

Linear Software Models for Well-composed Systems

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Abstract: Modularity is essential for automatic composition of real software systems from ready-made components. But given ready-made components do not necessarily correspond exactly to the units and functionality of designed software system architecture modules. One needs a neat composition procedure that guarantees the necessary and sufficient components to provide required units. Linear Software Models are rigorous theoretical standards subsuming modularity. The Linear-Reducible model is proposed as a model of well-composed software systems, above and beyond software variability. Indeed, case studies of representative systems recognized as well-composed, be they small, intermediate building blocks or large scale, are shown to be Linear-Reducible. The paper lays down theoretical foundations – upon exact linear independence and reducible matrix concepts – providing new precise meanings to familiar modularity ideas, such as the single responsibility theorem. The theory uses a Modularity Matrix – linking independent software structures to composable software functionals in a Linear Model.

1 INTRODUCTION

Significant progress since Parnas' classical paper (Parnas, 1972) – which posed the modularity issue in the software context and treated it by informal reasoning – paved the way for automatic run-time system composition/update from ready-made software components. Yet there remain fundamental obstacles to make this vision concrete.

Modularity's wisdom of low dependency among modules has been informally stated in innumerable ways: recommendations such as single responsibility, source-code dependency metrics, design patterns and tools. But recommendations, metrics, patterns and tools never crystallized into a systematic theoretical approach.

This is exactly the problem dealt with by this paper: to provide a solid and generic basis to treat software composition in a rigorous and consistent way, enabling theoretical models against which to check real systems. To this end, Linear Software Models are proposed upon well-established linear algebra techniques.

1.1 The Software Composition Problem

The software composition problem, analysed by this paper, is how to build a well-designed modular software system from available ready-made components that were not designed specifically for a particular system.

Components are deployment units needed to actually run software systems. They are loaded to the computer memory as indivisible wholes. Typical examples are a C++ dll (dynamically linked library), a jar (Java archive) or a C# assembly.

On the one hand, the software engineer, by using best practices, designs a modular software system in terms of desirable architectural units. Architectural units describe the structure and behaviour of a particular software system.

On the other hand, *components* are assumed to be mainly purchased as COTS (Commercial Off-The-Shelf) components from several manufacturers, and less frequently to be produced in-house. It is realistic to expect variability among COTS from distinct sources.

1.2 Linear Software Models: Structors and Functionals

This paper describes Linear Software Models as a theory of software composition. In this theory, the architecture of a software system is expressed by two kinds of entities: structors and functionals.

Structors – a new term reminding vectors – are architectural units, from the structural point of view. Structors generalize the notion of structural unit to cover diversity of types (structs, classes, interfaces, aspects) and hierarchical collections (sets of classes, as design patterns). Structors refer to types, not instances. Structors are loadable within components.

Functionals are architectural system units from a behavioural point of view. These are *potential functions* that can be, but are not necessarily invoked. Typically these are Java or C# methods, related functions (e.g. a set of trigonometric functions) or roles – supplying the functionality of a design pattern (Riehle, 1996). Note: we use *Functional* as a noun, similarly to the mathematical concept with this name in the calculus of variations, and to the grammatical use of *Potential* in physics.

Structors in general contain a finite set of functionals and are represented by finite vectors.

Modules are architectural units in a higher hierarchical level of a system. Modules are composed of grouped structors and their corresponding functionals.

Both components and systems contain structors. But the respective logics are different. Structors contained in a COTS component are fixed by the core technology of the component manufacturer and such structors must be assumed indivisible.

Structors within a system are determined by the software system purpose. Thus, not all structors of a component may be required by a system. Similarly, not all functionals provided by structors are needed by a system, and some of them may never be invoked. The numbers of structors or functionals provided by a component are not constrained.

Analysis starts with a list of structors and a list of functionals that must be in a system. If two structors provide distinct functionals, both are needed. If they provide the same functionals, one of them is redundant. For partial overlap, either one may be complemented by a third structor. But which one is preferable? Within a Linear Software Model the answer is clear: choose a linearly independent set of structors.

Linear models are usually formulated in terms of matrices. The Modularity Matrix is a Boolean matrix with columns standing for structors and rows for

functionals. A matrix element is 1-valued for a functional-structor link and 0-valued for no link.

1.3 Modularity Matrix: An Introductory Example

A simple yet useful system is here described in terms of a Linear Software Model to illustrate concepts found along the paper. It refers to OFB (Output Feed-Back) encryption of a long message (Fig. 1) before network transmission – e.g. (Kaufman, et al. 1997). A message is cut into N blocks of fixed length and each one is treated separately.

A random number, the *Initial Vector*, is generated. Each block of the message is pre-processed. A corresponding vector is also pre-processed to be of the size of a block, and then encrypted. The encrypted *output* vector is *fed back* into the encryption function, through the vector pre-processing. Feedback is done N times, obtaining one vector for each message block. The k^{th} encrypted vector is then *Xored* with the k^{th} message block.

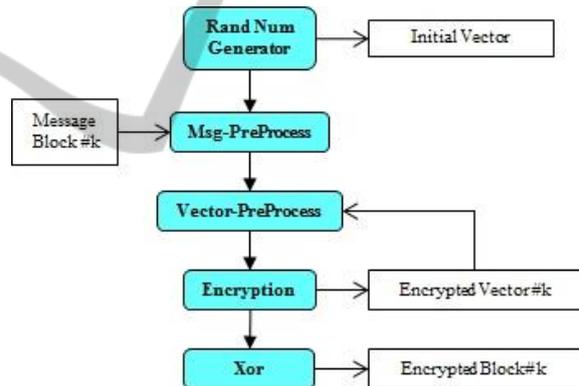


Figure 1: OFB encryption of long message. This is a functional calling dependency graph for a generic k^{th} message block. Functionals are displayed as (blue) rounded rectangles. Some of the input/output dataflow is shown in regular rectangles.

Table 1: OFB Functionals.

#	Functional	Description
1	Random number generator	from a distribution
2	Encryption function	e.g. RSA encryption
3	Xor	a logical function
4	Vector-Pre-Process	Fetch(vector,size)
5	Message-Pre-Process	fetch & pad message

The OFB software system has five functionals (Table 1) and the following four structors: S1) Rand - offers random distributions; S2) Crypto structor -

offers encryption protocols; S3) Logical - provides logical functions, say AND, OR, XOR; S4) Proc - a processor structor provides pre-processing functions.

The resulting OFB modularity matrix (Table 2) has 2 identical lowest rows. The matrix reflects that the pre-processing functionals are not independent. Both functionals are provided by the same structor, thus they are not distinguishable.

Table 2: OFB Rectangular Modularity Matrix with Linear Dependencies.

		S1=Rand	S2=Crypto	S3=Logical	S4=Proc
Rand Num Gen	F1	1	0	0	0
Encryption Func	F2	0	1	0	1
XOR	F3	0	0	1	1
Vector-Pre-Proc	F4	0	0	0	1
Msg-Pre-Proc	F5	0	0	0	1

An interesting feature of OFB encryption is that one can prepare all encrypted vectors in advance, before there are any message blocks to be sent. Therefore, one can rearrange the functional calling graph, enabling totally separate execution.

Two modules naturally appear, one dealing with vectors and the other with message blocks – see Fig. 2. In each module, the processing structor, besides pre-processing, invokes the respective processing functional.

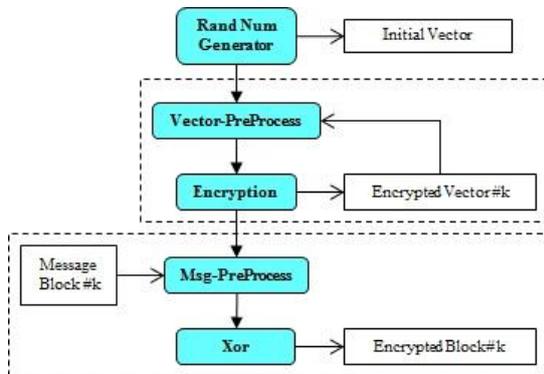


Figure 2: OFB functionals modularized. Rearranged functional calling dependency graph for a generic k^{th} message block. The upper dashed lines module pre-processes vectors, encrypts and saves them. The lower module fetches message blocks, pre-processes them and Xors each message block with its saved vector. This figure uses the same conventions as the previous figure.

Independent pre-processing functionals which are able to run in parallel in distinct machines need independent structors. So, one has a VectorProc structor and a separate MsgProc structor. These structors are aware – by means of associations – of the respective processing functionals which they are able to invoke.

Table 3: OFB Strictly Linear Model – Modularity Matrix.

		S1=Rand	S2=Crypt	S4=Vec Proc	S3=Logical	S5=Msg Proc
Rand Num Gen	F1	1	0	0	0	0
Encryption Func	F2	0	1	1	0	0
Vector-Pre-Proc	F3	0	0	1	0	0
XOR	F4	0	0	0	1	1
Msg-Pre-Proc	F5	0	0	0	0	1

Row and column swap operations lead to a block-diagonal matrix (Table 3). The diagonal blocks in this matrix match the modules in Fig. 2.

The system matrix in Table 3 strictly obeys a Linear Software Model. This means that all its modules have linearly independent rows and columns. Thus, each of its functionals is distinguishable. For instance, according to the matrix in Table 3 functional F_2 is provided by structors $\{S_2, S_4\}$.

The independent execution of VectorProc and MsgProc, in time and space, clarifies the sense of an independent software structor. It must be: a) *Loadable/Runnable* – in a virtual or real machine; b) *Separable* – i.e. able to run in *separate* machines.

The remaining of the paper is organized starting from the basic theory (section 2), through concrete case studies (section 3), to a discussion (section 4).

2 LINEAR SOFTWARE MODELS OF COMPOSITION

The aim of this section is to describe the new theoretical approach – the Linear Software Models of Composition.

Linear Software Models are the simplest theoretical models of software composition. Systems obeying such a model are composed just by addition of independent modules.

2.1 Modularity Matrices' Linear Independence

A Linear Software Model contains a list of software structors and another of software functionals. Its Modularity Matrix is defined as:

Definition 1 – Modularity-Matrix

A fully expanded *Modularity-Matrix* is a Boolean matrix asserting links (1-valued elements) between software functionals (rows) and software structors (columns). The absence of a link is marked by a 0-valued element.

By definition, software structors are elementary artifacts, which the software engineer decides to look at them as indivisible into smaller structors. Say, the OFB crypto structor (section 1.3) is not split into – prime factorization or modulo arithmetic – although decomposition is obviously possible. The same holds true for functionals.

Besides being indivisible, only structors with unique roles, as the OFB VectorProcessor and MsgProcessor needed for parallelism, are in the Matrix. Multiple structor copies, say by fault tolerance reasons, are not included in the Matrix.

Independent structors must be represented by distinct vectors. But, sets of differing vectors may still be dependent. The generic criterion for independent structors in any system subsets is linear independence:

Definition 2 – Independent Structor

A software *structor* is *independent* of other structors in the system, if it provides a non-empty proper sub-set of functionals of the system, given by the 1-valued links in the respective column, and is linearly independent of other columns in the Modularity Matrix.

We now look at the *links* from the functional point of view. In order to be able to distinguish a functional, its set of links in the functional row must be unequivocal. Again linear independence is the relevant criterion:

Definition 3 – Composable Functional

A software *functional* is independently composable or just *composable* in terms of structors, if it corresponds to a non-empty proper sub-set of system structors, given by the 1-valued links in the respective row, and is linearly independent of other rows in the Modularity-Matrix.

For instance, for OFB (Table 3) the XOR functional composition set is {S3, S5}.

2.2 Well-composed Modularity Matrices are Square

A well-composed modularity matrix has additive properties, i.e. it has composable functionals from independent structors. Then:

Theorem 1 – Well-Composed Modularity-Matrix

If in a Modularity-Matrix all its functionals are composable with independent structors, the Modularity-Matrix number of structors N_S is equal to its number of functionals N_F . The matrix is Square.

Such a matrix is called a *Well-Composed Modularity-Matrix*.

A proof sketch is (detailed proofs will be given in a long paper): Assume a matrix without empty rows/columns. First, structors are used as a basis for functional vectors. Then, functionals are a basis for structor vectors. By linear independence, in each case vector numbers cannot be greater than the basis. As both cases must be simultaneously true, follows the equality $N_S=N_F$. The matrix is square.

The theorem assumptions are very intuitive. An analogy is the symptom sets to diagnose illnesses. Say, high temperature, a very common symptom, is not enough to identify a disease. Diseases with the same symptom sets are indistinguishable. To unequivocally diagnose an illness one needs independent sets of symptoms.

This theorem does not force a one-to-one functional/structor match. It is obeyed by matrices with N_{VAL} 1-valued elements greater than N_S and N_F (see e.g. Table 4).

Table 4: Abstract square Modularity Matrix with $N_{VAL}=6$, while $N_S=N_F=3$.

	S1	S2	S3
F1	1	0	1
F2	0	1	1
F3	1	1	0

Algebraically, there are simple criteria for well-composed matrices: non-zero matrix determinant or matrix rank equal to the number of rows/columns.

2.3 Reducible Modularity Matrices

When a Modularity Matrix has disjoint dependency sets, structors and their functionals can be grouped into modules, such that vector sets in different modules are linearly independent. This reduces the matrix to block-diagonal (fig. 3), i.e. with smaller squares of any size along the diagonal (Rowland, Weisstein, 2006). All off-block elements are zero.

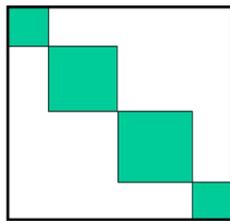


Figure 3: Schematic Block-Diagonal Matrix.

A “union-set” is the union of *composition sets* for a set of functionals. Then:

Theorem 2 – Modularity Matrix Reducibility

Any well-composed Modularity Matrix in which the union-set for a given set of functionals is disjoint to the other functional union-sets is reducible, i.e. it can be put in block-diagonal form.

A proof sketch is: apply row and column exchanges to bring the disjoint union-set to the upper-left matrix corner. As the whole Modularity Matrix is *Well-composed*, the disjoint union-set is itself square (by Theorem 1) and its diagonal is along the whole Matrix diagonal. The residual 1-valued elements are thus also square bounded, along the whole Matrix diagonal. This reasoning is extensible to any number of blocks.

The previous theorems naturally lead us to define linear software system models of composition. The two theorems are combined in the Linear-Reducible model, in the strict sense.

Definition 4 – Linear Reducible Model

The *Linear-Reducible* model of software system composition, in the strict sense, is characterized by a well-composed and reducible Modularity Matrix having at least two disjoint blocks.

We identify modules with disjoint diagonal blocks of structors/functionals, corresponding to the intuitive notion of modular software systems. One can express modularity quantitatively by the diagonality of a Modularity-Matrix M , telling how close its 1-valued elements are to the main diagonal. It is the difference between the Trace, the diagonal elements' sum, and *offdiag*, a new term dealing with off-diagonal elements:

$$Diagonality(M) = Trace(M) - offdiag(M) \quad (1)$$

Offdiag sums over 1-valued M_{jk} off-diagonal elements, times the absolute value of the difference

of the element's column k and row j indices. For each row j and column k :

$$RowOffdiag(M, j) = \sum_{k=1}^N M_{jk} \cdot |j - k| \quad (2)$$

$$ColOffdiag(M, k) = \sum_{j=1}^N M_{jk} \cdot |j - k|$$

The overall matrix offdiag is:

$$offdiag(M) = \sum_{j=1}^N \sum_{k=1}^N M_{jk} \cdot |j - k| \quad (3)$$

The next operations may unfold a whole hierarchy of block levels, where a level is defined by the matrix components explicit in that level:

- a- Block collapse – transforms a block into a single element *black-box*.
- b- Block expansion – restores a collapsed black-box back into a *white-box*.

To assure that these operations preserve diagonality, black-boxes are labeled by the <Trace, offdiag> diagonality value, of the original white-box. To be able to restore the white-box, hidden rows/columns should be stored elsewhere.

The modularity hierarchy is thus given by:

Definition 5 – Number of modules at a level

The number of modules of a software system at a hierarchy level is the number of blocks from the modularity matrix partition into disjoint union-sets at that level.

2.3 The Single Responsibility Theorem

The last theoretical piece of the Linear Models is a linear algebraic formulation of the single responsibility principle (Martin, 2003). It is valid after one obtains a block-diagonal matrix:

Theorem 3 – Single Responsibility

In a strictly block-diagonal modularity matrix each structor column intersects a single module. Similarly, each functional row intersects a single module.

This theorem is easily verified. It means that a single module is responsible for providing each functional exclusively by its structors. This is a plausible interpretation of the familiar principle,

since related structors and functionals are grouped in a single module.

3 CASE STUDIES

The goal of this section is to demonstrate the power of Linear Software models, showing by case studies a range of concrete software systems obeying such a model.

We have chosen representative software systems of disparate purposes and sizes to illustrate the Linear Models. It turns out that all systems are *well-composed*. Small systems are strictly Linear-Reducible – i.e. obey a model with a linearly independent and reducible matrix – and larger systems are bordered Linear-Reducible. In each case, a modularity matrix is obtained, block-diagonalized and analysed in terms of Linearity.

3.1 Small Systems are Strictly Linear Reducible

Small systems and intermediate reusable building blocks strictly obey the Linear-Reducible Software model. These are Parnas’ KWIC index and the *Observer* pattern. KWIC was thoroughly analysed by Parnas to be a canonical example. The *Observer* pattern was deliberately designed to be reused.

3.1.1 Parnas' KWIC Index

The 1972 Parnas paper (Parnas, 1972) described two KWIC index modularizations. The system outputs an alphabetical listing of all circular shifts of all input lines. Functionals were extracted from Parnas’ own system description. Structors are explicit in the paper – where they are called modules.

The Modularity Matrix of Parnas’ 1st modularization is in Table 5 (0-valued elements are omitted for clarity; blocks have a darker background). It is almost block-diagonal, but two features are problematic:

- a- the Master Control column vector is not a proper subset of the functionals;
- b- non-zero outliers break block-diagonality.

The matrix clearly hints to couplings to be resolved: the notably higher *RowOffdiag*=4 of the 3rd row, contains all outliers, besides the master control ones. The solution is: a- to delete the master control, as quoting Parnas, it “does little more than sequencing”, it is not a real structor; b- to add a new Line-Storage structor, as Parnas informally argued,

Table 5: Parnas' 1st Modularization – Modularity Matrix.

Structors →		Input	Circular Shifter	Master Control	Alpha-betizer	Output
Functionals		1	2	3	4	5
Input = ordered set of lines	1	1		1		
Does circular shift on a line	2		1	1		
Line= store line in word order	3	1	1	1	1	
Sort lines in alphabetical order	4			1	1	
Outputs circular shifted lines	5			1		1

to decouple the 3rd row "Line" functional from the Input, Circular-shifter and Alphabetizer. The 2nd Parnas’ modularization fits the strictly diagonal 5-module matrix in Table 6.

Table 6: Parnas' 2nd Modularization – Modularity Matrix.

Structors →		Input	Circular Shifter	Line Storage	Alpha-betizer	Output
Functionals		1	2	3	4	5
Input = ordered set of lines	1	1				
Does circular shift on a line	2		1			
Line= store line in word order	3			1		
Sort lines in alphabetical order	4				1	
Outputs circular shifted lines	5					1

We use *decouple* in a precise new meaning, of enabling linearly independent composition.

The matrices in Tables 5 and 6 are not equivalent, having different non-zero element numbers. Which one is more modular? The 1st matrix diagonality has a value –5. The 2nd diagonality equals the Trace and is 5. Thus, the 2nd is more modular.

We arrived at the same Parnas conclusions, by formal Modularity Matrix arguments. This system strictly obeys the Linear-Reducible Model.

3.1.2 Observer Design Pattern

The *Observer* Design Pattern abstracts one-to-many interactions among objects, such that when a "subject" – changes, all the "observers" – are notified and updated.

The *Observer* Modularity Matrix is based upon the Design Patterns’ GoF book (Gamma, et al., 1995). Its sample code refers to an analog and a digital clock, the "concrete observers", changing according to an internal clock – the "concrete subject".

The system structors (Table 7) are:

- generic *Observer* pattern entities, directly taken from the pattern list of “Participants”; (the first four rows of Table 7).
- specific application structors – a subject resource, a GUI for each external clock and an initiator which constructs the clocks.

Table 7: Observer - Structors.

	Structors	Description
1	abstract subject	interface to attach/detach observers
2	abstract observer	interface to update an observer
3	concrete subject	stores global state; sends notifications
4	concrete observer	implements the observer updating
5	subject resource	internal timer ticks the global state
6	analog GUI	inherits graphics analog clock widget
7	digital GUI	inherits graphics digital clock widget
8	initiator	to construct the clock objects

The Observer functionals were extracted from the complete set of functions that may be invoked in the sample code. These functions were trimmed by elimination of linear dependencies.

State is not maintained in the same way in the subject and in the observers. The unique subject sets the state ("ticks") by means of a Keep-Global-state functional, while observers, are updated at unrelated times and Keep-Local-state. Table 8 shows the Observer functionals.

Table 8: Observer – Functionals.

	Functionals	Description
1	Keep-observers-list	attaching/detaching observers
2	Notify observers	notify when subject changes
3	Update observers	update observers, after notify
4	Keep-Global-state	keep time to allow updates
5	Keep-Local-state	keep time in each observer
6	Draw-analog	specific analog clock draw
7	Draw-digital	specific digital clock draw
8	Constructor	construct the clock objects

Row/column reordering of a quite arbitrary initial matrix causes modules to emerge in a strictly Linear-Reducible matrix. (Table 9). These modules are a subject and an observer – the generic module roles (Riehle, 1996) for this design pattern – each one with the respective abstract/concrete structors, and the specific clock application modules.

Table 9: Observer Linear Reducible – Modularity Matrix.

Structors →		subject	concrete subject	subject resource	concrete observer	Observer	Gui analog	Gui digital	Init
Functionals		1	2	3	4	5	6	7	8
Keep observer list	1	1							
Notify observers	2	1	1						
Keep-global-state	3		1	1					
Keep-local-state	4				1				
Update observers	5				1	1			
Draw-analog	6						1		
Draw-digital	7							1	
Constructor	8								1

In Table 9, subject and observer modules emerged from basic structors. Alternatively one can deal with black-box collapsed modules (as defined in sub-section 2.3) in a higher hierarchical level matrix shown in Table 10.

In either format, one could analyse the pattern matrix (subject and observer) as a generic building block in separate from the specific clock application.

Table 10: Observer higher level Modularity-Matrix.

Structors →		Subject	Observer	Clock Appl.
Functionals		1	2	3
Keep-global-state & Notify observers	1	<3,2>		
Keep-local-state & Update observers	2		<2,1>	
construct, draw clocks	3			<3,0>

The Observer analysis illustrates that, despite arbitrary initial order, automatic reordering brings about a matrix accurately reflecting the pattern functionality.

The Observer pattern is a prototypical example. It would be desirable to have all reusable building blocks as the Observer, strictly obeying the Linear-Reducible Model, to allow linearly independent composition into larger software systems.

3.2 Larger Software Systems are Bordered Linear Reducible

To show the Linear Model applicability to real systems, we analysed larger projects from the literature. The novel result is that the examined systems are bordered Linear-Reducible.

3.2.1 Neesgrid Modularity Matrix

The NEESgrid “Network Earthquake Engineering Simulation” project enables network access to participate in earthquake tele-operation experiments. The infrastructure was designed by the NCSA at University of Illinois. Modularity Matrix functionals were extracted from a report (Finholt, et al. 2004) with exactly 10 upper-level structors.

An initial modularity matrix (Table 11) was obtained by straightforward linearity considerations: eliminating empty and identical rows; an empty column, “Electronic Lab Notebook”, was deleted. A column was assigned to a next-level sturctor *Data-Discovery* (DataDis) as it neatly fits the *SearchData* functional. The scattered non-zero elements are typical of initial matrices in this kind of analysis.

Table 11: NEESgrid initial Modularity Matrix.

Structors →		chef	Data Rp	Data Vu	Data Str	Data Dis	Tele pre	Data Ac	Hyb Exp	Sim Rep	Grid Infr
Functionals		1	2	3	4	5	6	7	8	9	10
Collect_Data	1	1		1				1			
Search_Data	2				1						
Manage_Data	3	1	1								
HybridExper	4								1		
Data_View	5			1							
Sync_Collab	6	1					1				
Async_Collab	7	1									
Other_Collab	8	1		1				1			1
SimulCodes	9									1	
HighPerfComp	10										1

Pure algebraic row/column reorder, without semantic concerns, brings about the almost block-diagonal Matrix (in Table 12). Its modules – diagonal blocks – are:

- Data manipulation – collect, search, manage, view – the upper-left block;
- Collaboration tools – synchronous, async, other – the middle block;
- Infrastructure – grid, codes – the lower-right block.

Table 12: NEESgrid Bordered Linear-Reducible Modularity Matrix.

Structors →		Data Str	Data Dis	Data Rp	Data Vu	Data Ac	Tele pre	chef	Grid Infr	Sim Rep	Hybr Exp
Functionals		4	5	2	3	7	6	1	10	9	8
Collect_Data	1	1		1		1					
Search_Data	2		1								
Manage_Data	3			1	1						
Data_View	5				1						
Other_Collab	8				1	1		1	1		
Sync_Collab	6						1	1			
Async_Collab	7							1			
HighPerfComp	10								1		
SimulCodes	9									1	
HybridExper	4										1

Module interpretation is hinted by the respective functional names – *Data* and *Collab*.

Outliers appear in the rows with maximal value $\text{RowOffDiag}=6$, *Collect_Data* and *Other_Collab*. The latter prefix "Other" hints at mixed functionals to be decoupled.

This case study shows that: a- algebraic reordering, without prior semantic knowledge, obtains plausible modules; b- outliers are amenable to interpretation, in particular matching project notes

(e.g. the row 8, column 10 outlier, marked "not within the project scope" in project documents).

The significant result, common to large case studies, is that there are few outliers, and all of them are in columns/rows adjacent to the Linear Model blocks. This is what we call bordered Linear-Reducible.

4 DISCUSSION

4.1 Main Contribution: Linear Software Models

This paper's main contribution is the Linear Software Models, as theoretical standards against which to compare real software systems. The models stand upon well-established linear algebra, as a broad basis for a solid theory of composition – beyond current principles and practices.

One can assert, from the Modularity Matrix properties of a system, which structors are independent and which functionals are independently composable. One can then infer which design improvements are desirable.

This view is very different from *design* models, such as UML, whose purpose is not to serve as theoretical standards. Design models freely evolve with design and system development. Design models have indefinite modifiability to adapt to any system, in response to tests of system compliance to design.

4.2 Related Work

Matrices have been used to deal with modularity. A prominent example is DSM (Design Structure Matrix) proposed by Steward (Steward, 1981), developed by Eppinger and collaborators e.g. (Sosa, et al. 2005), (Sosa, et al. 2007)) and part of the "Design Rules" approach by Baldwin and Clark – see e.g. (Baldwin, Clark, 2000) and also (Cai, Sullivan, 2006), (Sethi, et al. 2009). DSM and other matrices, such as Kusiak and Huang's (Kusiak, Huang, 1997) *hardware* modularity matrix, are meant to be evolving design models.

Linearity is the outstanding feature of our standard models, not found in DSM. Essential distinctions of Linear Software Models from DSM are:

- Theoretical Standards vs. Design Models – our models' goal is to serve as system standards and ultimately may lead to automatic software composition, as

opposed to DSM design models which emphasize design process and manufacturing.

- Functionals vs. Structures-only – our modularity matrices display structure to functional links, while both DSM matrix dimensions are labeled by the same structures.

Baldwin and Clark explicitly state in footnote 2, page 63 of their book (Baldwin, Clark, 2000) that “it is difficult to base modularity on functions... hence their definition of modularity is based on relationships among structures, not functions”. See (Ulrich, 1995) for a different view.

In practice, modularity matrices may be much more compact than DSM. For instance, Parnas’ KWIC DSM (Cai, Sullivan, 2006) has 20 rows/columns instead of just 5 in our Modularity Matrix.

Although diagonality has seldom been calculated within modularity, formulas have appeared in other contexts. Clemins in (Clemins, et al. 2002) used for speaker identification, the Frobenius norm – the sum of the squares – of all off-diagonal elements to measure diagonality. Our *offdiag* definition is better suited to modularity, as it directly reflects distance to the diagonal, while the Frobenius norm just sums Boolean elements.

An indirect coupling metric is “similarity coefficients” (Hwang, Oh, 2003), comparing matrix row pairs, over all columns. The similarity for each column is: both 1 elements, both 0, or different. These coefficients ignore distances from the diagonal.

The Modularity Matrix has a superficial similarity to a traceability table. But their purposes are definitely different. Traceability tables are used to *trace* code and tests to requirements, while our functionals’ essence is to obtain measures of linear independence.

The module detection literature is plentiful. Tools to improve legacy code use clustering to partition graphs (Mitchell, Mancoridis, 2006), metrics to increase cohesion (Kang, Bieman, 1999) and slicing of FDGs – Functional Dependence Graphs (Rodrigues, Barbosa, 2006), tools to detect modularity violations (Wong, et al., 2011). Even with a quantitative flavor, they clearly differ from the Linear Software Models’ approach.

4.3 Future Work

A mathematical characterization of modularity matrix outliers deserves further investigation. This

relates to the broader issue of determining block sizes, after exclusion of outliers, and module refactoring.

A practical issue is to systematically obtain a broad class of simple patterns strictly obeying the Linear-Reducible Model – like the Observer – as advocated for software building blocks.

Efficiency issues concerning modularity matrix generation and reordering (cf. Borndorfer, et al. 1998) for large scale systems will be investigated.

This work has found that small software systems are strictly Linear-Reducible, and some large software systems are bordered Linear-Reducible. This poses a variety of open questions.

The larger systems shown to be bordered Linear-Reducible were developed before the proposal of the Linear Model. It is conceivable, but still unclear, that in view of this model they could be modified in a natural way to comply with the strict Linear-Reducible model. Similarly, future large scale systems developed with awareness of linearity, may show that strictly Linear-Reducibility rather than limited to certain systems, is indeed applicable to a wide variety of software systems.

4.4 Conclusions

Software has been perceived as essentially different from other engineering fields, due to software’s intrinsic variability, reflected in the *soft* prefix. This versatility is often seen as an advantage to be preserved, even though software composition has largely resisted theoretical formalization.

We have found that Linear Software Models can be formulated, without giving up variability. Thus, software systems of disparate size, function and purpose, may have *Linearity* in common.

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