

Performance Evaluation of Meteor Key Distribution

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Abstract: The Wireless Key Distribution is a fast growing area of applied cryptography covering different approaches of secure distribution of secret encryption key based on channel stochastic properties in specific radio communication systems. To be applicable in key distribution purposes the radio channel must meet the three basic requirements: randomness, reciprocity and spatial selectivity. For a long time it was believed that only the fading channels in multipath environment could satisfy all the three requirements. However, several studies also considered a meteor burst channel as a potential candidate for the secure key distribution at distances up to 2000 km. Unfortunately, a rigorous analysis of meteor radio propagation processes shows that the channel has only approximate reciprocity. This may result in the two legitimate nodes would not be able to generate identical copies of the shared secret key indicating that the Meteor Key Distribution is impossible in practice. In addition, a complicated astronomical nature of the meteor burst channel causes deep diurnal variation in its performance. The aim of our research was a comprehensive evaluation of potential performance of the Meteor Key Distribution systems, taking into account all the physical layer effects of meteor burst channel. We also wanted to clarify, how strong the imperfection of a real meteor burst channel affects the key distribution feasibility.

1 INTRODUCTION

The problem of secure key distribution plays an important role in modern cryptography. In recent decades there has been a shift from mathematical algorithms of key distribution to the principles of creating a shared physical source of randomness. The Quantum Key Distribution (Bennett, 1984) along with the Wireless Key Distribution (Hershey, 1995) systems should be mentioned as notable examples of such approach. As for the Wireless Key Distribution (WKD), it covers different techniques of secure distribution of a shared secret encryption key based on channel stochastic properties in specific radio communication systems. In essence, the pair of legitimate nodes uses their radio channel as a source of common randomness to establish a shared encryption key.

For the radio channel to be applicable in key distribution purposes, it must meet the three basic requirements: 1) randomness, 2) reciprocity and 3) spatial selectivity. The channel randomness provides high entropy and unpredictability of the shared key. The channel reciprocity ensures that both legitimate

users to create identical copies of the key. The spatial selectivity, also referred as a sharp spatial decorrelation of radio signal parameters, prevents a key leakage to potential eavesdropper.

In (Sidorov, 2007) a meteor burst channel (MBC) has been proposed as a candidate for the secure key distribution over the distances up to 2000 kilometers. We call this method as the "Meteor Key Distribution" (MKD). The aim of our research was a comprehensive evaluation of potential performance of the Meteor Key Distribution systems, taking into account all the physical layer effects of meteor burst channel. First of all, we wanted to clarify, whether the imperfection of a real meteor burst channel violates the key distribution feasibility and if not, then how strong it decreases achievable MKD performance.

In (Sulimov and Karpov, 2014) a randomness of meteor radio propagation has been justified. In particular, an unpredictability of the radio propagation path has been proved. Such path unpredictability property ensures a carrier phase of received signal to be a random variable. A pair of legitimate nodes is able to create two copies of the

shared key simply by measuring values of the carrier phase at receiving a meteor radio reflection.

However, only one of the three WKD basic requirements has been justified for the Meteor Key Distribution to this moment. Moreover, there is a serious problem in MKD justifying at the second stage, where the channel reciprocity must be required. A rigorous analysis of meteor radio propagation shows that the meteor burst channel provides only approximate (non-perfect) reciprocity (Desourdis, 1993). The MBC reciprocity violation (or simply non-reciprocity) observed in practice causes some mismatch in copies of the shared secret key generated by the legitimate parties. This circumstance substantially decreases a key generation rate. That is the issue we are going to address in the following sections.

2 SIMULATION SCENARIO

An implementation of field experiments on channel non-reciprocity at real meteor radio link is extremely costly. This makes computer simulation method quite reasonable to study the MKD processes. As far as we know, KAMET (Karpov, 2001) is the only MBC computer model that simulates all basic effects of the channel non-reciprocity. This model uses an improved electrodynamic calculations unit based on rigorous diffraction theory of scattering of radio waves off the ionized meteor trails (Khuzyashev, 1984). Another advantage of the KAMET model is a high accuracy of the performance simulation of meteor burst communication systems. Such accuracy is achieved through accounting within simulation a long-term statistics of radar observations on the influx of meteoroids into the Earth's atmosphere.

Table 1: Technical specifications for the test meteor links.

Specification	Link 1	Link 2
Localization of the legitimate users	Paris-Colmar	Vienna-Colmar
Link length	379 km	669 km
Carrier frequency	50 MHz	
Transmitted power	2000 Watts	
Required signal-to-noise ratio	20 dB	
Amplitude threshold level	-6 dB μ	
Date of key distribution session	July 21, 2015	

Two test meteor links have been simulated to perform evaluation of the key generation rate. Table 1 summarizes their technical specifications. We intentionally used the test links of different lengths to show that the link geometry and especially the

link length affect crucially the performance of any type of meteor burst system including MKD. While the Link1 could be assigned as a short link, the Link2 has a quite moderate length. Besides, due to MBC complex astronomical nature its capacity varies greatly dependent on the session date and season. Without loss of generality concerning about MKD feasibility, we chose session date at July 21st.

The following scenario was implied during the MKD simulation. A pair of legitimate users (say, Alice and Bob) exchange with a series of sounding signals through the meteor burst channel in a full duplex synchronized mode. When an ionized trail left by burning in the upper atmosphere meteoric particle arises at the altitudes of 80-110 km, the sounding signals transmitted simultaneously in opposite directions by Alice and Bob are reflected from it to be received at another side of the link. At receiving the signals Alice and Bob measure their carrier phase. Due to the MBC stochastic properties all these measurements are random numbers. By repeating such exchange with Bob for N times, Alice collects the $\{\varphi_A\}_N$ sample of the carrier phase measurements. Bob collects the $\{\varphi_B\}_N$ sample of size N in the same manner. Further, Alice and Bob achieve an agreement on the parameters of the so-called "bit extraction" procedure (Croft, 2011) to extract a shared key from their samples. After that, Alice processes her own sample $\{\varphi_A\}_N$ to generate the K_A key, and Bob generates the K_B key by extracting the random bits from his sample $\{\varphi_B\}_N$ in the same way as Alice. From each single measurement of the carrier phase Alice and Bob were extracting m random bits, where m is the so-called "codeword length".

In practice, due to the non-perfect channel reciprocity and noise factors, the $\{\varphi_A\}_N$ and $\{\varphi_B\}_N$ samples always have some mismatch $\Delta\varphi = |\varphi_A - \varphi_B|$. As a result, the K_A and K_B instances of the secret key also contain some bits in mismatch. To eliminate them, Alice and Bob had to perform a key reconciliation procedure. Unfortunately the key reconciliation leads to some loss in the length of the final shared key. If we denote its effectiveness as η , then we can state that $\eta < 100\%$. According to (Smolyakov, 2013), the key generation rate R_K can be estimated by the formula (1), where T is the sample collecting time.

$$R_K = m \cdot N \cdot \eta / T \quad (1)$$

The two test meteor links (Paris-Colmar and Vienna-Colmar, respectively) we simulated had essentially different lengths. As will be shown in the

following sections, this difference resulted in very different MKD performance indication.

3 NON-RECIPROcity OF METEOR BURST CHANNEL

The channel reciprocity is one of the basic properties allowing secret key distribution between a given pair of legitimate nodes. Unfortunately, an absolute reciprocity is only an idealization of observed reality. Specifically, the fact that the meteor burst channel has non-perfect reciprocity inevitably leads to some key mismatch between Alice and Bob. This limits achievable key generation rate.

There are three main physical layer mechanisms that cause the channel non-reciprocity:

- 1) polarization phenomena at meteor radio propagation (Khuzyashev, 1984);
- 2) multipath fading effects due to scattering from meteor trails with several reflecting points (Weitzen, 1987);
- 3) random drift of the trail reflecting point due to ionospheric turbulent winds (Weitzen, 1987).

All these mechanisms are included into the KAMET simulation model allowing them to be accounted within our MKD performance evaluation.

The polarization phenomena are the dominant factor of the MBC non-reciprocity. This follows from the fact that the characteristics of scattered signal depend mainly on the polarization of the incident radio wave. Looking deeper, we find that two radio waves transmitted by Alice and Bob in the opposite directions experience different polarization rotations during their propagation. This is due to non-symmetric geometry of the uplink propagation accompanied with the influence of the Faraday Effect in a magnetized ionosphere. As a result, the incident waves transmitted by Alice and Bob have different polarizations. This causes non-symmetric scattering of opposite signals off the meteor trail and in the end it results in the channel non-reciprocity.

We also took into account effects of the meteor trail decay with continuous changing in time of its scattering properties. Therefore, level of channel non-reciprocity constantly varies during the signal detection time. For the key distribution purposes, Alice and Bob should choose the time point that corresponds to the lowest value of phase non-reciprocity (Sulimov, 2014). By accounting all the above mentioned non-reciprocity effects, we were able to properly simulate limitations on the key

distribution process caused by physical layer of channel. Let us consider the simulation results next.

Figure 1 shows a probability distribution histogram of the phase mismatch $\Delta\varphi = |\varphi_A - \varphi_B|$ simulated for the test Link 1 (Paris-Colmar). As can be seen, the phase mismatch $\Delta\varphi$ typically shows a symmetric shape of probability distribution with a zero mean value. The histogram standard deviation was $\sigma(\Delta\varphi) = 16.1^\circ$, and we should admit here that this is a relatively high value with non-negligible impact. The test Link 2 (Vienna-Colmar) showed a similar histogram but with slightly lower standard deviation: $\sigma(\Delta\varphi) = 14.2^\circ$. It should be noted that even such a small decrease in the standard deviation $\sigma(\Delta\varphi)$ allowed 24% higher key reconciliation efficiency η than achieved at the Paris-Colmar link.

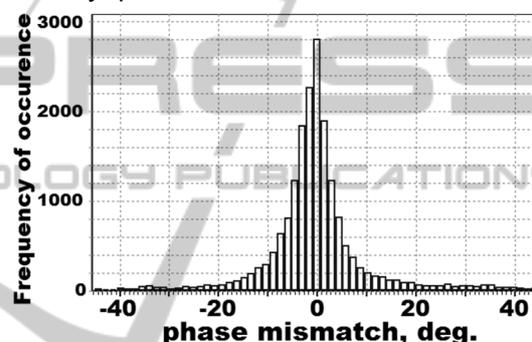


Figure 1: Probability distribution histogram of mismatch in the phase measurements of legitimate users observed at the Paris-Colmar meteor link (session time is 3 a.m.).

The simulation results also revealed that a certain amount η_{NR} of detected meteor trails had an extremely high non-reciprocity. Obviously, such "extremely" non-reciprocal trails could not be used for the key distribution purposes. About $\eta_{NR} = 3.2\%$ of all the trails detected at the Paris-Colmar link were recognized as extremely non-reciprocal. The same indicator for the Vienna-Colmar link showed the value about $\eta_{NR} = 1.2\%$.

As will be shown in the Section 5, the presence of phase mismatch significantly limits achievable key generation rate. However, a meteor activity during the key distribution session is another essential factor that defines the key generation rate. The next section presents our simulation results of the meteor activity for both test meteor links.

4 DIURNAL VARIATIONS OF METEOR ACTIVITY

The meteor activity and performance of meteor burst

channel experience a strong diurnal variation (McKinley, 1961). According to formula (1), the key generation rate should follow these variations too. The diurnal dependence is a legacy of the MBC astronomical nature. In practice, current meteor activity is usually characterized by the number of meteor trails N_H detected per 1 hour (the hourly number of meteors). In this case, the (N/T) multiplier in the formula (1) should be substituted by $(1 - \eta_{NR}) \cdot (N_H / 3600)$. The KAMET simulation model we used is based upon a long-term (about several decades) statistics of radar observations on the meteoric matter influx. This allows accurate prediction of diurnal variations in the meteor activity for a given radio link. Let us consider the simulation results obtained for the test meteor links.

Fig. 2 presents the diurnal variations in the meteor activity simulated for both test links. As all the legitimate nodes had close geographic latitudes, the curves reveal very similar profiles but differ greatly in the number of detected meteor trails. The maximum meteor activity with hourly number $N_H = 280$ at the Paris-Colmar link was observed at 3 a.m. local time. The minimum meteor activity with $N_H = 63$ was observed at 12 a.m. Similarly, at the Vienna-Colmar link the maximum meteor activity with $N_H = 598$ was observed at generating the shared key at 2 a.m. local time. The minimum activity with hourly number $N_H = 181$ was observed at 3 p.m.

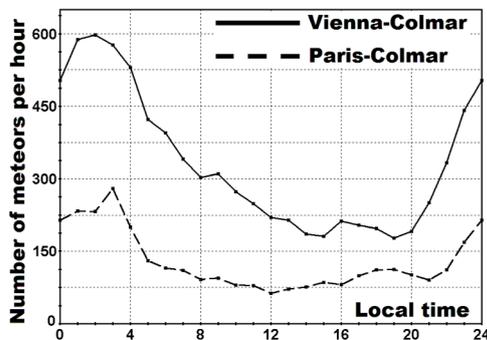


Figure 2: Diurnal variations in meteor activity simulated for the test meteor links.

As we can see, the depth of diurnal variations in meteor activity may reach up to 400%. Another feature of meteor burst communications is a low performance at shorter radio links. Moreover, an accurate analysis shows that the MBC non-reciprocity reveals some daily changes too. However, these changes affect the MKD performance much weaker than the variations in meteor activity. To be precise, it's about 15% versus 400% in depth. Thus, we should seek a maximum

possible hourly number N_H in the first place in practice.

From the practical point of view, the greatest interest is to estimate a gap between the maximum and minimum MKD performance indications. Considering this, for each test link we present the performance evaluation corresponding only to the maximum and minimum values of meteor activity.

5 EVALUATION OF KEY GENERATION RATE

After collecting the $\{\phi_A\}_N$ and $\{\phi_B\}_N$ samples of carrier phase measurements for all detected meteor radio reflections Alice and Bob were able to generate the K_A and K_B copies of the secret key, respectively. To generate them, Alice and Bob implemented a so-called "bit extraction" procedure following the algorithm presented in (Sulimov and Smolyakov, 2014). Since (Sulimov and Karpov, 2014) proved an absence of autocorrelation within the sample of carrier phase measurements, no sample decorrelation procedure was required to create the key.

As long as meteor burst channel always causes some mismatch in phase measurements of Alice and Bob, a key reconciliation procedure is required. We used the cyclic redundancy codes (CRC-16) for the K_A and K_B reconciliation. This is similar to a more secure approach based on the privacy amplification (Bennet, 1995) but easier to implement. In this case, the choice of optimal codeword length m^* for the bit extraction procedure becomes a crucial moment. Increasing the value of m we can extract more random bits from each single carrier phase measurement but causing a higher key mismatch rate p_e . It's a trap, because high key mismatch rate may result in total rejection of the key at its reconciliation.

Fig. 3 presents the key mismatch rate (i.e. the probability of mismatch in a single bit) as a function of the codeword length m used in the bit extraction procedure. We can see an asymptotic rise of the key mismatch rate up to the $p_e = 50\%$ value as the codeword length m is increased. Due to the channel non-reciprocity and noise factors, even at the smallest possible value of $m = 1$ bit per measure a very high key mismatch rate is observed: $p_e = 4.7\%$ (for the Paris-Colmar link) and $p_e = 3.5\%$ (for the Vienna-Colmar link), respectively. At first glance, the obtained values of the key mismatch rate are unacceptably high, but such mismatch levels are

quite typical for the WKD practice. We can refer to the (Liu, 2014) research as an example of the recent experimental results.

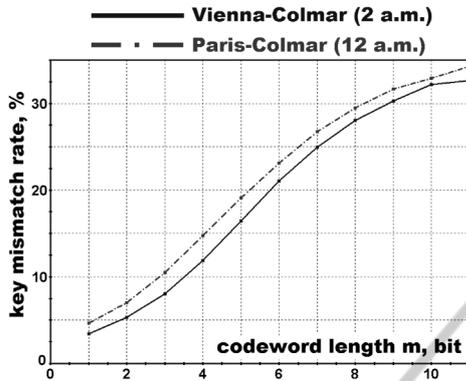


Figure 3: Key mismatch rate as a function of the codeword length.

During the key reconciliation each block of l_b key sequence bits in length had been verified with the CRC-16 code. After successful block verification we excluded 16 arbitrary bits from this block to prevent a key leakage caused by the CRC-code transmission. For each value of the codeword length m we were looking for an optimal verification block length l_b^* that provided a maximum efficiency of the key reconciliation η^* . This was done, because a short verification length l_b leads to a strong loss in the verified key size after excluding 16 arbitrary bits from the reconciled block. At the same time, the longer the block size l_b the higher probability of the mismatch and reconciliation failure.

Table 2: Optimization of the codeword length (Link2 at 2 a.m.).

m , bit	$corr(K_A, K_B)$	p_e , %	l_b^* , bit	$\eta(l_b^*)$, %	R_K , bph
1	0,931	3,45	33	15,46	91,4
2	0,893	5,34	32	13,05	154,3
3	0,840	8,02	30	8,88	157,4
4	0,763	11,87	24	4,50	106,3
5	0,671	16,43	23	1,82	53,7
6	0,579	21,06	22	0,763	27,1
7	0,501	24,95	21	0,295	12,2
8	0,439	28,03	20	0,114	5,4
9	0,394	30,31	20	0,047	2,5
10	0,356	32,19	20	0,036	2,1
11	0,344	32,64	20	n/a	n/a

Table 2 shows an example of the codeword length optimization for the Vienna-Colmar test link (session time 2 a.m.). The second column presents the cross-correlation coefficient $corr(K_A, K_B)$ of the K_A and K_B key instances, which is a linear function of the key mismatch rate p_e . The maximum key

generation rate (in bits per hour) was achieved at $m^* = 3$ bits per measure. The table string with corresponding performance indicators is highlighted in bold.

Resulting estimates of the key generation rate R_K based on the formula (1) are presented in Fig. 4 (test Link 1) and Fig. 5 (test Link 2), respectively. For each test link we present the codeword optimization curves at the minimum and maximum levels of meteor activity. Due to a high mismatch in the phase measurements of Alice and Bob the optimal codeword length m^* had very low values: $m^* = 2$ or $m^* = 3$. Low values of meteor activity resulted in a very humble MKD performance with maximum key generation rate R_K less than 160 bits per hour.

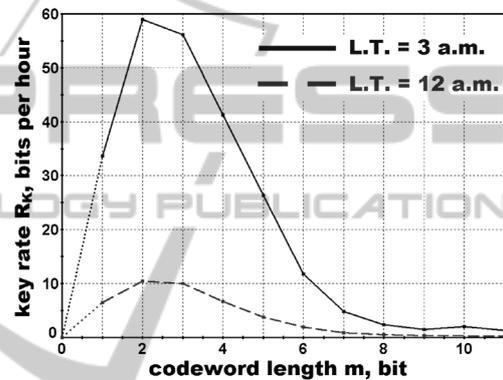


Figure 4: Key generation rate as a function of the codeword length (Paris-Colmar meteor link).

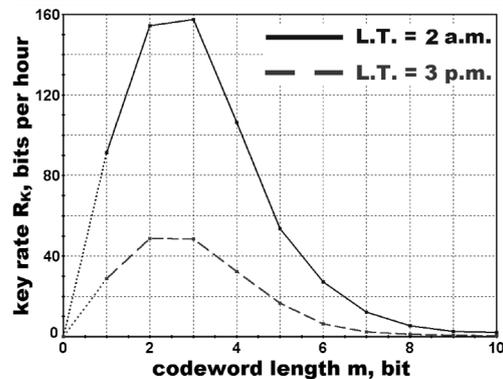


Figure 5: Key generation rate as a function of the codeword length (Vienna-Colmar meteor link).

Taking into account all the diurnal variations in the MKD performance observed at the Paris-Colmar link, Alice and Bob were able to generate a shared key with total length of 571 bits for a 24 hour period. Thanks to a higher meteor activity, similar estimates performed for the Vienna-Colmar link show the key length about 2081 bits per 24 hours day.

Summarizing our results, we can conclude that, despite non-perfect reciprocity of meteor burst channel, the Meteor Key Distribution is feasible in practice. Physical parameters of each individual meteor link essentially affect the performance of the key distribution process making it unique for the given radio link. Despite the humble key generation rate, we believe that MKD systems could be useful in practice. Their advantages are much greater key distribution distances (up to 2000 km), low cost equipment, vitality and ability to operate under conditions of severe climate such as in polar and other remote regions.

6 CONCLUSIONS

Our study presents the very first attempt to evaluate a performance of Meteor Key Distribution systems with an account of full complex of the physical layer effects of meteor burst channel (MBC). To deal with the channel non-reciprocity and diurnal variations effects, we used a MBC simulation model with improved electrodynamic and astronomical calculations unit.

Our simulation results proved that, despite non-perfect reciprocity of meteor burst channel, the Meteor Key Distribution is feasible in practice. Physical parameters of each individual meteor link essentially affect the performance of the key distribution process making it unique for the given radio link. The main factor limiting key generation rate is low meteor activity and its diurnal variation. As show presented estimates, a shared key of 571 bits in length could be generated for a 24 hour period at the Paris-Colmar meteor link with a peak key generation rate about 59 bits per hour. Similar estimates performed for the Vienna-Colmar meteor link show 2081 bits as a daily key length with a peak rate about 157 bits per hour, respectively.

Despite the humble key generation rates presented, we believe that MKD systems could be useful in practice. Specifically, they might be used for key distribution at up to 2000 km distances in such applications as AES, lightweight cryptography like CLEFIA, and hash-based post-quantum subkeys.

As we stated in the beginning, for the radio channel to be applicable in key distribution purposes, it must meet the three basic requirements: 1) randomness, 2) reciprocity and 3) spatial selectivity. Hence, the spatial selectivity of meteor burst channel must be addressed at the next stage of justification of the Meteor Key Distribution systems.

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