

QUANTIFIED ONTOLOGIES FOR REAL LIFE APPLICATIONS

Lucia Vacariu¹, George Fodor²

¹*Department of Computer Science, Technical University of Cluj Napoca, 26 Baritiu str, Cluj Napoca, Romania*

²*ABB AB Process Automation, Vasteras, Sweden*

Gheorghe Lazea, Octavian Cret

Department of Automation, Department of Computer Science, Technical University of Cluj Napoca, Cluj Napoca, Romania

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Abstract: Industrial applications are using run-time symbolic approaches only when formal methods can assign useful meaning to symbols by computationally inexpensive algorithms. However, most reasoning methods are either computationally prohibitive or may compute indefinitely; thus such methods have limited use in industrial applications. In many practical situations, the uncertain environment in which an “intelligent” control system acts consists of the symbolic space of some other “intelligent” control system, both networked in the same name space. The result of such interaction is to establish relations between heterogeneous vocabularies and reasoning agents, and between symbols and the physical environment in which the connected systems act. This paper introduces and motivates the necessity for on-line quantification of the degree to which symbols in a system have their intended meaning.

1 INTRODUCTION

The theory presented here has relevance for distributed real-time systems such as those used in multi-robot applications or in distributed manufacturing industries. These systems have a large set of symbols in the form of names for components, signals, process states or configuration parameters. Heterogeneity manifests by units being of different specialization and of different make. Having the right meaning of each symbol is essential for a correct operation of the system. The costs for matching all signals, communication protocols and sub-products during a tender process for a complex system turn out to be a significant part of the total cost of the system. Moreover, after delivery, ensuring that the final system behaves according to specifications can be a lengthy and highly qualified process. The solution to this problem is to establish an ontology for the given industrial domain. These tools need to reduce symbol complexity by automatic information processing, such as via semantic web and ontological languages.

It is an uncommon situation today that such configuration tools work across dissimilar firms or

markets, though many core technologies and standards are available. On a theoretical level, the operations needed, such as ontology merging, alignment composition, union and intersection are still under research (Furst, 2008).

We stress here that formal design verification cannot replace the ontological compliance presented in this paper: even a perfectly designed system that is formally proven to follow a design might encounter a complex environment that does not follow the assumptions in the specification.

Seen as software architecture, ontologies are implemented at the current level of technology as services. These can be organized as local services in each unit or as a combination of a hierarchic set of services – local and specialized - with indirections provided by name servers. Without specifying details, we call in this paper a generic ontology service as the “Industrial Ontology Server” (IOS) (Figure 1).

Practically, an IOS should be able to infer the structure of any type of distributed industrial application. Of course, this is a very ambitious claim, well beyond the forefront of what is available today in academia or industrial research institutes.

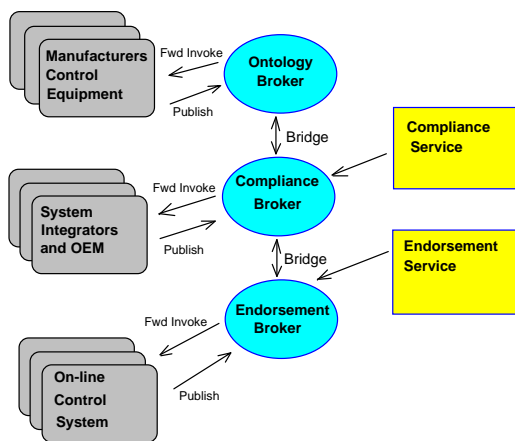


Figure 1: Federated Service Architecture IOS.

The solution proposed here is to use a minimal, uniform ontology associated to products and that each actual systems to update automatically the required operational and dynamic information into IOS-es. With this approach, bidding costs would be substantially reduced and market participation would guarantee more objectivity and flexibility.

Clearly a multi-agent system architecture could implement such requirements. Tools such as JADE (Java Agent DEvelopment framework) with Protégé (Ontology editor and knowledge-base framework) could handle the design of systems based on a common design ontology (Tomaiuolo et al., 2005).

2 SOME RELEVANT TECHNOLOGIES AND THEORIES

An IOS is relevant if it considers all levels of an industrial unit included in a system: from signals and actuators, up to overall goals, including states, alarms, resource allocation, synchronization with other units, etc.

2.1 Ontologies in Process Control

Ontologies for automation and process control applications have several specific layers. At the *design phase*, there is available a formal or informal description of the domain and of the constraints of the design (*Design Model*). The standard for IEC61499 prescribes the Engineering Support System (ESS) that can perform certain syntactic and semantic verifications. Valuable research is being conducted for improving ESS tools (Thamboulidis, Koumoutsos, and Doukas, 2007).

The *framework* (middleware) in which the program, agents or components execute has own ontology and semantics that limits what the application program can execute (*Execution Model*).

Application Ontology. This is the ontology that effectively decides on goals and actuation. The domain of discourse is not necessarily the same as the domain used for the design ontology.

Visualized Ontology. The ontology is typically visualized on a human-readable interface. The domain of discourse, taxonomy among objects and object properties are represented graphically. Automatic generation of visualization using as input ontologies expressed in XML/RDF would be an important advance in technology.

Communicated Ontology. From the execution model and design data, system designers extract an ontology used for communicating among cooperating systems. This ontology may not have the domain of the Design Model, nor of the Execution Model.

Ontology does not enter in a formal, verified way in the design of large control products, as tools are not mature enough. Another conclusion is that a system has several ontologies that should be aligned. There are no commercial tools that can align ontologies. Interesting research results are reported using category theory (Zimmermann et al., 2006).

3 METRICS FOR ON-LINE ONTOLOGIES

3.1 The Decision-Control Space

Essential for taxonomy of process control systems is the type of actuators used. Actuators are performing changes in the real world; their semantics is determined by physical laws.

Independently of the software architecture type, a control system has two essential parts: (a) a decision (information) level and (b) a physical, energy-related level of actuator and plant changes. All the relevant information from sensed signals used for decision forms a hyperspace with each coordinate being one kind of decision information. Let this space be H with N dimensions, $H \in R^N$. Chains of decisions generate chains of action trajectories $T_i (i = 1, \dots, M)$ in this space.

Trajectories may not be continuous as disturbances create ‘jumps’ from one possible trajectory to another. A valid place on this trajectory is often not a point but a hyper-sphere (a topology) since usually control decisions are taken within intervals and not on discrete points.

In traditional control, the space of all relevant signals is called a *state space* and the trajectory a *goal path*. Each point reached during control in the state space is a *state*. For each state space, a controller (or agent or component) has mapped a decision procedure that result in some action being taken. We are not concerned here with what kind of decisions or reasoning a controller is doing, but only with the mapping between actions $u_{i,j}$ and spheres $h_{i,j}$. Here the variable u denotes an action, and the variable h denotes a sphere in H . We call the pair $S_{i,j} = (h_{i,j}, u_{i,j})$ as a state and use first index to denote the sphere i and the second index to denote a goal path j .

3.2 State Semantics

The *active state* is the one currently materialized by sensors. Even if the system consists of many agents or components, there is always a unique active state since the hyperspace covers the whole possible space. However, there may be multiple actions corresponding to each active state. We represent here all actions for one point in the hyperspace as a single action. Some typical state transitions we are interested in are the following: case (1) - normal control with no disturbances, case (2) - control with disturbances and case (3) - lockout.

Performing no action may be a legal, correct operation of the controller, however if inaction is due to decision lockout, then this case is distinct and should be detected.

We seek here moreover to quantify the level of true semantics states have.

3.3 Quantifying State Semantics

3.3.1 Intra-state Distance

The degree for how ‘‘strong’’ is a state with a current place $h_{i,j}$ in the space H is the inverse of the distance from $h_{i,j}$ to the centre of the sphere intended for that state. Therefore, closer $h_{i,j}$ is to the state sphere boundary, less correlation it has with the current state and its action. This is the typical situation when the true state is somewhere in between two states, none fully reached; fuzzy logic can quantify and correct this situation (Grantner and Fodor, 2002). For a measured probability distribution Q , the

Kullback Leibler divergence of Q from P is:

$$D_{KL}(P \parallel Q) = \sum_i P(i) \log \frac{P(i)}{Q(i)} \quad (1)$$

3.3.2 Inter-state Endorsement

The following levels of endorsement for a state are defined to characterize how well predictions are built into the semantics of a state materialize.

Void-Endorsement. A state is void-endorsed if it is materialized. That means for a state $S_{i,j} = (h_{i,j}, u_{i,j})$

$$ve(S_{i,j}) \text{ iff } (h_{i,j} \wedge u_{i,j}) \quad (2)$$

This is the simplest form of endorsement, but it tells an important think: that the program semantics about the environment matches at least once a real instance of the environment. The properties matched are those in $h_{i,j}$.

Weak State Endorsement. A state is weakly endorsed if a consecutive state of a void-endorsed state placed on the same goal path is weakly endorsed at the next instance of time.

$$we(S_{i,k}) \text{ iff } ve(S_{i,k}) \wedge \mathbf{O} ve(S_{j,k}) \quad (3)$$

Here ‘ \mathbf{O} ’ is the ‘next time’ logical operator; both states are on the same path T_k as the second index k shows. A $we()$ state is not a goal state. Weak endorsement means that if a state has materialized and the controller has executed an action at that state, then the expected outcome really turned out to be true in the environment.

Strong State Endorsement. A state is strongly endorsed if a consecutive state on the same path materializes and both states are weakly endorsed.

$$se(S_{i,k}) \text{ iff } we(S_{i,k}) \wedge \mathbf{O} we(S_{j,k}) \quad (4)$$

A $se()$ state is again not a goal state. This state endorsement tells that after a state materialize and the action executed, a next expected state indeed materializes as well and moreover the action from that second state has the expected effects.

3.3.3 Goal Path Endorsement

Goal paths are endorsed in similar way as states.

Void Endorsed Goal Path. A goal path T_k is void endorsed if there exists a state that is not the goal state of the path and which is weakly endorsed:

$$\text{vpe}(T_k) \text{ iff } \exists i \text{ ve}(S_{i,k}) \quad (5)$$

A void-endorsed goal path has some $h_{i,j}$ of some state that materializes in the environment, moreover the corresponding action is being executed, but there is no evidence that any of the following expected states on the same goal path have been materialized. Note that there may be states that are not on any goal path, so a void endorsed state may not necessarily mean a void endorsed goal path.

Weakly Endorsed Goal Path. A goal path T_k is weakly endorsed if there exists some state on the goal path that is weakly endorsed and which is not the goal state of the path.

$$\text{wpe}(T_k) \text{ iff } \exists i \text{ we}(S_{i,k}) \quad (6)$$

A weakly endorsed goal path has at least one state that when acting on the path, get expected effects on the same goal path. However, it is not sure that the expected state has the required quality that even its action will get expected results and thus the semantics of the second reached states is not entirely sure.

Strongly Endorsed Goal Path. A goal path T_k is strongly endorsed if there exists a state that is strongly endorsed on the goal path.

$$\text{spe}(T_k) \text{ iff } \exists i \text{ se}(S_{i,k}) \quad (7)$$

More generally, a goal path is n-strongly endorsed if there are n states which are strongly endorsed on the path. N-strong endorsement tells that many states on the goal path are semantically right, but there may be disturbances that materialize states interleaved with disturbances, on some other goal paths. The condition that one full goal path is traversed without interruption is given by the full-goal path endorsement: a goal path is full-goal endorsed if all the states of the goal path materialize in expected order up to the goal state. Clearly all states of a path that has full-goal endorsement are strongly endorsed, except the goal state and the state immediately before the goal state that is weakly endorsed.

3.3.4 Global Semantic Norms

Many types of norms can be conceived to quantify the level of true semantics using the endorsements given above. For example if $|h_{i,j}|$ is a normalized distance from the center of a state hyper-sphere to $h_{i,j}$ so that $|h_{i,j}| \leq 1$ and the norm $|\text{se}(S_{i,k})|$ gives the

number of states on the current goal path from the state i to the goal state, then a measure of the semantics of the current goal path, SM, is:

$$\text{SM}(T_i) = |h_{i,j}| + |\text{se}(S_{i,k})| \quad (8)$$

SM is a continuous, real valued function that shows how much of the current goal path has been completed.

4 CONCLUSIONS

Complex systems such as mobile robots systems, or distributed industrial control systems need to communicate and use ontological information about their environments and about the tasks they perform. Symbolic operations using formal methods are as yet prohibitive due to computational reasons while manual work raises substantially the costs of such systems. This paper presents a method that combines ontological operations defined formally with automatic updates for control ontology based on on-line direct sensory and actuation data.

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