

UNCOUPLED PARALLEL VIEW DEPENDANT LEVEL OF DETAIL RENDERING OF BINTREE TRIANGULATIONS

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Abstract: In this paper we present an uncoupled parallel technique for view dependant continuous level of detail rendering of regular height field terrain. To this end, the terrain is globally modelled by a simple bintree patch-based representation of right-triangles, which is adaptively divided into uncoupled meshes for parallel processing. This is easily performed by uncoupling the mesh along the observer line. Each uncoupled mesh is then recursively approximated to the corresponding level of detail in the terrain. Cracks are avoided by constraining the refinement levels at the boundaries of adjacent meshes. The level of detail is created on-the-fly with a low amount of CPU overhead, allowing a good representation of the terrain and high frames per second performance. This implementation shows significant improvements in CPU load and frames per second performance over the serial method when executed on machines with multiple processors.

1 INTRODUCTION

Several approaches for multiresolution representation, adaptive modelling, level of detail control, and real time rendering of terrain data have been proposed and studied in the last 10 years. Pioneer work on general mesh simplification and multiresolution modelling are discussed in (Heckbert and Garland 1997; Hope 1996, Hope 1998). In particular for the multiresolution modelling of regular height-field data, methods based on right-triangle triangulations have been developed and widely used. The most important of these methods are quadtree-based and bintree-based representations methods (Duchaineau et al, 1997; Pajarola 2002). These methods are reported to provide a more compact representation of the terrain, better spatial access, faster level of detail (LOD) triangulation and rendering and are easier to implement than more general methods. (Pajarola 2002).

Recently (Holst and Schumann 2006) combine the use of a reduced multiresolution hierarchy based in (De Florianí et al 1997) together with triangle strips for patches. They do not use right bintree triangle representation, nor perform any parallel work.

As a basic tool of this research we use a right-triangle multiresolution bintree algorithm as

discussed in (Duchaineau et al, 1997), which is a special case of the longest edge refinement algorithms for general triangulations discussed in (Rivara 1984, Rivara 1997).

Essentially, the bintree multiresolution method works as follows: (1) Each right triangle in the bintree structure is splitted by its longest edge producing right and left right-triangle children in the bintree; (2) The adaptive local splitting of every target triangle is performed by using a sequence of simple mesh operations over couples of triangles of the same bintree level, which share a longest edge in the mesh; (3) The local splitting of a general target triangle having a longest-edge neighbour of a different level requires splitting propagation, which is a particular case of the longest edge propagation path discussed in (Rivara 1997).

When rendering terrain a specific problem arises. The horizon of visibility is much bigger than in other real time rendering with fixed scenes. If we were to render every triangle at a fixed resolution we would either have a very low quality terrain or a low performance with a high quality terrain. To solve this particular problem the amount of triangles close to the observer has to be maximized and be reduced as the distance increases. This is known as continuous level of detail and the problem has been

addressed by some authors with different techniques. The core of the problem is to find a mesh for each frame that will realistically represent the terrain. The main problem of Continuous Level of Detail is the huge amount of CPU overhead it produces. The ROAM (Real Time Optimally Adapting Meshes) algorithm (Duchaineau et al, 1997), which will be briefly explained on this paper, is a View Dependant Level of Detail algorithm, which adapts a mesh to an optimal triangulation for the camera point of view.

In order to improve the adaptive terrain rendering performance we propose a parallel uncoupled method similar to the one discussed in (Rivara et al, 2006) which takes advantage of the multiprocessor architecture of current computers. Traditional VLOD (Visual dependant Level of Detail) algorithms run all this CPU work through only one processor and therefore make the CPU overhead a big problem. This can be reduced by mesh subdivision and parallel tessellation (Padrón et al, 2005). Most algorithms present fairly complicated solutions for this. In this paper we consider simple parallelization of real time view dependant continuous level of detail.

2 BASIC TECHNIQUES USED IN THIS PAPER

2.1 View Dependant Level of Detail Algorithms

VLOD algorithms are designed to interactively perform view-dependent locally adaptive terrain meshing. They rely on a multi-resolution terrain representation that is used to build the adaptive terrain representation of a frame. To accomplish this, the algorithm used in this paper is a basic implementation of ROAM. VLOD algorithms take the complete terrain data usually, in form of a height field, and generate a mesh that includes every vertex of the height map near the observer and discard vertices as they generate mesh sections further away from the observer. ROAM accomplishes this by generating a mesh for each frame, while other algorithms, like GeoMipMaps (H. de Boer, 2000) do this by pre-calculating a certain amount of meshes. The advantage of dynamically generating the mesh for each frame is that the representation is much better than on pre-calculated meshes. The disadvantage is that they require a lot more processing than pre-calculated meshes do. To make up for this a simple parallelization mechanism is shown in this paper.

2.2 Mesh Representation

Multiresolution representations based on right-triangle triangulations are the most suitable for adaptive terrain modelling of regular height field data, where a multi-resolution surface is stored in a data structure that can be recursively refined on demand. Among these we can mention the quadtrees and bintree methods (Lario, Pajarola and Tirado, 2003). Quad-Trees have the disadvantage though, that they are inherently based on recursive quads which can have up to four children. For our purpose a structure that is inherently based on triangles is much better suited given the refinement technique described later. In a Binary Triangle Tree, each of two possible children are triangles formed by dividing the parent triangle in two by longest edge bisection as illustrated in Figure 1. For a discussion and evaluation of the bintree multiresolution representation see (Duchaineau et al, 1997; Pajarola 2002).

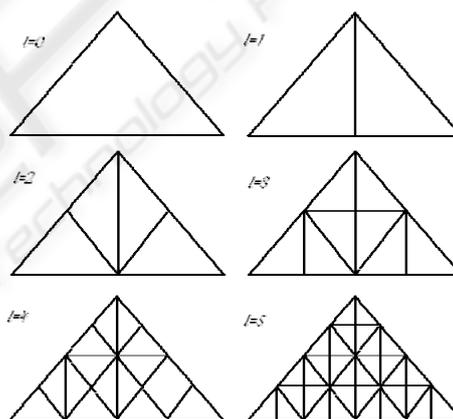


Figure 1: First 6 levels of a Binary Triangle Tree.

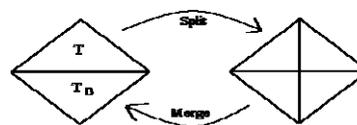


Figure 2: Split and Merge operations.

3 NON-PARALLEL BINTREE PATH-BASED METHOD

The serial method chosen to represent terrain is a simple version of ROAM. The algorithm can be described as follows:

Serial Tessellation Algorithm:

Input: A high resolution height map terrain data.

- a) The terrain data is divided into N square patches.
- b) Each terrain patch is initialized with two bintree triangle structures at the coarsest level (every patch contains two triangles). The height coordinate of each triangle vertex is given by the corresponding value in the height map.
- c) All triangles are linked together by pointers to their three neighbours.

For each frame adaptively do:

- d) The landscape is initialized and the adaptive view dependant level of detail tessellation process is performed.
 - e) The frame is rendered.
- For ends.

The terrain is subdivided into bintree patches in order to keep the depth of the tree structures controlled.

During the tessellation process (Duchaineau et al, 1997) the mesh is dynamically calculated for each frame. This is done by recursively traversing the two BTT structures in each visible patch and refining it until the desired level of detail is reached. This is done with two basic mesh operations called merge and split (see figure 2).

When two adjacent triangles are both from the same level we refer to it as a diamond. Split replaces a triangle T with its children $T0$ and $T1$ and does so with the adjacent triangle as well. This introduces a new vertex at the centre of the diamond, resulting in a new continuous triangulation.

This tessellation process is by far the most CPU expensive part of the algorithm. During this process, each visible patch is visited and their BTT structures recursively traversed and refined. A triangle T in a triangulation cannot be split immediately when its base neighbour Tb is of a coarser level. It would produce a crack in the mesh, so we have to force T to split. To force T to split, Tb must be forced to split first, which may require other triangles to split first. As long as we recursively split all required triangles, the mesh will remain consistent. This is also known as the Longest Edge Propagation Path or LEPP (Rivara, 1997), which finishes finding a terminal edge, which is the base edge for two adjacent triangles or an edge of the border of the

mesh. Only then we start the splitting process (see figure 3).

If we go on refining too much, we can eventually run out of memory. To avoid this, a pool of possible vertices to be allocated is defined. Once a new vertex is introduced, it is eliminated from the pool and becomes part of the triangulation. Once we run out of vertices in pool, no further refinement is possible.

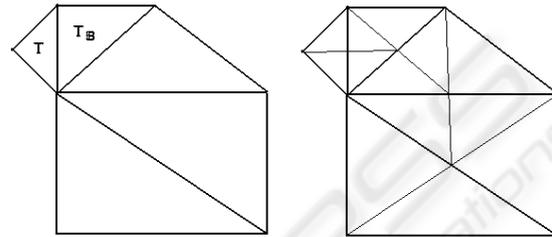


Figure 3: In order to split T all the triangles along the longest edge path have to be split first.

Since we want different levels of refinement on different areas of the mesh, we need to decide how far to refine the mesh in a given point. This is done by introducing an adaptive tolerance parameter t which depends on the distance of the given point to the observer.

Each time that a triangle in the BTT is analyzed we compute t for the given position and compare its value with the distance between the middle point of the triangles base and the corresponding point in the Height Map. If the distance is bigger than our calculated t we split the triangle. In this way we achieve our goal of continuous level of detail.

4 PARALLEL TESSELLATION

With multiple processor technology the tessellation could be done in a fraction the time if the meshing problem is decoupled. The problem that arises is that given the recursive nature of the process, the CPUs might want to allocate the same vertex at the same time, causing access violations and ending the process. The first approach to solve this problem is to work on separate meshes and thereby distributing the process completely. This causes a problem of continuity at the bounding of adjacent meshes. For illustration purposes, let us focus on dual processor scenario. As both meshes are refined independently, we have no guarantee the union is going to be continuous and therefore cracks may appear at the union.

Other algorithms solve the problem by concurrent vertex insertion (Chernikov and Chrisochoides, 2004), scheduling the point insertion in such a way that no access violations will occur during the refinement. The problem could be solved by using a different mesh structure but we want to keep the efficiency of the BTT structure. The solution implemented in this paper works by decoupling the mesh, and working on it as if it were two separate meshes, but maintaining consistency. This way we do transform the process into a distributed problem but with shared memory resources.

The solution implemented parallel tessellation algorithm works as follows:

Input: a high resolution height map terrain data. The terrain data is divided in N square patches. Each terrain patch is initialized with two bintree triangle structures.

While the observer traverses the terrain do
 For each frame (observer position) do
 - the mesh is partitioned along the observer line L and common parameters t are calculated along L .
 - each submesh and associated parameters t are assigned to each processor for tessellation.
 - each submesh is adaptively calculated in parallel using bintree triangle structures until the desired level of detail is achieved.

For ends.

While ends.

The reason to partition the mesh along the line of observation is to balance processor load. This can be easily done by identifying the patches at each side of the observer and accessing their BTT structures.

At the initialization stage all triangles of the BTT structures would be linked together in the serial algorithm, instead each half of the mesh is linked together, leaving an unconnected division of it running through the middle of the mesh (see figure 4). This way we assure that it is a decoupled problem because the edges in the middle of the mesh become terminal edges, where the refinement stops. More on decoupled parallel refinement can be found in (Rivara et al, 2006).

After splitting the mesh two separate pools for vertex allocation are defined. This way both halves can be refined independently without sharing resources.

The LOD is computed by two tessellation threads. These can be assigned to separate CPUs. The tessellation process is the same as in the non-parallel case using the method described above.

Each of the tessellation threads has its own pool of vertices and set of patches containing the BTT structures, the only shared resource is the height

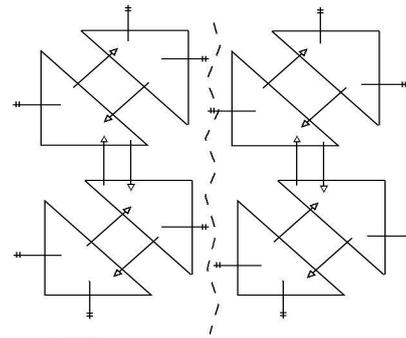


Figure 4: The initial mesh is divided. The triangles running along the middle are uncoupled.

map. If we would duplicate the height map and include a copy of the height map to each thread, the process would be completely distributed and therefore have no shared memory resources. This could be useful to run the algorithm on a cluster of machines.

For the parallel processor problem, the access to the height map does not present any problems, because each thread is dealing with a different part of the terrain and therefore there is no possibility that both threads would try to access the same value.

Once the tessellation process is complete, and before the rendering step can start, the two halves have to be consistent with each other to avoid cracks at the initial division line of the mesh. At this point we have no guarantee that both parts of the meshes fit together without leaving cracks at the union. To ensure consistency we force the tolerance factor t at both adjacent borders to be exactly the same. This means the same level of detail will be used on both halves and the condition to stop refining is the same at both sides of the union as long as we had enough vertices in the pool. This produces a mesh with duplicate vertex at the union and thus producing a few extra triangles, but with no cracks in the mesh. We are also over-refining the neighbouring mesh. This causes a few extra triangles to appear at the coarser side of the mesh. The triangle overhead is a small price to pay, given the fact we have done the tessellation in almost half the time.

5 RESULTS

To test the performance of the algorithm, a parallel version and a non-parallel version of the same algorithm were implemented and compared in terms of CPU load, memory usage and frames per second performance.

The testing method used can be described as follows. A 230 second long flight over the terrain

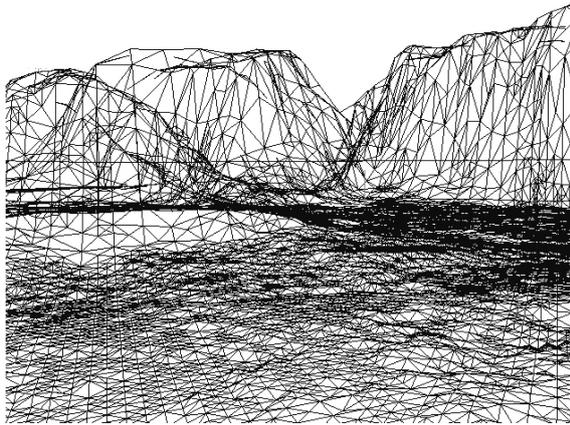


Figure 5: Terrain mesh.

was programmed and executed ten times. Each of the flights was done rendering the scene at different levels of detail, starting with 180000 vertices and ending at 5000. Each flight was performed with the serial and the parallel algorithm and CPU load and frames per second were measured for both algorithms. The number of desired vertices per frame, maximum allowed time and circuit were given. Spatial steps were calculated in terms of the elapsed time, total time allowed and position. The test platform was a dual Turion64 processor computer, running at 1.6 GHz with 1 GB of RAM and an Nvidia GeForce 6100 graphics card running Windows XP.

The results are illustrated in the following table:

Table 1: Performance of sequential and parallel algorithm.

Vertices per Frame	Average Sequential frames/sec	Average Parallel frames/sec	% Increase in Performance
180000	17.03	23.78	39.64
160000	17.42	25.31	45.29
140000	20.22	33.10	63.70
120000	21.05	35.71	69.64
100000	23.61	39.85	68.78
80000	28.60	44.57	55.84
60000	39.42	57.35	45.48
40000	52.19	68.16	30.60
20000	71.04	73.06	2.84
5000	81.03	81.18	0.19

The same experiment executed over different terrain data yielded similar results.

As shown in figures 6 to 8, the average CPU load was well balanced. The load balance was not exactly the same due to differences on the terrain geography on each side of the mesh division.



Figure 6: CPU load at 180 000 vertices per frame.

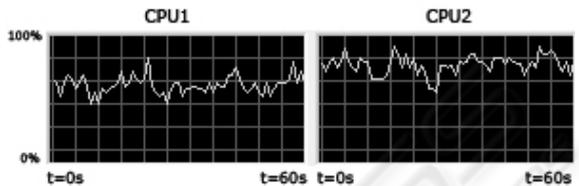


Figure 7: CPU load at 100 000 vertices per frame.

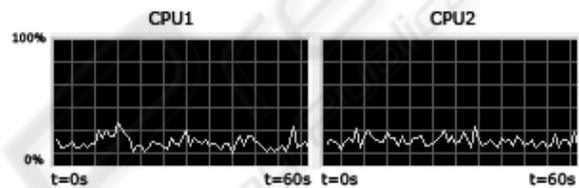


Figure 8: CPU load at 20 000 vertices per frame.

Although one might be tempted to think that the performance should increase by a bigger factor, we need to have in mind the fact that all we have done is parallelizing the tessellation process, but other processes like computing texture coordinates are done sequentially. Given this, the increase in frames per second performance shown in table 1 is considerable. As shown in table 1, the increase in performance was at its maximum between 100 000 and 160 000 vertices per frame. When decreasing the amount of vertices, the parallel algorithm shows smaller improvement over the sequential algorithm. This happens because at low workloads one CPU is able to manage all of the work, improvement starts to be noticeable at 40 000 vertices. On the other end, with over 160 000 vertices the improvement of the parallel algorithm over the sequential one starts decreasing until it stays at around 39 %.

The memory usage showed no considerable change between the serial and the parallel algorithm. In both algorithms the entire structure could be stored in the same amount of memory. This was to be expected as the mesh size did not change and no extra storage space was needed. The structure took about 1.2 MB of storage space at low vertex rates. This increased to a maximum of 17.2 MB at 180 000 vertices per frame.

6 CONCLUSIONS AND DISCUSSION

We have presented a parallel continuous level of detail technique for rendering bintree based structures. The mayor improvements in performance are achieved by turning the problem into an uncoupled refinement process, which allows a terrain mesh to be generated on multiple processors and thereby increasing the performance in terms of frames per second considerably.

The issue of keeping the load balanced on more than two processors posts some difficulty because the mesh partitioning strategy needs to be generalized. Due to the decreasing level of detail towards the observers horizon, and the changes in complexity of the mesh, the next partition strategy is much harder to determine and will be part of future research.

This paper focused on showing the increase in performance by parallel mesh tessellation, but many improvements can still be implemented, like node caching, triangle fan generation, priority queues and occlusion culling among other techniques.

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