COINCIDENCE BASED MAPPING EXTRACTION WITH GENETIC ALGORITHMS

Vahed Qazvinian, Hassan Abolhassani and S. Hossein Haeri Web Intelligence Lab, Computer Engineering Department Sharif University of Technology, Tehran, Iran

Keywords: Ontology Matching, Ontology Aligning, Coincidence Based Matching, Alignment Extraction, Genetic Algorithm.

Abstract: Ontology Aligning is an answer to the problem of handling heterogenous information on different domains. After application of some measures, one reaches a set of similarity values. The final goal is to extract mappings. Our contribution is to introduce a new genetic algorithm (GA) based extraction method. The GA, employs a structured based weighting model, named "coincidence based model", as its fitness function. In the first part of the paper, some preliminaries and notations are given and then we introduce the coincidence based weighting. In the second part the paper discusses the details of the devised GA with the evaluation results for a sample dataset.

1 INTRODUCTION

Semantic web is made up of distributed information that has resulted in designing ontologies to lessen the heterogeneity. Yet another problem exists: ontologies themselves may cause heterogeneity. This is when two ontologies are trying to express same knowledge or concepts but they use different languages or words (Euzenat and Valtchev, 2004). This has led to the problem of ontology aligning.

Ontology aligning has been discussed for a while, to enable a mapping of two heterogenous information in different domains. This will make agents which are using different ontologies, establish interpretability. This alignment will give a correspondence between concepts and semantics of two ontologies.

To do an alignment it is customary to first apply some measures (simple or complex) to reach to some initial guesses. Then the problem is how to form an ideal mapping. This problem which is referred to as Mapping Extraction is the target of this paper.

After having an explanation of related works in section 2, section three shows some definitions and notations used in the paper. Then section 4 discusses our coincidence based theory which is the basis for our GA algorithm discussed in section 5. Section 6 explains evaluations done on the algorithm and finally section seven concludes the paper.

2 RELATED WORKS

Unfortunately works on ontology extractions are not so many, as stated in (Bouquet et al., 2004). In (Melnik et al., 2002), to extract a reasonable extraction, *Stable Marriage* (Gibbons, 1985) problem is discussed. There are some other approaches, e.g. a machine learning approach to the problem is discussed in (Doan et al., 2003), and (Mitra et al., 2003) describe a probabilistic based model. Some methods tend to a trade off of different features such as efficiency and quality, as in QOM (Ehrig and Staab, 2003) and some have used approaches to integrate various similarity methods (Ehrig and Sure, 2004).

(Kalfoglou and Schorlemmer, 2003) have come with a comprehensive review and presentations on the methods and approaches and the state of art in ontology aligning.

According to (Haeri et al., 2006): "no [much] work is so far done on the problem of Ontology Alignment or Ontology Matching in which graph theoretic backbone of problem is scrutinized.". With the use

176 Qazvinian V., Abolhassani H. and Hossein Haeri S. (2007).

COINCIDENCE BASED MAPPING EXTRACTION WITH GENETIC ALGORITHMS.

In Proceedings of the Third International Conference on Web Information Systems and Technologies - Web Interfaces and Applications, pages 176-183 DOI: 10.5220/0001271901760183

Copyright © SciTePress

of graph theory and such a modeling we believe that there is a vast area for new work on the problem of ontology aligning.

3 DEFINITIONS

In this section we will define the notations that are used throughout the paper. We will start by defining necessary mathematical backgrounds.

A graph G_i , by definition, consists of two sets: $V(G_i), E(G_i)$, where $V(G_i)$ is the set of vertices, and $E(G_i)$ is the set of edges. The size of a graph is $|V(G_i)|$, which is denoted by |G|. Lets assume that labels assigned to nodes are chosen from a finite alphabet Σ . Let $\lambda \notin \Sigma$ be a null character, and $\Sigma_{\lambda} = \Sigma \cup \lambda$.

3.1 Typed Graph

An ontology O_i , in this paper, is considered as a *typed* graph G_i . A *typed* graph, as defined in (Haeri et al., 2006), is denoted by $G_i(V, E, T)$, where E is of type: $E : V \times V \rightarrow T$. labels in T are from Σ_{λ} . In such a graph an edge e between two vertices v_{i_j}, v_{i_k} with type t is denoted by: $e(v_{i_1}, v_{i_k}) : t$.

In this paper each ontology O_i is modeled using a typed graph G_i where concepts in O_i are nodes of G_i and relations and properties of O_i are *typed* edges of the graph.

3.2 Distance

The distance of two concepts belonging to two different graphs is described as the distance of their labels in a metric space. Usually this metric distance is described by a *distance function*, $\delta : (\Sigma_{\lambda} \times \Sigma_{\lambda}) \setminus (\lambda, \lambda) \to R$. So the distance of two nodes v_i, v_j is denoted by $\delta(label(v_i), label(v_j))$. For simplicity we will show it by $\delta(v_i, v_j)$. This metric distance, δ , for any $v_1, v_2, v_3 \in \Sigma_{\lambda}$ should have the following properties (Haeri et al., 2006):



- 2. $\delta(v_1, v_2) = \delta(v_2, v_1)$ Symmetry
- 3. $\delta(v_1, v_2) + \delta(v_2, v_3) \ge \delta(v_1, v_3)$ Transitivity

This distance function can either be a *String-Based* distance or any other possible one. In this research, the distance of two vertices v_i, v_j , say $\delta(v_i, v_j)$, is considered to be the *Levenshtein Distance* (Levenshtein, 1966) of their labels.



Figure 1: A sample matching of two graphs G', G.

3.3 Ontology Alignment

In this section we will discuss our own understanding of a one to one matching between ontologies.

A one to one matching¹ of two ontologies O_{i_1}, O_{i_2} is denoted by $m : O_{i_1} \to O_{i_2}$ and is a one to one correspondence between nodes of the two graphs of O_{i_1}, O_{i_2} : (G_{i_1}, G_{i_2}) .

(Ehrig and Sure, 2004) define the mapping function the following way:

- $m: O_{i_1} \to O_{i_2}$
- $\forall v \in G_{i_1} : m(v) = v' \text{ if } v' \in G_{i_2} \text{ and } \delta(v, v') < t$, for t being a threshold

v' is the *corresponding node* of v under the mapping m. It is clear that with this definition, correspondence of edges, is determined by the correspondence of nodes:

 $\forall e(v_{1_j}, v_{1_k}) : t \in E(G_{i_1}) : \text{if } m(v_{1_j}) = v_{2_j} \neq \\ \emptyset, m(v_{1_k}) = v_{2_k} \neq \emptyset, \text{and } e(v_{2_j}, v_{2_k}) : t \in \\ E(G_{i_2}) \text{ then } m(e(v_{1_j}, v_{1_k})) = e(v_{2_j}, v_{2_k})$

Figure 1 illustrates a sample alignment for two sample ontologies G', G.

3.4 Edge Preservation

We will call an edge $e(v_{1_j}, v_{1_k}) : t \in E(G_{i_1})$ is preserved under the matching *m*, iff there is an edge $e(m(v_{1_j}), m(v_{1_k})) : t \in E(G_{i_2})$. In other words an edge, *e* is preserved under matching *m* if and only if $\exists e' : e' = (m(v_{1_j}), m(v_{1_k})), m(e) = e'$, and is not preserved otherwise

The preservation of edges between corresponding nodes is the key point to find an ideal matching. In fact in an ideal alignment most of the edges of one ontology are preserved in the second one.

¹We will sometimes call it Alignment or mapping of two ontologies

4 COINCIDENCE BASED WEIGHTING

In this section we introduce and discuss a new weighting model for an alignment, with which we will later design our genetic algorithm. The coincidence based alignment weight function is sufficiently discussed in (Haeri et al., 2006), and here, we will have a brief introduction to it. Before talking about the weight itself, lets take some time, and discuss the matter.

Consider a mapping *m*, between two ontologies with graphs G_{i_1}, G_{i_2} , and also consider two nodes $v_{1_j}, v_{1_k} \in V(G_{i_1})$ and their matches $m(v_{1_j}), m(v_{1_k})$. The weighting system should result a high weight if v_{1_j} is close to $m(v_{1_j})$ and also is v_{1_k} to $m(v_{1_k})$ and besides, $e = (v_{1_j}, v_{1_k}) : t \in E(G_{i_1})$ is preserved under *m*. In this case v_{1_j}, v_{1_k} are close to $m(v_{1_j}), m(v_{1_k})$, and there is an edge both between v_{1_j}, v_{1_k} , and between $m(v_{1_j}), m(v_{1_k})$. This case is considered to be the most desired one and should be given the highest value.

In the second case lets suppose the edge is not preserved. Here, a negligible negative point should be given. The reason for negative point is the fact that, the edge is not preserved and the structural matching of the graphs is interrupted. In this case the nodes are very close but the edge is missing.

The farther any of the nodes is, from its match, the lower should be the positive value of the matching. If the edge is preserved, we give this matching a low positive value. But when the edge is not preserved, in fact it is an undesired matching. So we give it a negative point. In this case not only the nodes are far from their matches, but also the edge is not preserved.

According to above considerations there should be six different categories: (Suppose G, G' are graphs of two ontologies O, O' to be aligned. a, b are concepts from G, and a', b' from G')

- **Category I.** *a* and *a'* are too close², and *b, b'* are close as well. That means, the distance between *a*, *a'* is low, and so is the distance between *b, b'*. The edge between *a, b* is preserved so this category is of much importance. This is because actually the two edges coincide too much.
- **Category II.** In this category the edge is preserved but only one of the *a* or *b* is close to its match. This is good but not as much as the previous category.
- **Category III.** The two peers of an edge are close to their matches, that means, *a* is close to *a'* and *b* is close to *b'*. But the edge is not preserved. This category should not be penalized much, be-

cause at least concepts are close to their matches and vertices coincide.

- **Category IV.** The edge between *a*, *b* is not preserved, and *b* if far from *b'*. The only positive point of such a matching is the fact that *a* and *a'* are close.
- Category V. and VI. in these categories, both *a*, *a*' and *b*, *b*' are far from one another, and the difference is in the preservation of edges. Both cases are not desired and should obtain low points.

According to the above sort and discussions, the following weigh function is suggested:

$$w(m) = w_0(m) - w_l(m) - w_r(m)$$

$$w_0(m) = \sum_{(v_1, v_2): t \in E(G), (m(v_1), m(v_2)): t \in E(G')} f(v_1) + f(v_2)$$

$$w_l(m) = \sum_{(v_1, v_2): t \in E(G), (m(v_1), m(v_2)): t \notin E(G')} g(v_1) + g(v_2)$$

$$w_r(m) = \sum_{(v_1, v_2): t \notin E(G), (m(v_1), m(v_2)): t \in E(G')} g(v_1) + g(v_2)$$

The functions f and g, referred to as **Normaliza**tion Functions, are in the form:

$$f: R \to R^+$$
$$g: R \to R^+$$

f,g are related to the distance function. In fact, f should be a positive decreasing function, so that if $\delta(v,m(v))$ grows, it decreases to reduce the positive point. And on the other hand g should be a positive increasing function to grow with the growth of $\delta(v,m(v))$ to increase the negative point for that match. Normalization functions are defined by tuning the system. This will be described again later.

5 GENETIC ALGORITHM

This section describes the designed genetic algorithm. Matching two general graphs in polynomial running time algorithms is impossible, because the problem in its general case is MAX SNP-Hard (Arora et al., 1992). So a random search algorithm could be a good idea when designed carefully. This led us to the idea of using genetic algorithm.

²in terms of a distance function described before

5.1 Coding a Matching

To code a matching we used hashmaps (Cormen et al., 2001) in which keys are concepts of one ontology and entries are concepts of another. Entry for each key is actually the corresponding node of that key in the mapping

5.1.1 Pairs

According to the coincidence-based model, we defined *Pair*, as two concepts from one ontology, between which there is a relation (So there is an edge in the graph of that ontology). Figure 1 shows the alignment of two ontologies. $(v_1, v_2), (v_1, v_3), (v_3, v_4), (v_2, v_4)$ are pairs of *G*. A pair also has a weight according to the alignment it involves in.

Clearly speaking, a pair is a function of the form:

$$P: V \times V \times T \longrightarrow R$$

where *V* is the set of vertices in ontology graph *G* and *T* is a set of labels in Σ_{λ} . So an ontology in a matching has a limited number of pairs.

The weight of a pair depends on the alignment in which the ontology is involved. Let G_{i_1}, G_{i_2} be two graphs of two aligned ontology, and $v_{1_j}, v_{1_k} \in V(G_{i_1})$. Also assume an edge between v_{1_j}, v_{1_k} to be of type *t*, $e_{1_{j_k}} = e(v_{1_j}, v_{1_k}) : t$.

 $P(v_{1_j}, v_{1_k}, t)$ in the alignment of two ontologies is given by:

$$P(v_{1_j}, v_{1_k}, t) = \begin{cases} f(v_{1_j}) + f(v_{1_k}) & e_{1_{jk}} \text{ preserved} \\ -(g(v_{1_j}) + g(v_{1_k})) & e_{1_{jk}} \text{ not preserved} \\ -\infty & \text{if } e_{1_{jk}} \notin E(G_{i_1}) \end{cases}$$

For the couple of concepts which do not form a *pair* the value of *P* function is set to be $-\infty$. Definition of pairs, seemed necessary for further crossover function which will try to improve the structural matching.

In the alignment of two ontologies, $O_{i_1}(G_{i_1})$, $O_{i_2}(G_{i_2})$, say $m : O_{i_1} \to O_{i_2}$, we also define the *weight* of a single concept from one ontology, $W(v_{1_j})$ where $v_{1_j} \in V(G_{i_1})$, as follows:

$$if v_{1_j} \in V(G_{i_1}), m(v_{1_j}) \in V(G_{i_2})$$
$$W(v_{1_j}) = \sum_{\forall v \in G_{i_1}: e(v_{1_j}, v): t \in E(G_{i_1})} P(v_{1_j}, v, t)$$

5.1.2 An Example

To make things clear about the definition of *pair* and its corresponding weights described above, we give an example on how to compute these weights. In the Figure 1 we have:

$$\begin{split} P(v_1, v_2, t_2) &= f(v_1) + f(v_2) \\ P(v_1, v_3, t_1) &= f(v_1) + f(v_3) \\ P(v_3, v_4, t_3) &= -g(v_3) - g(v_4) \\ P(v_2, v_4, t_4) &= -g(v_2) - g(v_4) \\ P(v_1, v_4, t_i) &= P(v_2, v_3, t_i) = -\infty \\ W(v_1) &= (f(v_1) + f(v_2)) + (f(v_1) + f(v_3)) \\ W(v_2) &= (f(v_2) + f(v_1)) - (g(v_2) + g(v_4)) \\ W(v_3) &= (f(v_3) + f(v_1)) - (g(v_3) + g(v_4)) \\ W(v_4) &= -(g(v_4) + g(v_3)) - (g(v_4) + g(v_2)) \\ \text{Now, with these definitions, it is the time, to clarify the steps of the genetic algorithm. \end{split}$$

5.2 Initialization

As of any other genetic algorithms, a population is needed. A population is made up of some individuals each of which is a solution to the problem (an alignment in this problem). The start population, is initialized randomly, with an initial size of 1000 individuals. The ideal matching can be reached more quickly if the initial individuals, are made on the basis of the labels of concepts, that is if v_{1_j} in G_{i_1} and v_{2_j} in G_{i_2} have same labels, then let v_{1_j} corresponds to v_{2_j} in the initial mapping.

5.3 Selection

In each iteration, we sorted the 1000 individuals according to their fitness described in section 4 (coincidence based weight function, and we selected the 500 best individuals as parents of next step. From these 500 individuals, with the use of crossover and mutation functions (as we will see later), 1000 new individuals are created. These 1000 individuals are sent to the next iteration.

5.4 Crossover

In the crossover function, single nodes are compared according to their weight. As we described before the weight of a single node in an alignment is the sum of weights of pairs in which, that node is included. Consider two parents between two ontology graphs

Consider two parents between two ontology graphs G_{i_1}, G_{i_2} . To make an offspring from two parents, for every node in G_{i_1} , say v_{1_j} , the mapping with larger $W(v_{1_j})$ is copied to the offspring. if $m(v_{1_j})$ in G_{i_2} is already assigned with some other node of G_{i_1} , then v_{1_i} is put in a reserved list. At the end of the

complete iteration of nodes in G_{i_1} , the nodes in that reserved list are randomly mapped to the unassigned nodes of G_{i_2} . The random assignment is not done in the middle of the iteration, to prevent nodes of G_{i_2} to be assigned to some random nodes, where they can be assigned to better nodes later in the iteration. So this random assignment is postponed until all nodes of G_{i_1} are examined to map to nodes in G_{i_2} .

As an example suppose $v_{1_m} \in V(G_{i_1})$ should map to $v_{2_m} \in V(G_{i_2})$ and v_{2_m} is already mapped by some node from G_{i_1} , so if at the time we assign v_{1_m} to some random node like $v_{2_n} \in V(G_{i_2})$. This will prevent a possible good mapping of v_{1_n} to v_{2_n} later in the iteration, because that will make v_{2_n} assigned. So this random matching is delayed until no more assignment is possible. In Figure 2 two mapping between two ontologies O, O' are shown, and we want to decide the match node for $v_i \in V(G)$ in the offspring. In parent 1 we have:

$$W(v_i) = P(v_i, v_j, t_2) + P(v_i, v_k, t_1) = (f(v_i) + f(v_j)) + (f(v_i) + f(v_k))$$

and in parent 2 we have:

 $W(v_i) = P(v_i, v_j, t_2) + P(v_i, v_k, t_1) = -(g(v_i) + g(v_j)) + (f(v_i) + f(v_k))$

f, *g* are positive functions so amount of $W(v_i)$ in parent 1 (Figure 2 (a)) is greater than that of in parent 2 (Figure 2 (b)). So as is seen, the corresponding node of $v_i \in V(G)$ in offspring is chosen by the mapping shown from parent 1, and therefore is $v'_i \in V(G')$

This kind of crossover seems reasonable because the mapping of a single node in the offspring is not worst than that of the two parents. So by this assumption, little by little, mappings of nodes will converge to an ideal ones.

5.5 Mutation

A proportion of the population are mutated with some probability, different in various situations. In mutation of a mapping of two ontologies with graphs G_{i_1}, G_{i_2} , two random nodes from G_{i_1} are chosen, and their matches in G_{i_2} are substituted with each other. Let $v_{1_j}, v_{1_k} \in V(G_{i_1})$ are chosen randomly, also let $m(v_{1_j}) = v_{2_j} \in V(G_{i_2}), m(v_{1_k}) = v_{2_k} \in V(G_{i_2})$. In the mutation process we just substitute the match nodes of the selected ones. So the new mapping will be $m(v_{1_j}) = v_{2_k} \in V(G_{i_2}), m(v_{1_k}) = v_{2_j} \in V(G_{i_2})$.

5.6 Continuation

500 best children from a previous step are sorted decreasingly, and form the parents of current step. These parents produce 1000 new individuals, and the best 500 of these 1000 children are selected as parents



Figure 2: crossover . (a)part of parent 1 matching. (b)part of parent 2 matching (c) part of offspring.

of next step. The i^{th} and $i + 1^{th}$ parents involve in crossover and mutation, and create offsprings. From the set of all these offsprings and their two parents, the two best are chosen as children. The last parent is mixed with the first one to produce last offsprings. As stated before, after the necessary 1000 children are made, they are sorted according to the fitness function and the best 500 are selected for the next step (Figure 3).

5.7 End Condition

To end the iteration of GA, we used a threshold for convergence. The sequential GA continued until the best mapping among all individuals in the population



Figure 3: Population generation in each step of GA.

did not improve for more than 15 steps. Such mapping is reported as the answer, and the alignment of two ontologies is then finalized.

6 EVALUATION

To evaluate our Genetic Algorithm, we designed three kids of experiments. In the first experiment, we tested the Genetic Algorithm with diverse mutation probability. In the second experiment, we tried to align to identical ontologies (actually we aligned one ontology with itself). This experiment helped us examine the efficiency and accuracy of the algorithm, when two ontologies are more similar to each other. To verify our contribution, we come with a third experiment, in which we used a naive local search alignment method.

6.1 Limitations

The coincidence-based ontology matching appears to be an innovative idea, however there are essential limitations with this method. The most important limitation is the available ontologies and test collections. Most of them do not have a large taxonomic structure and so the method does not have enough merit for them. To test this method we needed large taxonomic ontologies. We found "Tourism" ontologies (tou,) a suitable test bench with approximately 340 classes and concepts.

6.2 Various Experiments Characteristics

For the tourism ontologies, an ideal alignment was given, so that we could have the ideal alignment of our graphs. The precision measure (Baeza-Yates and Ribeiro-Neto, 1999) was calculated based on the given information.

• Experiment 1

As discussed previously, in this experiment we aligned "TourismA" with "TourismB". This is the main experiment of our method, to check the efficiency of our coincidence-based genetic algorithm.

In this experiment, from each two parents, we made one offspring with the use of crossover function. From the three different individuals (parents and the offspring) we chose two best of them, to introduce as children of this amalgamation.

Normalization Functions are as follows:

$$f(v) = \frac{1}{e^{\delta(v,m(v))}}$$
$$g(v) = \frac{1}{e^{\max(5,15-\delta(v,m(v)))}}$$

These functions actually satisfy the characteristics expected from f, g (explained in Sec. 4). fis a decreasing function and decreases with the growth of δ and g is increasing. Exponential functions were chosen for f, g so that f, g would have close and comparable values. In fact, these functions match the discussions on positive and negative points for different categories of a coincidence based weight.

- Experiment 1.1

After the 1000 individuals were created, we mutated the **lower half** of the them (with the mutation function described before) with a probability of **0.7**.

- Experiment 1.2

After the 1000 individuals were created, we mutated the **lower half** of the them (with the mutation function described before) with a probability of **0.3**.

- Experiment 1.3

Mutation was done on each one of the 1000 individual in the **all** of the 1000 children with the probability of **0.5**.

• Experiment 2

In this experiment we are aligning "TourismA" with itself. This actually is a verification that shows how efficient the genetic algorithm will work, if two ontologies are more similar and actually more coincident.

The generation summary is similar to the previous experiment, and mutation was done on the **lower**



Figure 4: Precision result of experiments.



Figure 5: Convergence of results in GA.

half of the individuals, with the probability of 0.5.

Normalization Functions are like the previous experiment,

$$f(v) = \frac{1}{e^{\delta(v,m(v))}}$$
$$g(v) = \frac{1}{e^{\max(5,15-\delta(v,m(v)))}}$$

• Experiment 3

This experiment actually provides a baseline comparison of the GA method with a naive local search method. In this part, we implemented a naive hill-climbing local search method. For the start point, we made an initial alignment. In this initial alignment, all concepts in "TourismA" where matched with concepts in "TourismB". For a node v_j in TourismA if there were a node v'_j with label $label(v_j)$ in TourismB, we matched v_j with v'_j . Otherwise we mapped v_j to a random node of TourismB.

After that, in each iteration, the best single change (mutation) was preformed to improve the weight value of alignment. We iterated the method until almost 1000 steps, where, the results did not improve for more than 15 steps.

6.3 Results

Figure 4 shows the result of the above experiments in precision (Baeza-Yates and Ribeiro-Neto, 1999). As it is shown, with equal graphs Genetic algorithm finds the best matching and precision is 1.

In our experiments, the distance threshold, which we talked about in section 3.3, is set to be 4. We chose this number by experience, however there could be

other solutions to determine this number, like machine learning techniques, etc. The discussion on how to define this threshold and the distance function, is beyond the scope of this paper. With other experiments, however, the result is a little inaccurate in comparison with the ideal alignment and precision is approximately 0.8.

We also did an investigation on iteration number and convergence of the result in this genetic algorithm for Experiment 1. The results are shown in figure 5.

7 CONCLUSION AND FUTURE WORK

Genetic Algorithm seems efficient in the problem of ontology alignment extraction. It also converge rapidly, for example after approximately 40 iterations in our experiments. This number can even be reduced by choosing a biased initial population, where labels can be involved to choose better initial mappings.

Coincidence Based approach, when improved and used as a fitness function of a genetic algorithm might be useful when ontologies have a more taxonomic structure. There is also some weakness with genetic algorithms. One of them is the dependency of results to initial population. The more important weakness is when two ontologies are in the form of sparse graphs or even forests, in that case the domain for crossover is not a soft one, and small changes in an individual in crossover or mutation might take it to a very far point, and most of the time, an out of goal point of course. Work is now being done on tree ontologies, and relations in them. Once we can align tree ontologies, we can model ontologies as trees and align these trees. We are also interested to extend our theory and mechanisms for matching ontologies based on their shapes, graph areas, etc.

REFERENCES

- Tourism ontology FOAM. available online: http://www.aifb.unikarlsruhe.de/WBS/meh/foam/ontologies/.
- Arora, A., Lun, C., Motwani, R., Sudan, M., and Szegedy, M. (1992). Proof verification and hardness of approximation problems. In 33rd IEEE symposium on the Foundations of Computer Science (FOCS), pages 14– 23.
- Baeza-Yates, R. and Ribeiro-Neto, B. (1999). *Modern Information Retrieval*. Addison Wesley.
- Bouquet, P., Euzenat, J., Franconi, E., Serafini, L., Stamou, G., and Tessaris, S. (2004). Specification of common network framework for characterizing alignments. *Knowledge web NoE*, 2.2.1.
- Cormen, T. H., Leiserson, C. E., Rivest, R. L., and Stein, C. (2001). Introduction to Algorithms, second edition. MIT Press.
- Doan, A., Madhavan, J., Domingos, P., and Halevy, A. (2003). Ontology matching: A machine learning approach. *Handbook on Ontologies in Information Systems. Springer-Velag*, 2003.
- Ehrig, M. and Staab, S. (2003). Qom quick ontology mapping. In *Proc. ISWC-2003*.
- Ehrig, M. and Sure, Y. (2004). Ontology mapping an integrated approach. In *1st European Semantic Web Symposium (ESWS)*, pages 76–91.
- Euzenat, J. and Valtchev, P. (2004). An integrative proximity measure for ontology alignment. In *Proc. 3rd International Semantic Web Conference (ISWC2004)*
- Gibbons, A. (1985). *Algorithmic Graph Theory*. Cambridge University Press,.
- Haeri, H., Hariri, B., and Abolhassani, H. (2006). Coincidence-based refinement of ontology alignment. In 3rd Intl. Conference on Soft Computing and Intelligent Systems.
- Kalfoglou, Y. and Schorlemmer, M. (2003). Ontology mapping: the state of the art. *The Knowledge Engineering Review*, 18:1–31.
- Levenshtein, V. I. (1966). Binary codes capable of correcting deletions, insertions and reversals. Sov. Phys. Dokl., 6:707–710.
- Melnik, S., Garcia-Molina, H., and Rahm, E. (2002). Similarity flooding: a versatile graph matching algorithm. In Proc. 18th International Conference on Data Engineering (ICDE), San Jose (CA US), pages 117–128.
- Mitra, P., Noy, N. F., and Jaiswall, A. R. (2003). Omen: A probabilistic ontology mapping tool. In *4th international semantic web conference (ISWC 2005)*, volume 3729, pages 537–547.