as high as possible. The second speed limit results in an appropriate deceleration, but note that the implementation in Listing 4 allows the vehicle to have a velocity too high for the duration of one transition. This could be adapted by an additional constraint for the variable representing the velocity of the previous configuration.

The second situation is visualized in Figure 3. It is similar to the first situation but with the second speed limit removed. Here, trajectory planning has been queried for the time instances $3\text{s}$, $6\text{s}$, $9\text{s}$ and $12\text{s}$ so that an overtaking maneuver can be demonstrated. Due to the low velocity of the objects on the two rightmost lanes, the results show that it is favourable to change to the second leftmost lane despite its low strategic advice in order to obtain a higher velocity. This planning behaviour can of course be influenced by the design of the costs that have been discussed in

![Figure 2: Resulting trajectory (blue) calculated by CLPTP for a simulated traffic situation consisting of 4 lanes, 4 other dynamic objects with predicted movements (dark rectangles) and 2 speed limits (red). Planning has been queried for the discrete time instances $2\text{s}$, $4\text{s}$, $6\text{s}$ and $8\text{s}$. At $t = 0\text{s}$, the vehicle is in initial configuration.](Image 73x306 to 523x535)

![Figure 3: Resulting trajectory (blue) calculated by CLPTP for a simulated traffic situation consisting of 4 lanes, 4 other dynamic objects with predicted movements (dark rectangles) and 1 speed limit (red). Planning has been queried for the discrete time instances $3\text{s}$, $6\text{s}$, $9\text{s}$ and $12\text{s}$. At $t = 0\text{s}$, the vehicle is in initial configuration.](Image 73x306 to 523x535)

<table>
<thead>
<tr>
<th>Traffic situation</th>
<th>Execution time</th>
</tr>
</thead>
<tbody>
<tr>
<td>Situation 1 (Figure 2)</td>
<td>1427 ms</td>
</tr>
<tr>
<td>Situation 1 with $a_{\text{max}} = 1 \text{ m}^2\text{s}^{-1}$</td>
<td>360 ms</td>
</tr>
<tr>
<td>Situation 2 (Figure 3)</td>
<td>500 ms</td>
</tr>
<tr>
<td>Situation 2 with $a_{\text{max}} = 1 \text{ m}^2\text{s}^{-1}$</td>
<td>219 ms</td>
</tr>
</tbody>
</table>

Section 3.

**Execution Time.** Due to the dynamic environment of real world automated vehicles, trajectories must be planned cyclically at a high enough rate to account for any changes in the environment. Thus execution time is a critical characteristic for any trajectory planning program.

Table 2 gives a summary of execution time measurements for CLPTP. The results show that the computational cost of our implementation is too high for real world vehicles, where typical planning frequencies lie in the range of 10 Hz to 20 Hz. However, this does not mean that Prolog or CLP in general is inappropriate for solving the problem of trajectory planning, as the performance of the system is solely determined by the labeling, search and optimization algorithms of the underlying interpreter. For example, if it is more appropriate to conduct a breadth-first search (as the A* algorithm does) instead of a depth-first search, this can be implemented in the interpreter without modifying the declarative code it executes. This inherent separation of logic describing a problem to solve and operational control flow can be seen as one of the advantages of CLP over other, e.g. im-
perative, programming paradigms when solving complex problems.

The results in Table 2 also illustrate the heavy influence the size of the quantized configuration space has on execution time performance. Due to the safety-criticality of trajectory planning, this is a significant issue because international standards for functional safety such as ISO 26262 demand worst-case execution time estimations for any safety-relevant software. Worst-case execution times can only be estimated for trajectory planning programs if upper bounds are applied for the number of lanes, dynamic objects, speed limits etc. The automated driving system can then only be applied safely in traffic situations not exceeding the estimated limits and mechanisms (e.g. handover to a human driver) ensuring functional safety in unexpectedly complex situations must be implemented.

Verifiability. Two important methods are commonly used to verify or validate safety-critical systems (ISO, 2011): On the one hand, there are formal techniques like theorem proving. On the other hand, there is systematic testing. Both methods can profit from the characteristics of CLP.

Theorem proving is based on first-order logic (Ray, 2010). Assertions on a program are proved using logic inference. Since constraint logic programs are also based on first-order logic, the application of theorem proving to them is more straightforward compared to imperative programs.

To demonstrate how CLP supports writing systematic tests efficiently, Listing 6 outlines a simple test case for CLPTP. The interpreter executing the test case infers whether the lateral coordinate of the first vehicle configuration of any trajectory calculated by CLPTP lies within certain lane boundaries.

Listing 6: Test case for CLPTP checking that lateral coordinates lie within lane boundaries. Partial data structures and double negation are used to make the test case as generic as possible. The predefined predicate `assertz/1` is used to set up a test scenario. `fd_dom/2` is used to check the domain of the constraint variable `D`.

```prolog
:- assertz(object(_, _, _, _, _, _)).
:- assertz(speed_limit(_, _)).
:- assertz(right_boundary(_, _)).
:- assertz(left_boundary(_, _)).
:- assertz(trajectory(_, _, [conf(_, D, _,-,-,_)]), _, fd_dom(D, -1..2))).
```

Note that the double negation used in the last directive of Listing 6 is the usual Prolog method to prove a test goal for all possible solutions and ensure that variable bindings (in this case `D`) are unchanged after executing the test goal. This test case example uses the clause database of the interpreter for defining a test scenario and shows how partial data structures allow generic tests. The test scenario in Listing 6 is generic in so far that it includes traffic situations with arbitrary objects and speed limits by asserting facts for the predicates `object/6` and `speed_limit/2` that unify with any query. Thereby the independence of staying within lane boundaries on the one side and speed limits and objects on the other side is shown. In other languages, multiple concrete test cases would be needed to achieve an equivalent test coverage.

5 CONCLUSIONS AND FURTHER WORK

The implementation of a general trajectory planning module presented in this paper shows that CLP reduces the complexity of implementing a software solving the problem of trajectory planning. Although several simplifications have been applied for the ease of demonstration, compared to imperative programming languages, the translation of requirements and constraints of the planning problem into program code is more straightforward and easier to accomplish for domain experts. Less operational details in the code and the abstract, declarative way of specifying knowledge about a problem as program code make CLP a promising technology for the implementation of software solving complex planning problems in automated driving systems. Compared to approaches based on neural networks, well-known verification methods demanded by standards such as ISO 26262 can be applied more easily.

The execution time evaluation presented in Section 4 shows that algorithms for constraint solving and optimization currently available in interpreters of CLP languages such as SICStus Prolog are not capable of fulfilling the performance requirements of automated driving systems in real world situations. However, the underlying algorithms of interpreters could be adapted and improved without having to modify the declarative code describing the problem to solve.

Future work should include an investigation on how to improve execution time performance, e.g. by using parallel implementations of constraint solving and optimization algorithms. If execution time can be reduced considerably, the approach of this paper could be applied to implement not just general, but also detailed trajectory planning. Smaller quantization step sizes, high-dimensional configuration spaces and more accurate vehicle models could be used.
REFERENCES


