

Mapping Textual Scenarios to Analyzable Petri-Net Models

Edgar Sarmiento¹, Eduardo Almentero², Julio C. S. P. Leite¹ and Guina Sotomayor³

¹*Departamento de Informática, PUC-Rio, Rio de Janeiro, Brazil*

²*Departamento de Matemática – DEMAT, Universidade Federal Rural do Rio de Janeiro - UFRRJ, Rio de Janeiro, Brazil*

³*Instituto Nacional de Matemática Pura e Aplicada - IMPA, Rio de Janeiro, Brazil*

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Abstract: With the growing use of user-oriented perspectives at requirements engineering, transforming requirements models into executable models is considered to be significant. One of the key elements in this perspective is the notion of scenarios; scenarios are used to describe specific behaviors of the application through a flow of events based on user-perspective. Since scenarios are often stated in natural languages, they have the advantage to be easy to adopt, but the requirements can then hardly be processed for further purposes like analysis or test generation; partly because interactions among scenarios are rarely represented explicitly. In this work, we propose a transformation method that takes textual description of scenarios as input and generates an equivalent Petri-Net model as output. The resulting Petri-Net model can be further processed and analyzed using Petri-Net tools to verify model properties, to identify concurrency problems and to optimize the input and output models. Demonstration of the feasibility of the proposed method is based on two examples using a supporting tool.

1 INTRODUCTION

Scenario-based representations are used in Requirements Engineering mainly because it improves communication among clients and developers. In this context, requirements are stated as a collection of scenarios and described by specific flows of events in the system. The use of scenarios helps in understanding a specific situation in an application, prioritizing their behavior. The most prominent languages to describe scenarios are restricted-form of use case descriptions (Cockburn, 2001; Gutiérrez et al., 2008), UML dynamic behavior diagrams and Message Sequence Charts (Andersson and Bergstrand, 1995).

The graphical notation based languages are very attractive and user-friendly; however, they can be difficult to design, and domain experts cannot reasonably be asked to draw them (Gutiérrez et al., 2008). Although the mentioned languages provide an accessible visualization of models, they lack formal semantics to support the analysis of structural and behavioral properties of the application.

For practical reasons, and in order to allow for an easy communication with stakeholders, requirements are written using natural language-based textual

templates. Textual scenario-based approaches offer several practical advantages: **(1)** Scenarios are easy to describe and understand. **(2)** They are scalable; the behavior of a large and complex system can be stated as a collection of independently and incrementally developed scenarios. **(3)** It is easy to provide requirements traceability throughout the design and implementation (Lee et al., 1998).

Unfortunately, textual scenarios exhibit some shortcomings: **(1)** Scenarios informally specified are usually hard to analyze, because natural language is by definition ambiguous. **(2)** Modularity is poorly supported because the interactions among scenarios are rarely represented explicitly. **(3)** There are currently no systematic approaches to make explicit interactions by concurrency among textual scenarios. Requirements are rarely truly independent, they interact (Lee et al., 1998).

Scenario languages are either informal or semi-formal and cannot be used for further analysis of the application. In order to automatically analyze the requirements it is necessary to translate them from informal languages to formal languages like Petri Nets. Thus before developing, requirements engineer needs to create requirements specification in two formats. One format is to communicate with

customers (textual scenarios) and other format is for analysis or testing (Petri-Nets).

The original contribution of this work is an automated transformation method that takes textual description of scenarios (conform to a metamodel defined in this work) as input and generates an equivalent Petri-Net model (conform to a restricted Petri-Net metamodel) as output. The generation is performed by a model transformation, defined as mapping rules and implemented in the C&L (C&L, 2014) prototype tool. It might eliminate the redundancy of writing specification twice.

This transformation allows to benefit both from graphical and textual scenario advantages, and it allows an easier integration to available Petri-Net tools (PIPE2, 2014). On the basis of this transformation, it is possible: **(1)** Analyze some properties of the executable model, i.e. to verify whether the scenario is consistent, complete and correct. **(2)** Identify concurrency problems such as deadlocks among concurrent scenarios that compete with each other for common resources. **(3)** Use the Petri-Net for further treatments like test generation.

The details of our proposal are presented in 6 Sections, from the related work, the description of the source and target metamodels, the strategy we propose, to the case study and conclusions.

2 RELATED WORK

Many researches have shown the importance to formalize the informal aspects of scenarios in order to be useful in automated analysis. Some researches focused on developing the formal semantics for scenario representations, recent researches focus on developing techniques to transform scenarios into executable models with rigorous semantics.

Hsia et al., (1994) used a BNF-like grammar to formally describe scenarios. Scenarios are represented like scenario trees and scenario schema. A tree is constructed to represent all the scenarios for a particular user view. While scenario trees are defined, each of scenarios is converted into an equivalent regular grammar. This approach is only effective when applied to a small number of relatively simple scenarios (Lee et al., 2001).

UML Sequence Diagrams (Lee et al., 2001) and Message Sequence Charts - MSCs (Andersson and Bergstrand, 1995) are frequently used as formalisms for scenarios. The main problem tackled by these approaches is the interactions among scenarios, and these have advantages over the grammar-based approach in terms of scalability/understandability.

However, these models are either informal or semi-formal and can not be used for automated analysis.

In our work, we describe scenarios using a restricted form of the natural language; then, scenarios are transformed into Petri-Nets, which are used as the mechanism to enable the analysis. Other approaches based on natural language include (Lee et al., 1998; Somé, 2007; Zhao and Duan, 2009).

In (Lee et al., 1998), is proposed a systematic procedure to convert use case descriptions into Constraint-based Modular Petri-net models, and to analyze use cases. To facilitate the transformation, use cases are described in relation to formal definition of *pre and post-conditions* (like Action-Condition tables). Use cases are considered as a collection of interacting and concurrently executing units of functionalities. However, intermediate models are created and alternative/exception flows of use cases are not considered.

In (Somé, 2007), is proposed a semantics for use cases based on Petri-Nets. However, the syntax to describe use cases does not deal with non-sequential relationships (concurrency) and only deals with sequential relationships (include and extend).

In (Zhao and Duan, 2009), is proposed an approach to formalize use cases with Petri-Nets. A semi-formal language is proposed for use case syntax. This syntax is based on message sender and receiver objects, and the events in use cases can be sequential, selection, iteration and concurrent. Petri-Nets are derived extracting objects and messages between objects. However, it is necessary to create “*event frames*” for the extraction of the objects and the message from each one the sentence events.

The related Petri-Net based approaches exhibit the following shortcoming: **(1)** Scenarios are described in relation to formal definition of pre and post-conditions. **(2)** There is a lack of systematic procedures on how to represent the given scenarios. **(3)** The transformation of scenarios to Petri-Nets is not automated (intermediate models). **(4)** Scenarios do not provide constructs to support *modularity*.

On opposite, our approach: **(1)** Use a semi-structured natural language to write scenarios. **(2)** Define an abstract and concrete syntax for scenarios. **(3)** Implement automated mapping rules. **(4)** Provide powerful characteristics to deal with *modularity* and identify concurrency problems.

3 BACKGROUD

The natural language representation of scenarios is based on a previous work where were defined an

abstract syntax (metamodel) and a concrete syntax (restricted form of natural language) for scenarios (Leite et al., 2000). Differently to previous work, in this work our focus is the transformation of scenarios to executable models. For further purposes like analysis or test generation, the syntax was updated: the *result* attribute (or expected result) was added to the scenario syntax.

3.1 Scenario

Scenario is a language used to help the understanding of the requirements of the application; it is easy to understand by the developers and other stakeholders. Scenario represents a partial description of the application behavior that occurs at a given moment in a specific geographical context - a *situation* (Leite et al., 2000).

In this work, the scenario modelling is based on a semi-structured natural language proposed by Leite et al., (2000), and it is composed of the entities described in Table 1.

Use case (Cockburn, 2001) is a particular model of scenario. Use cases describe the interaction between the users and the system through its interface. Scenarios describe: (1) situations in the environment and the system, (2) interactions among objects or modules and (3) procedures or methods. Table 1 explains how a scenario (Leite et al., 2000) can be also used as a use case (Cockburn, 2001).

Table 1: Comparing scenario and use case.

Scenario	Description	Use Case
Title	Identifies the scenario. Must be unique.	Use Case #
Goal	Describe the purpose of the scenario.	Goal In Context
Context	Describes the scenario initial state. Must be described through at least one of these options: pre-condition, geographical or temporal location.	Scope
		Level
Resources	Passive entities used by the scenario to achieve its goal. Resources must appear in at least one of the episodes.	Preconditions
		Trigger
Actors	Active entities directly involved with the situation. Actors must appear in at least one of the episodes.	Actors
Episodes	Sequential sentences in chronological order with the participation of actors and use of resources.	Description
Exception	Situations that prevent the proper course of the scenario. Its treatment should be described.	Extensions
		Sub-Variations
Constraint	Non-functional aspects that qualify/restrict the quality with which the goal is achieved. These aspects are applied to the context, resources or episodes.	
Result	Internal condition satisfied by an episode/exception, and described as a message or information of the state of some resource.	

Figure 1 shows a metamodel for scenario description used in this work. It defines an abstract syntax for a scenario using a class diagram.

Definition 1: According to our metamodel, a *scenario* is a 7-tuple $S = (Title, Goal, Context, Resources, Actors, Episodes \text{ and } Exceptions)$ and the attributes *Constraint* and *Result*.

A *scenario* S must satisfy a *goal* that is reached

by performing its *episodes*. The episodes describe the operational behavior of the situation, which includes the main course of action and possible alternatives. An *exception* can arise during the execution of episodes, and indicates that there is an obstacle to satisfy the goal. The treatment to this exception does not need to satisfy the scenario goal.

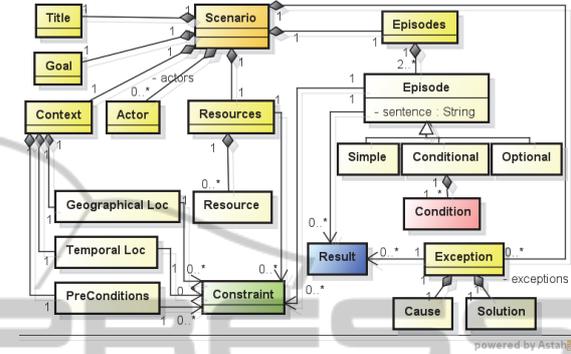


Figure 1: Scenario metamodel.

The *episodes* of a scenario can be of three different types: *simple*, *conditional* and *optional*. Simple episodes are those necessary to complete the scenario; Conditional episodes are those whose occurrence depends on internal or external condition, internal conditions can come from scenario pre-conditions, resources, actors, constraints or previous episodes; Optional episodes are those that may or may not take place depending on conditions that cannot be detailed.

A sequence of episodes implies a precedence order, but a non-sequential order can be bounded by the symbol “#”. This is used to describe parallel or concurrent episodes (#<Episode Series>#).

The scenario language makes explicit the sequential interactions among scenarios. *Scenarios* can be connected to *other scenarios* through links, yielding a complex network of relationships:

Integration Scenario gives an overview of the relationship among several scenarios of the application, since each integration scenario episode corresponds to a scenario.

Sub-scenario is defined when an episode of a scenario can be described by another scenario. This allows the decomposition of complex scenarios, facilitating both its writing and understanding.

Pre-condition is a relationship defined within the context element of a scenario. A scenario that is pre-condition to other must be executed first and so on.

Exception relationship is defined when a scenario is used to detail the exceptional behavior of another. The main scenario should be executed and an

exception must occur to execute the other one.

Constraint relationship is defined when a scenario is used to detail non-functional aspects that prevent the proper execution of another, which also give us an order among the scenarios.

3.2 Linguistic Pattern Syntax

In order to reduce ambiguity in natural language requirements descriptions, we have defined a concrete syntax based on linguistic patterns for describing scenario elements conform to this metamodel (Figure 1).

Table 2 shows the different linguistic patterns (template) for describing scenarios based on natural language. The scenario model should be seen as a syntax and structural guidelines to facilitate the automated analysis (Leite et al., 2000).

In Table 2, + means composition, {x} means 0 or more occurrences of x, () is used for grouping, | stands for or and [x] denotes that x is optional.

Table 2: Linguistic patterns for describing scenarios.

TYPE	DESCRIPTION
Title	Phrase ([Actor Resource] + Verb + Predicate)
Goal	{Actor Resource} + Verb + Predicate
Context	{Geographical Location} + {Temporal Location} + {Pre-condition}
Geographical Loc.	Phrase + {Constraint}
Temporal Location	Phrase + {Constraint}
Pre-condition	{Subject Actor Resource} + Verb + Predicate + {Constraint}
Resources	{Name} + {Constraint}
Actors	{Name}
Episodes	{ (<Sequential Group> <Non-Sequential Group>) }
Sequential Group	<Basic Sentence> <Basic Sentence> <Sequential Group> <Basic Sentence>
Non-Sequential Group	#<Episode Series>#
Episode Series	<Basic Sentence> <Basic Sentence> <Episode Series> <Basic Sentence>
Basic Sentence	<Simple Episode> <Conditional Episode> <Optional Episode>
Simple Episode	<Id> <Episode Sentence> CR
Conditional Episode	<Id> IF <Condition> THEN <Episode Sentence> CR
Optional Episode	<Id> [<Episode Sentence>] CR
Exceptions	{Exception}
Exception	<Id> IF <Cause> THEN <Solution>
Id	Identifier
Episode Sentence	((([Actor Resource] + Verb + Predicate) ([Actor Resource] + [Verb] + Title)) + [{"with the result"} {"such that"}] {Result}) + {Constraint}
Solution	{(Verb + Predicate) Title} + [{"with the result"} {"such that"}] {Result}
Condition/Cause/Result	{[Subject Actor Resource] + [Verb] + Predicate} Phrase
Constraint	{[Subject Actor Resource] + Must [Not] [Verb] + Predicate} Phrase

A **simple episode** is described as follows:

<Id> (([Actor | Resource] + (Verb + Predicate) | ([Actor | Resource] + [Verb] + Title)) + [{"with the result"} | {"such that"}] {Result}) + {Constraint}

An **episode** accesses or modifies **resources** and it is executed by **Actors**. The relevant information for an **episode** is the action performed (**episode sentence**). Optionally, it is possible to add non-functional requirements (**Constraint**) related to the episode and to add expected results (**Result**).

A **result** is not a post-condition, because a post-

condition is a successful response of the system when the main flow of episodes is carried out. A **Result** or **expected result** is an internal condition and it is important in Model-based Testing context.

An **exception** is described as follows:

<Id> IF <Cause> THEN ((Verb + Predicate) | Title) + [{"with the result"} | {"such that"}] {Result}

The first element of an **exception** is the **identifier**. This is composed by the identifier of the **episode** followed by the number of the exception (an episode can throw several exceptions). The second element is the **Cause** that triggers the exception, the third element is the **Solution** to treat the exception, and the **Result** attribute are the expected results at the end of performs the **Solution**.

3.3 Analyzable Petri-Net Model

Petri-Net is a graphical and mathematical modeling and analysis language for describing and studying systems that are characterized as concurrent, asynchronous, distributed, parallel, nondeterministic, and/or stochastic.

A Petri-Net (Figure 2) is composed of nodes that denote places (Place) or transitions (Transition). Nodes are linked together by arcs (Arc).

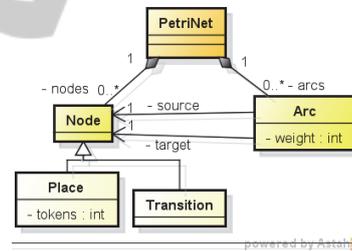


Figure 2: Petri-Net metamodel.

Transitions are active components. They model the activities that can occur – **events**, thus changing the state of the system. Transitions are only allowed to fire if they are enabled, which means that all the pre-conditions for the activity have been fulfilled.

Places are passive components and placeholders for tokens. They model communication medium, buffer, geographical location or a possible state (condition). The current state of the system being modeled is called **marking** which is given by the number of tokens in each place.

Arcs are of two types: Input arcs start from places and ends at transitions, while output arcs start at a transition and end at a place.

When the transition fires, it removes tokens from its input places and adds some at all of its output

places. The number of tokens removed/added depends on the *cardinality (weight)* of each arc.

Definition 2. A *place-transition* Petri-Net is a five-tuple $PN = (P, T, F, W, M_0)$ where $P = \{p_1, p_2, \dots, p_n\}$ is a finite set of places, $T = \{t_1, t_2, \dots, t_m\}$ is a set of transitions, $F \subseteq (P \times T) \cup (T \times P)$ is a set of arcs, $W : F \rightarrow \{1, 2, \dots\}$ is a weight function, $M_0 : P \rightarrow \{0, 1, 2, \dots\}$ is the initial marking and $P \cap T = \emptyset$.

4 MAPPING RULES

This section defines the transformation (Figure 3) to generate a place-transition Petri-Net (instance of the metamodel of Figure 2) from a scenario description (instance of the metamodel of Figure 1). By an automatic transformation, we can have more precise requirements through the analysis of Petri-Nets. Detailed mapping rules are described below.

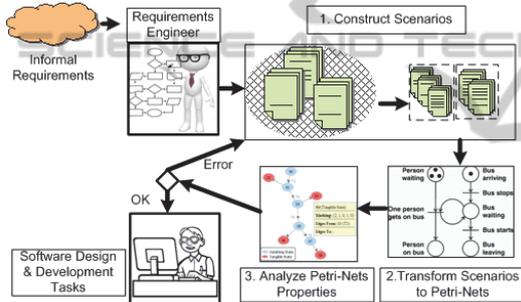


Figure 3: Overview of the proposed method.

4.1 Initialization

The initial state (context) and the resources used by the scenario are mapped into a Sub-Petri-Net composed of *Places* generated from geographical location, temporal location, pre-conditions and constraints as follows:

Initially, one *Place* p with an initial token is generated for the Scenario S :

→ *Place*, p with $p.name = S.title$; $p.tokens = 1$;

For every *pre-condition* pc in *Pre-conditions*:

→ *Place*, p with $p.name = pc.name$; $p.tokens = 1$;

For every *Constraint* c in *Context*:

→ *Place*, p with $p.name = c.name$; $p.tokens = 1$;

For every *Constraint* c in *Resources*:

→ *Place*, p with $p.name = c.name$; $p.tokens = 1$;

Places generated from the context and the resources are “*input places*” of the first episode of the scenario.

4.2 Mapping Episodes

Initially, each one of the *episodes* is mapped into the “*transition*” t ($t.name = episode\ sentence$) and its internal “*dummy places*” (*input dummy place* p_{id} and *output dummy place* p_{od}) of a Sub-Petri-Net. The *Conditions* or *Option* to trigger the *transition* t are mapped into “*input places*” with an initial token.

When episodes are performed, exception can arise. An exception is mapped into the “*transition*” t_{ex} ($t_{ex.name} = solution$) and its internal “*places*” (*output places* for the results). The “*output dummy place*” p_{od} is linked to the *transition* t_{ex} . The “*condition*” or “*cause*” to trigger the *transition* t is mapped into an “*input place*” with an initial token.

A *Constraint* is an “*input place*” (non-functional requirement, resource and also time constraints) that are needed in order to perform the *transition* t .

A *result* is an “*output place*” satisfied by an internal condition of the *transition* t .

Places and *Transitions* are generated as follows:

For every Episode e in Episodes:

→ *Place*, p_{id} with $p_{id.name} = "IDummy_" + e.id$;

→ *Transition*, t with $t.name = e.episode_sentence$;

→ *Place*, p_{od} with $p_{od.name} = "ODummy_" + e.id$;

→ *Arc*, a with $a.source = p_{id}$; $a.target = t$;

→ *Arc*, a with $a.source = t$; $a.target = p_{od}$;

For every Constraint c in Episode e :

→ *Place*, p : $p.name = c.name$; $p.tokens = 1$;

→ *Arc*, a with $a.source = p$; $a.target = t$;

For every Result r in Episode e :

→ *Place*, p with $p.name = r.name$;

→ *Arc*, a with $a.source = t$; $a.target = p$;

For every Exception ex in Exceptions:

IF $ex.id$ starts with $e.id$:

→ *Transition*, t_{ex} with $t_{ex.name} = ex.solution$;

→ *Arc*, a with $a.source = p_{od}$; $a.target = t_{ex}$;

IF $ex.cause \neq \emptyset$:

→ *Place*, p : $p.name = ex.cause$; $p.tokens = 1$;

→ *Arc*, a : $a.source = p$; $a.target = t_{ex}$;

For every Result r in Exception ex :

→ *Place*, p with $p.name = r.name$;

→ *Arc*, a : $a.source = t_{ex}$; $a.target = p$;

→ Remove exception ex from *Exceptions*;

IF $|e.conditions| > 0$:

→ *Transition*, t_{else} with $t_{else.name} = "ELSE_" + e.episode_sentence$;

→ *Arc*, a with $a.source = p_{id}$; $a.target = t_{else}$;

→ *Arc*, a with $a.source = t_{else}$; $a.target = p_{od}$;

For all Condition cd in Episode e :

→ *Place*, p : $p.name = cd.name$; $p.tokens = 1$;

→ *Arc*, a with $a.source = p$; $a.target = t$;

IF $|e.conditions| > 0 \wedge e$ has exceptions:

→ Remove *Arc* a between t_{else} and p_{od} ;

\rightarrow Transition, t_d : $t_d.name = "TDummy_" + e.id$;
 \rightarrow Arc, a with $a.source = p_{od}$; $a.target = t_d$;
 \rightarrow Place, p_{od} : $p_{od.name} = "ODummy2_" + e.id$;
 \rightarrow Arc, a with $a.source = t_d$; $a.target = p_{od}$;
 \rightarrow Arc, a with $a.source = t_{else}$; $a.target = p_{od}$;

Transformation of episodes is shown in Figure 4.

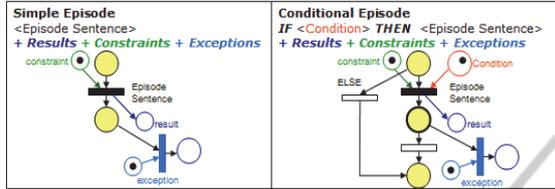


Figure 4: Episode mapping rules.

4.3 Mapping Exceptions

Each one of the remaining exceptions are mapped into the “transition” t ($t.name = solution$) and its internal “input dummy place” p_{id} of a Sub-Petri-Net. The “condition” or “cause” to trigger the transition t is mapped into an “input place” with an initial token.

A *result* is an “output place” generated at the end of perform the transition t .

Places and Transitions are generated as follows:

For every Exception ex in Exceptions:

\rightarrow Place, p_{id} with $p_{id.name} = "IDummy_" + ex.id$;
 \rightarrow Transition, t with $t.name = ex.solution$;
 \rightarrow Arc, a with $a.source = p_{id}$; $a.target = t$;

IF $ex.cause \neq \emptyset$:

\rightarrow Place, p : $p.name = e.cause$; $p.tokens = 1$;
 \rightarrow Arc, a with $a.source = p$; $a.target = t$;

For every Result r in Exception ex :

\rightarrow Place, p with $p.name = r.name$;
 \rightarrow Arc, a with $a.source = t$; $a.target = p$;

The mapping of the exceptions or alternative flows of the episodes is illustrated in Figure 5.

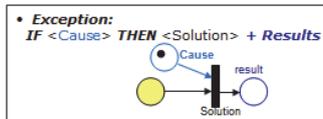


Figure 5: Exception mapping rules.

4.4 Mapping Concurrency Constructs

If the *episode sentence* of an *episode* e starts or ends with the symbol “#”, this symbol describes the start or the end (synchronization) of multiple concurrent episodes, respectively. These situations are mapped into two Sub-Petri-Nets composed of the transition *fork* and the transition *join*, respectively (and their internal *input dummy place* p_{id} and *output dummy*

place p_{od}).

For every Episode e in Episodes:

IF $e.episode_sentence$ starts with “#”:

\rightarrow Place, p_{id} with $p_{id.name} = "IFork_" + e.id$;
 \rightarrow Transition, t with $t.name = "Fork_" + e.id$;
 \rightarrow Place, p_{od} with $p_{od.name} = "OFork_" + e.id$;
 \rightarrow Arc, a with $a.source = p_{id}$; $a.target = t$;
 \rightarrow Arc, a with $a.source = t$; $a.target = p_{od}$;

IF $e.episode_sentence$ ends with “#”:

\rightarrow Place, p_{id} with $p_{id.name} = "IJoin_" + e.id$;
 \rightarrow Transition, t with $t.name = "Join_" + e.id$;
 \rightarrow Place, p_{od} with $p_{od.name} = "OJoin_" + e.id$;
 \rightarrow Arc, a with $a.source = p_{id}$; $a.target = t$;
 \rightarrow Arc, a with $a.source = t$; $a.target = p_{od}$;

The mapping of concurrency constructs is illustrated in Figure 6.

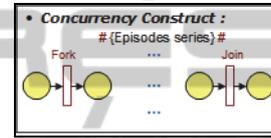


Figure 6: Concurrency mapping rule.

4.5 Composing the Elements

This section explains the steps to compose a complete Petri-Net model from the Sub-Petri-Nets obtained of the elements of a scenario description.

- Fusion Places:** If an “input/output place” derived from a scenario element appears like “input/output place”, it should be merged into the “input/output place” of the later one.
- Linking Concurrent Episodes:** The Sub-Petri-Nets derived from the episodes between a “fork” Sub-Petri-Net and a “join” Sub-Petri-Net must be linked as concurrent Sub-Petri-Nets and composed into a complete Sub-Petri-Net. The “input dummy place” of the “Sub-Petri-Nets” of the episodes are linked to the “transition” of the “fork” Sub-Petri-Net. The “output dummy place” of the “Sub-Petri-Nets” of the episodes are linked to the “transition” of the “join” Sub-Petri-Net.
- Composing The Sub-Petri-Nets:** First, all the places generated from the context and the resources are “input places” of the first transition derived of the first episode. Second, the “output dummy place” of the previous Sub-Petri-Net is merged into the “input dummy place” of the later one (Fusion Place).

4.6 Integration of Petri-Net Models

Scenarios are related to other scenarios by sequential and non-sequential interactions. Two techniques, “*Fusion Places*” and “*Substitution Places*” are used to obtain a complete Petri-Net model of the application (composed of Petri-Net models derived from different scenarios).

Sequential Interactions: As described before, these interactions could be of five types (integration scenarios, pre-condition, constraint, sub-scenario and exception) and determine the order in which the scenarios should be executed (Section 3).

Non-sequential Interactions: Scenarios interact by shared resources described as: pre-condition, constraint and result. Through these relationships it is possible to identify the scenarios that could be executed concurrently.

If an “*input/output place*” derived from *Pre-Conditions*, *Constraints* or *Results* of a scenario appears like “*input/output place*” on other Petri-Net model derived from other scenario, it should be merged into the “*input/output place*” of the later one (*Fusion Places*).

5 CASE STUDIES

In this section is described how the template and mapping rules, detailed in Section 3 and 4 were used to create the Petri-Net models of some classical concurrency problems, using as basis their scenario descriptions. The complete Petri-Net models were obtained through the interconnection of Sub-Petri-Nets and the integration of Petri-Net models.

The order which Petri-Net models are constructed is: **First**, for each scenario, transforming the context (“*Pre-condition*”) and the resources of the scenario into a Sub-Petri-Net, transforming “*Episodes*” of scenario into Sub-Petri-Nets sequentially, linking “*output places*” (episode’s results) into related episodes Sub-Petri-Nets, and composing the “*Sub-Petri-Nets*” into a complete Petri-Net model; **Second**, integration of Petri-Net models by sequential & non-sequential interactions.

5.1 Bounded Producer-consumer

In multithreaded programs, there is often a division of labor between threads. In one common pattern, some threads are producers and some are consumers. The producer generates one item and puts it in the buffer. The consumer gets one item of the buffer and extinguishes it (Downey, 2005).

Problem: Make sure that the producer won't try to add an item into the buffer if it is full and that the consumer won't try to remove an item from an empty buffer.

- The *buffer* is shared and it has only *N* positions for items.

Solution: A “*Customer*” performs his consume only when a “*Producer*” puts an item in the Buffer (“*Buffer is not empty*”). A “*Producer*” produces an item only when the “*Buffer is not full*”.

Figure 7 depicts the scenarios to describe a simple solution for the “*Producer-Consumer Problem*” and the corresponding Petri-Net model.

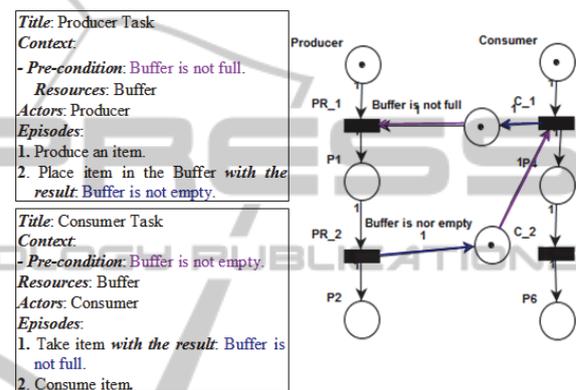


Figure 7: Bounded Producer-Consumer Petri-Net.

The *output place* “**Buffer is not empty**” of the “*Producer*” is merged with the *initial input place* “**Buffer is not empty**” of the “*Consumer*”. The *output place* “**Buffer is not full**” of the “*Consumer*” is merged with the *initial input place* “**Buffer is not full**” of the “*Producer*”.

5.2 Readers-Writers Problem

The Reader-Writer Problem pertains to any situation where a data structure, database, or file system is read and modified by concurrent processes: *Readers* and *Writers*. While the data structure is being written or modified it is often necessary to bar other processes from reading, in order to prevent a reader from interrupting a modification in progress and reading inconsistent or invalid data (Downey, 2005).

Problem: Allow multiple readers to read at the same time. Only one single writer can access the shared data at the same time.

- *Reader*: only read the data set; it does not perform any updates.
- *Writer*: can both read and write.

Solution: Once a “*Writer Task*” is ready (“*Data Set is available*”), it gets to perform its write as soon as

possible and locks (“Data Set is not available”) the other “Writer” and “Reader Tasks”.

Figure 8 depicts the scenarios to describe a simple solution for the “Readers-Writers Problem” and the corresponding Petri-Net model.

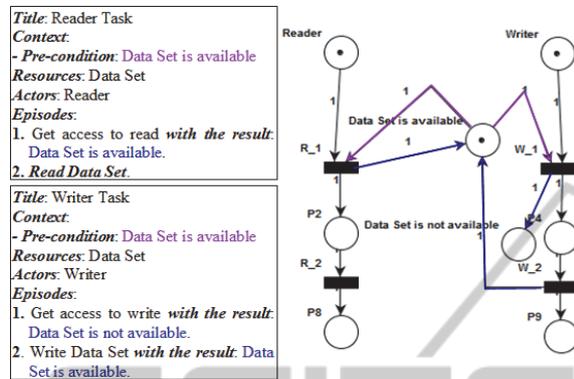


Figure 8: Readers-Writers Problem Petri-Net Model.

The initial input place “Data Set is available” of the “Reader” is merged with the initial input place “Data Set is available” of the “Writer”. The output place “Data Set is available” of the “Reader” is merged with the output place “Data Set is available” of the “Writer”.

6 CONCLUSIONS

This work proposes mapping rules for transforming a specific natural language-based scenario model into a Petri-Net model. The resulting Petri-Net model can be analyzed using available tools (PIPE2, 2014). Such tools are able to identify the following problems in a Petri Net model: boundedness, safety and deadlock. As such, if this earlier feedback is available, it is possible to trace backwards to scenario descriptions and fix problems earlier on.

The main contributions of this work are summarized as follows: (1) Describe and update an abstract and concrete syntax for describing scenarios. (2) Define mapping rules from Scenario descriptions to Petri-Nets. (3) Express each mapping rule in terms of “Data related issues” (Constraints, Pre-Conditions and Results). (4) Define rules for integrating Petri-Net models derived from different scenarios of the application (modularity). (5) Provide examples considering concurrent scenarios in order to show the interactions by concurrency.

The contributions of this work provide convenient ways to make explicit the non-sequential interactions among scenarios by shared resources. This fact can be used for further analysis of

properties (like deadlock and conflict) and testing.

Our future research plan will consider the following tasks: (1) Define criteria to identify inconsistency, incompleteness and incorrectness for intra-scenario and inter-scenario relationships based on analysis of structural and behavioral properties of Petri-Nets. (2) Investigate strategies, which automatically traverse the Petri-Net model to generate the test scenarios based on path analysis strategies. This strategy will take into account interactions by “shared resources”.

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