A Stream Clustering Algorithm using Information Theoretic Clustering Evaluation Function

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Abstract: There are many stream clustering algorithms that can be divided roughly into density based algorithms and hyper spherical distance based algorithms. Only density based algorithms can detect nonlinear clusters and all algorithms assume that the data stream is an ordered sequence of points. Many algorithms need to receive data in buckets to start processing with online and offline iterations with several passes over the data. In this paper we propose a streaming clustering algorithm using a distance function which can separate highly nonlinear clusters in one pass. The distance function used is based on information theoretic measures and it is called Clustering Evaluation Function. The algorithm can handle data one point at a time and find the correct number of clusters even with highly nonlinear clusters. The data points can arrive in any random order and the number of clusters does not need to be specified. Each point is compared against already discovered clusters and each time clusters are joined or divided using an iteratively updated threshold.

1 INTRODUCTION

Rapid increase of Cloud Systems and Internet of Things concept have started to generate a huge amount of data as streams. The processing of stream data requires different approaches as data mining algorithms that require the whole data set for processing are having problems in terms of memory and speed. The data that is already collected and called Big Data can also be considered as a stream since reading and processing Big Data in one step is not possible because of memory constraints. One of the important data mining algorithms is clustering the data into several similar groups. There are several challenges that needs to be addressed here. One of the problems is the unknown number of clusters in the data set where many clustering algorithms require this information before processing. Another problem is the arriving order of the data. Each data point from any cluster can arrive in any order. The evolving nature of some real-time generated data features also increases the difficulty of clustering such streams.

During the analysis no special processing is done for categorical data. One of the basic assumptions in this paper is that the total data set contains groups that can be clustered offline; even theoretically; by a clustering algorithm. The motivation of this assumption is that although the data arrives as a stream one at a time, the total set should contain clusters otherwise no algorithm; streaming or not; can cluster the data.

In this paper we propose an algorithm that requires one pass over the previous data to cluster a stream of data where total number of clusters are not known and not needed. To separate clusters properly, we need a distance metric between clusters. The Cluster Evaluation Function (CEF) developed by the author (Gokcay, 2002) is chosen as a metric since CEF can cluster nonlinear data sets. Many distance measures are using hyper spherical distances and cannot cluster nonconvex data sets. The algorithm is called CEFStream and it can analyze data streams containing nonlinearly separated clusters.

The order of data points are considered as random. The order of arrival should not change the results. Each arriving data point is tested against the clusters created so far. If the distance is closer than a threshold then the point is joined to that cluster. If not, a different cluster is created. After each addition, the cluster distances are measured again to find out cluster groups that need to be merged together.

The paper is organized as follows. Section 2 gives information about previous work. Section 3
introduces the CEFStream Clustering algorithm. Section 4 presents simulation of the algorithm using synthetic data. Finally Section 5 presents future work and Section 6 presents conclusion.

2 RELATED WORK

In this section we briefly review different approaches in clustering data streams.

In the past decades, many stream clustering algorithms have been proposed based on adaption of off-line algorithms, such as BIRCH (Tian, 1997), CluStream (Charu, 2003), DenStream (Feng, 2006), StreamKM++ (Marcel, 2012), etc. These algorithms can be generally summarized into two steps: an online learning step that derives a data abstraction from input data stream, and an offline clustering step that generates the final result. During the online learning step, input data stream is reduced to smaller size representations so that storing the whole data is not needed. Obviously these representations should represent the original distributions properly to be effective. Different algorithms gave different names to these abstract representations like “micro-clusters” in DenStream, or “coreset” in StreamKM++. After receiving new data, there is an online step which determines micro-clusters and an offline step that uses abstract data to finalize clustering. In this step, K-Means (James, 1967) is used by BIRCH, CluStream and StreamKM++, while DBSCAN (Martin, 1996) is used by DenStream.

Although many stream clustering algorithms adopt K-Means or DBSCAN in their offline clustering step. However, both methods have their disadvantages. K-Means needs to know the number of clusters k, and tends to generate similar sized spherical clusters. DBSCAN and similar methods (Xu 2016; Lin 2009; Alazeem 2017) can detect a suitable number of arbitrary shaped clusters, but it has parameters to be decided, which is critical for the result but difficult to decide.

A good review of stream clustering methods is given in (Reddy, 2017).

3 STREAM CLUSTERING USING CEF

Many clustering algorithms are using hyper ellipsoidal distance functions resulting in clusters of similar shape. One exception is density based algorithms (Xu 2017; Chen 2016; Khan 2016; Hassani 2016) where there is no assumption about the shape of the clusters. The problem with density is how to decide the density value so that points are considered to be part of a cluster. In this paper the Clustering Evaluation Function (CEF) developed by the author using Information Theory is chosen. The derivation is described in (Gokcay, 2002) and will not be repeated here. The distance measured by CEF can differentiate nonlinear separated clusters and therefore it is suitable to cluster data with highly nonconvex regions. The CEF function defined for two clusters is given in (1).

\[
CEF(C_p, C_q, \sigma) = \frac{1}{N_p N_q} \sum_{i=1}^{N_p} \sum_{j=1}^{N_q} G(x_i, x_j, 2\sigma^2)
\]

In (1), \(C_p\) and \(C_q\) are clusters of size \(N_p\) and \(N_q\) respectively where \(x_i \in C_p\) and \(x_j \in C_q\). The Gaussian kernel (G) needs a parameter \(\sigma\) for the kernel size which should be determined using the scale of the data.

3.1 Data Stream

The algorithm assumes that each data point \(x_m\) arrives in a random order where \(m \in \text{randperm}(N)\) and \(N\) is the total number of points in the stream. Obviously \(N\) is not restricted in a real data stream but this will not change the assumption about random arrival. Another assumption is that the data has \(d\) dimensions and it is scaled between \([-1 .. 1]\). Although this seems as a strong assumption, it would be easy to scale the data online using a pre calculated scaling factor if the data has known maximum and minimum limits. The time stamp is not important in the analysis.

3.2 Data Structures

During processing we need to create several data structures besides the data itself. For each generated cluster \(C_m\) in the dataset, we have to store number of points in the cluster, an array holding the CEF distance to other clusters and data points that belong to the cluster. The need to store the data stream will increase the processing time since although there is no requirement to copy all data to memory, disk access time will increase total access and processing time. As a future study to reduce the storage and processing requirements, we can replace these points with data skeleton centers (Gokcay, 2016).
The combined storage complexity of the data stream is $O(N)$ which is proportional to the number of points in the stream. The extra information that needs to be stored has a complexity of $O(M)$ and the distance calculation has a storage complexity of $O(M^2)$ where $M$ is the number of clusters discovered so far. We also assume $M \ll N$.

### 3.3 CEFStream Algorithm

The CEFStream algorithm will use one pass over the currently stored dataset. There is no online and offline parts of the algorithm.

For each newly arrived sample point $x$, the distance to already formed clusters is measured once as in (2).

$$D_q(x) = CEF(x, C_q) \quad p \in [1..M]$$

(2)

The next step is to find the closest cluster as in (3). CEF function is inversely proportional to the distance between clusters. Therefore we have to maximize the function.

$$q = \arg \max \{ D_q(x), p \in [1..M] \}$$

(3)

The next step is to check if the closest cluster is larger than a threshold $T$. If $D_q(x) > T$ then the new sample point is added to $C_q$. Otherwise a new cluster is created.

After processing the sample point, the data structure of each cluster needs to be updated. If a new cluster is generated, $M$ will be increased by one for each cluster data structure. The size of CEF distance array is also increased by one to include the new cluster. The distance from the new cluster to other clusters are added to the array. Since these distances are already obtained in the previous step, there is no need to recalculate CEF distances again.

When a new sample point is added to an existing cluster, $M$ is not changed but the distances between clusters still need to be updated. Since the distance from the new point to other clusters are already obtained, we only need one step operation to update each distance.

Once the new sample is processed, it is time to check if there are clusters that need to be merged together using the threshold $T$. Using the updated distance array in each cluster, we will check if there are clusters closer to each other than the threshold $T$ as in (4). The cluster joining operation will continue until there are no clusters left to be joined.

$$\text{join}(C_p, C_q) \text{ if } CEF(C_p, C_q) > T$$

(4)

### 3.4 Threshold Estimation

The threshold $T$ used in the calculation to differentiate clusters from each other can be estimated using an iterative calculation by averaging all distances obtained. Since most of the points are close to each other to form a cluster, the average will give a nice estimation to decide if a point can be considered part of a cluster or not.

### 3.5 Processing Complexity

The distance calculation using the new point is done only once using the available dataset. Therefore the algorithm can be considered as a one-pass algorithm and the complexity is $O(N)$ where the total number of points stored so far is $N$. When there is a new cluster, updating the existing array and adding the distance of the new cluster to other clusters has a complexity of $O(M)$. When an existing cluster needs to be updated by adding the sample point, we have to recalculate all distances. This calculation needs one iteration over all clusters and it should be repeated for every cluster, hence it requires a complexity of $O(M^2)$. Although the CEF distance calculation between clusters $C_p$ and $C_q$ requires a complexity of $O(N_p \times N_q)$, the complexity can be reduced to $O(1)$ by using the already calculated distances and the distance calculation of new point.

Assume that we have a new point $x$ and we already calculated the distances $D_x(p), p \in [1..M]$ from the point to all clusters during cluster assignment. Assume that the new point $x$ belongs to cluster $C_q$ and we have to update all distances from $C_q$ to all other clusters. For example assume that we want to update the distance $CEF(C_q, C_k)$ from cluster $q$ to cluster $k$. When we update cluster $q$ by adding $x$, the new distance is given as $CEF(C_{q+x}, C_k)$.

Cluster $C_{q+x}$ has $N_q + 1$ points and cluster $C_q$ has $N_k$ points. The computational complexity of CEF is $O(N_q \times N_k)$. The calculation can be simplified using...
the previously calculated distance iteratively. The iteration is shown in (5).

\[
CEF(C_{q+x}, C_k) = \left( (N_q \ast N_k \ast CEF(C_q, C_k)) + D_q(k) \right) / ((N_q + 1) \ast N_k) \tag{5}
\]

In the calculation \( CEF(C_q, C_k) \) is already calculated in the previous iteration and \( D_q(k) = CEF(x, C_k) \) is calculated during one pass cluster assignment operation. So the complexity of the calculation is \( O(1) \).

The final complexity is \( O(N) \ast O(M^2) \) where \( N \) is the # of points in the stream and \( M \) is the number of clusters in the stream. Although \( N \) is not limited, \( O(N) \) means that the algorithm is a one-pass algorithm. Also we assume \( M \ll N \) and not chance drastically (will not increase during the flow of the stream), so that practically \( O(M^2) \) can be considered as constant.

### 3.6 Algorithm

The CEFStream algorithm is given below. As long as there is a new sample, the calculation continues as indicated by the while loop.

**Algorithm 1: CEFStream algorithm.**

Input : Data stream \( X \)
Output : Cluster Data Structure

\[
\text{while } x_{\text{rand}} \in X_{\text{STREAM}} \text{ do}
\]

for each Cluster \( C_k, k \in [1..M] \)

\[
D[k] = CEF(x, C_k)
\]

update \( T \)

end

\[
q = \max arg C_k, k \in [1..M]
\]

if \( (D(q) > T) \)

Join \( C_k \) and \( x \)

else

create a new cluster using \( x \)

end

for each Cluster \( C_k, k \in [1..M] \)

for each Cluster \( C_q, k \in [1..M] \)

update \( C_k.D(C_q) \)

end

end

for each Cluster \( C_k, k \in [1..M] \)

for each Cluster \( C_q, k \in [1..M] \)

if \( C_k.D(C_q) > T \)

join \( C_k \) and \( C_q \)

end

end

end

### 3.7 Outliers

There is no special processing for outliers like in (Thakran, 2012). Outliers can be detected by testing the number of points in each cluster. When the number is below a certain threshold we can mark that cluster as an outlier.

### 3.8 Evolving Data Streams

No special processing is used to track changing clusters as the algorithm can track the clusters as long as the change is not drastic and sudden. But if the change of a cluster starts overlapping with other clusters, the algorithm may combine these overlapping clusters. To overcome this problem a window can be applied to the data stream to slowly discard old data.

### 4 EXPERIMENTS

The algorithm is tested using several synthetic datasets. Each time the data arrival is randomly modified to test whether the clusters depend on the arrival order or not. Using a simple dataset given in Fig. 1 the dependency to the order is tested.

![Figure 1: Sample data set.](image)

There are 10 points in the data set and each point is taken from the set using a random permutation simulating a stream. The result after each random generation is the same with correct clustering as in Fig. 2.

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The number of clusters generated during iterations can be seen in Fig. 3. During the start of the computation number of clusters increase, which is expected as the sample order is random. After processing several samples the clusters start to merge and stays the same.

The algorithm is tested using several times where each time the arrival order is changed using a random permutation as given in Table 1. The precision is 100% each time which means that the result matches with the original labels of the data.

Table 1: Different arrival orders of 10 sample points using random permutations.

| Arrangement | 15967810423 | 49853267110 | 41921078635 | 65278910134 | 84911037526 | 86179421035 |

4.1 Test Results

The CEFStream algorithm is tested using several different datasets and each resulting cluster is displayed using a different color. The change of number of groups is also given in Fig. 4. Each set is tested several times using random permutations to change the arriving order of points.

The algorithm is also tested with outliers where each outlier is assigned to a different cluster which will be very easy to eliminate by checking the cluster size. The result is given in Fig. 5.
4.2 Difficult Datasets

There are cases where points from two different clusters start forming a single cluster during the early phase of the iterations. For example occasionally the dataset given in Fig. 6 generates clusters incorrectly depending on the order of arrival.

The original algorithm can merge similar clusters but what we need is to separate a cluster into 2 clusters if the addition of a new point start forming a different sub-cluster. The extra step will come with additional computational complexity.

After adding the new point to a cluster, that cluster will be tested if there is a need to separate the clusters into two different clusters. For this operation we will use the OptimalNumberofClusters algorithm developed (Gokcay, 2017) to find the number of clusters in any dataset. The derivation of the algorithm will not be repeated here. The motivation is not to determine the number of clusters but to test the minimum point of the distance plot created by the algorithm against the threshold T and separate the clusters if necessary. Assuming that the points are arriving one at a time, the running complexity of this algorithm is $O(N_q)$ where $q$ is the current cluster.

5 FUTURE WORK

As a position paper there is work that needs to be completed. The threshold calculation needs to be improved because in some cases the average calculation may not be enough to detect the boundary between clusters. The other improvement will be in the incremental version of the data skeleton algorithm to reduce the storage requirements. Also the algorithm needs to be tested using real data sets as well. Although the algorithm performs well against random arrivals with nonlinearly separated synthetic clusters, the case will be different with more complicated cluster shapes.

6 CONCLUSIONS

In this paper we have developed a one-pass stream clustering algorithm where the clusters are independent of the arrival order and highly nonconvex cluster distributions pose no problem. The distance measure used in the algorithm can cluster nonlinearly separable clusters efficiently. This is not the case with K-means and all its derivatives since the distance measure is hyper-ellipsoidal. No assumption is needed about the possible number of clusters where many algorithms require this number. Each new sample point is processed once and a snapshot can be taken from the algorithm at any time since there are no on-line and off-line iterations.

REFERENCES


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