Crowd Simulation on a Graphics Processing Unit based on a Least Effort Model

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Abstract: Large crowd simulation is becoming a very important field of study for many researchers. In this paper we endeavour to study large crowds and their interactions with each other and the environment. The interaction is based on a very simple Least Effort Model inspired from a real world pedestrian modelling scenario. The simulation of the crowd is computationally very expensive and GPU modelling and simulation is a viable alternative computational means to accelerate the simulation process. Compute Unified Device Architecture or CUDA is used for GPU implementation.

1 INTRODUCTION

Simulation of crowd movement is very important in the areas of traffic and pedestrian interaction, game design, animation and within crowd management industries. Pedestrian crowds are ubiquitous and can be observed in many different scenarios such as crosswalks, streets, sporting and other mass gathering events, etc. When crowd densities become very large the situation can become vulnerable to unpredictable behaviour with a chance for disaster and so it is very important to obtain knowledge of this type of situation. As it is not possible to emulate this situation in the real world, simulation is possibly the only solution for understanding behaviour of large crowds. Crisis situations such as building evacuation (Pelechano and Badler, 2006) are also very important areas of study with added difficulty due to the potential of panic. Researchers are finding simulation of pedestrian movement to be very promising and as such, there are several models to simulate the movement of pedestrians in routine and/or emergency situations. There are several pedestrian models such as, cellular automata (Blue and Alder, 2001; Weifeng and Hai, 2007), social force models (Helbing, 1991; Helbing et al., 2002) and agent based models (Cherif and Chighoub, 2010) each having their own advantages and disadvantages.

While modelling the movement of pedestrians there are several things that are important to keep in mind. Pedestrians in their movement typically have a goal towards which they move but at the same time they try to avoid collisions with other pedestrians. Taking these factors into account pedestrians often try to make least effort decisions to reach their goals or destinations. In this manner pedestrians are able to have a somewhat optimized path to their destination.

In our model the environment is divided into regular cell grids and each agent or pedestrian occupies a cell grid. The model is based on Least Effort Model (LEM) (Sarmady et al., 2009). The agents move by minimizing deviation along an optimal path to their destination. Modelling pedestrian movement is computationally expensive and the computation time increases with the number of agents. To gain an advantage in computation, a Graphics Processing Unit (GPU) (Sanders and Kandrot, 2010; Kirk and Hwu, 2010) is often used. GPU were developed for graphics purposes and take advantage of many processing cores. The introduction of Compute Unified Device Architecture or CUDA (NVIDIA, 2012) by NVIDIA opened avenues for general purpose GPU (GPGPU) parallel programming. GPGPU programming is becoming a viable alternative in fields which are very computationally expensive. Modelling of crowds is one of those applications well suited to GPUs. This work addresses modelling difficulties and accelerated
crowd behaviour modelling using a GPU. Considerable emphasis is placed on attempting to make CPU and GPU ABMs match (Wilensky and Rand, 2007).

The paper is organized as follows, in section 2 a discussion about the LEM is provided, in section 3 there is a brief discussion about GPU and CUDA, in section 4 the detailed GPU implementation of the crowd model is given. Section 5 contains the speedup obtained using GPU and section 6 contains the simulation results.

2 LEAST EFFORT ALGORITHM AND MODEL DESCRIPTION

In our model, the environment is divided into a lattice of rigid regular cells. Pedestrians or agents (Hayes, 1999) are placed in the environment and occupy a single cell. For simplicity all agents are of same size. At the beginning of a simulation the agents have an initial position and also a target or goal. They move in the environment and try to reach their goal or destination using a least effort model (LEM).

A situation is emulated where the agents are on one side of the environment and they are trying to reach the other side of the environment. Agents (pedestrians) have 8 neighbour cells as shown in Figure 1.

\[ C_i = (1-n_i)\frac{D_{\text{min}}}{D_i} \]  

(1)

In equation (1) \( D_i \) is the distance of all the neighbouring 8 cells from the target and \( D_{\text{min}} \) is the minimum \( D_i \) value. This distance only gets calculated for those cells which are empty as \( n_i \) becomes 1 occupied cells and \( C_i \) evaluates to zero. So, essentially \( C_i \) gets calculated for the cells which are empty but becomes 0 for the occupied cells.

Then all the calculated \( C_i \) are ranked according to their distances in ascending order. A random number is generated for each of the agent to select a \( C_i \). In order to generate a random number, a normal distribution is used with a mean value of 0 and a standard deviation of 3. By using the above configuration numbers from 0 to 7 are generated.

Before calculating the \( C_i \)s, a check is performed to determine whether the agent is already lying in the target column and if the cell in the next row is empty. If it is empty then there is no further checking is done and agent makes move forward towards their target.

As mentioned earlier, in this model the pedestrians are placed on one side of the environment and try to move to the other side of the environment. In this scenario, the pedestrians have the goal to reach the other side of the environment only. So to achieve this goal the traditional LEM is modified and the target is simply the opposite side. As such, whenever they are in a certain column that column is their target column.

3 OVERVIEW OF GPU ARCHITECTURE AND CUDA PROGRAMMING MODEL

NVIDIA® introduced the first GPU in 1999. But endeavours to exploit the GPU for non-graphical purposes started in earnest from 2003. With the introduction of Compute Unified Device Architecture (CUDA) in 2007 NVIDIA® broadened their scope of general purpose GPU computing. GPUs are now used for many computationally intensive tasks. In this section we briefly discuss the
GPU features and CUDA architecture. Figure 2 shows a modern CUDA capable GPU architecture. GPUs are mainly based on Single Instruction Multiple Data (SIMD) architecture massive numbers of threads for data level parallelism can be launched. GPUs are organized into highly threaded streaming multiprocessors (SMs) which are the building blocks of the device.

The number of SMs varies from one GPU to another. This project is based on the FERMI™ architecture (GEFORCE™ 560ti). This particular card consists of 14 multiprocessing units, each consisting of 32 cores for a total of 448 cores. Each of these SIMD cores is capable of carrying out large data parallel applications.

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Figure 2: CUDA scalable GPU architecture (NVIDIA 2012).

Inside the GPU there is a 2 gigabytes of graphics double data rate (GDDR) DRAM memory which is an off chip memory also known as device or global memory. Apart from device memory there is other memory like shared memory which is much faster. This memory behaves like L1 cache on-chip memory and it is shared among the SIMD cores. Each multiprocessor has a fixed number of registers. There are also other memory like constant memory and texture memory used for specific applications.

3.1 CUDA Programming Model

Figure 3 shows the CUDA programming model. CUDA is basically a C/C++ program with GPU extensions. The programming model of CUDA offers the GPU as data-parallel co-processor to the CPU. In the CUDA context, the GPU is called the device, whereas the CPU is called host. At first the data is copied to the global memory of the device and then kernel functions are launched from the host program and get executed on the GPU. The threads of the CUDA application are arranged into 3D blocks and the blocks are arranged into 3D grids. Each thread has their unique identity and they can be accessed based on their block identity and size. Each of these thread blocks is computed in one of the multiprocessors and multiple thread blocks can be executed on same multiprocessor. However, there is a limit to the maximum number of threads that can be launched. Threads in a thread block are arranged in a group of 32 threads which is known as warp. This is the smallest unit in which the threads are scheduled and executed on a multiprocessor.

4 GPU IMPLEMENTATION OF LEAST EFFORT MODEL

In our model the agents are placed on the sides of the environment and each tries to reach the opposite side of the environment. The whole environment is made up of rigid cells of the same size, dividing the whole environment into rows and columns. We have considered an environment which is square of size 500x500 and the environment is divided into cells each of unit size. Each of agents is considered to occupy a space of 1 unit. Two cases are considered for simulation. In the first instance, the agents are placed on two opposite sides and in the second instance; the agents are placed on four opposite sides of the environment. In this section we briefly discuss the first case only as the basic implementation method remains same.

4.1 Implementation Details

At first the agents are placed in the environment in a random fashion up to a maximum number of rows.
Then the agents are updated with their target. For an agent placed in the bottom of the environment the initial target is the top row. In the CUDA program each agent is considered to be a single thread. The total number of threads launched is equal to total number of agents present in the environment. First, it is required to ensure that enough threads can be launched to cover all the agents present in the environment. Then the threads are arranged into grids and blocks. In each of the simulation steps a kernel function is launched which carries out the agent movement.

Inside the kernel function, before carrying out the LEM algorithm, an agent checks whether it is placed in the target column. If true, then it verifies whether the immediate cell in the forward direction is empty or not. If empty then no further calculation is performed and the agent proceeds forward. Otherwise, a cell gets selected using LEM. The pseudo code is shown below:

```plaintext
if(target_col == present column &&
    forward_cell == 0)
    move one cell forward
else
    Calculate LEM.
```

No movement is observed when all neighbouring cells are occupied.

### 4.1.1 Challenges in Implementation

The primary motivation of the implementation of the model on the GPU was to gain speed in the simulation while emulating the situation where the pedestrians are capable of decision making. Decisions of each agent are considered to be independent of the other. The biggest challenge in implementation of the algorithm was to keep the total number of agents same, without any loss. As mentioned, each of the agents is launched as single thread and all of them are executed in parallel. So a situation could arise when two agents try to access the same environment position at the same time and a loss of agent could occur. To avoid this kind of scenario, atomic features of CUDA are used.

In the atomic operations when one of the threads is performing an operation on a particular memory location, residing either in global or shared memory, then it does not get interfered with by operations of other threads. An atomic function performs a read-modify-write operation on 32 or 64 bit word which is residing in global or shared memory. Inside the kernel function, the agent first finds out the number of neighbouring empty cells, calculates the distance of the empty cell from the target and arranges them in the ascending order. After that a random number is generated which decides the cell number to be selected. Once the cell is chosen, the movement of the agent is achieved by calling the `atomicExch()` function. In this way the numbers of agents are kept intact in every simulation step and no agent (memory) gets overwritten throughout the operation.

### 4.1.2 Generation of Random Numbers

Inside the kernel function, after the agents calculates the distance of the available neighbouring empty cells to the target and arranging them in the ascending order, a random number is generated. Random numbers are generated from a normal distribution. A separate random number is obtained for each agent. This is obtained by using cuRAND library of which comes with the CUDA SDK. Before the start of the simulation, a setup kernel is launched once for the total number of pedestrians present in the environment. Inside this setup kernel there is the cuRAND application programming interface (api) `curand_init()`. This kernel function is essentially responsible to generate a seed which is later fed to the main random number generator. The random number from the normal distribution is generated by using the `curand_normal()` api. This normal distribution api is responsible for generating a random number having a mean of 0 and standard deviation of 1. But the result is multiplied by 3 to obtain the desired standard deviation.

### 5 SPEEDUP

Simulation of the LEM for a large number of agents is very computationally intensive. The primary motivation of using GPU for the implementation of the LEM algorithm is to accelerate the simulation process. The time for the simulation process is measured for both the CPU and GPU implementations. The time is only measured for the simulation and not for any memory transaction in case of the GPU. As mentioned, the GPU used is NVIDIA® GEFORCE™ 560ti and CPU used is CPU is Intel Xeon E3-1280 (server grade).

For the GPU, time for the simulation process is measured by using `cudaevent` functions. The computation time both for CPU and GPU is given in the Figure 4a and the speedup for the simulation is provided in the Figure 4b. In Figure 4a the time is measured in seconds along the Y-axis with the the number of agents along the X-axis. This speedup is measured using two groups of agents only (bi-
directional movement). The measurement of the computation speed is initialized with 1,000 agents and it is measured up to 100,000 agents with an increment of 1,000 agents in each comparison. In Figure 4b illustrates the speedup graph. When the number of agents is 1000 the speedup is 1x. A speedup of 10x is achieved when the number of agents is 7000. The typical speedup is approximately 8x as the agent population increases further.

6 SIMULATION AND RESULTS

In this section, a brief discussion of the simulation is provided followed by a discussion about the results.

6.1 Simulation

The visualization of the simulation is performed using MATLAB™. In Figure 5a the scenario depicts the initial situation of agents placed at opposite ends of the environment. In this instance, the total number of agents in the environment is 100,000. The objective of black agents is to move to the bottom and blues objective is to do the opposite. Figure 5b and c depicts the simulation in the time step of 3,000 and 6,000 respectively.

The asymmetry seen is somewhat disconcerting as similar results were not apparent from simulations when run on the CPU. At this time suspicions are on synchronization issues associated with each agent being run as a thread. Although difficult to debug on the GPU it may also be implementation of the next cell that is providing a directional bias.

More interesting dynamics are shown in Figures 6a, 6b and 6c when agents are placed on 4 sides of

![Figure 4a: Time computation of CPU and GPU.](image)

![Figure 4b: Speedup graph.](image)

![Figure 5a: Initial placement of 100,000 agents.](image)

![Figure 5b: Simulation of 100,000 agents in time step 3,000.](image)

![Figure 5c: Simulation of 100,000 agents in time step 6,000.](image)
the environment. Figure 6a is the initial position of the agents; Figure 6b and 6c are the simulation of agents in time step of 3,000 and 6,000 respectively. The total number of agents placed on 4 sides of the environment in this scenario is 160,796. The asymmetry of the flow is apparent resulting from the asymmetry of the initial conditions as well as from the underlying model implementation.

Figure 6a: Initial placement of 160,796 agents.

Figure 6b: Simulation step 3,000.

Figure 6c: Simulation step 6,000.

6.2 Results

In this section a brief discussion of throughput of the agents is provided. The throughput is measured by considering only two groups of pedestrians at opposite sides. The number of rows up to which the agents would occupy the environment initially is kept fixed. The time required for the throughput i.e. the agents to cross the other side of the environment is measured. The total number of time steps is also kept fixed at a maximum of 6000. Initially the throughput results are obtained for only 1000 agents (500 agents on each side). Then the next result is carried out for 2000 agents (1000 agents on each side) so on and so forth until the number of agents reaches to 100,000. Figure 7a depicts the throughput when the number of agents is 1000.

In Figure 7a, when the number of agents is 1000 the distribution seems to be almost uniform. All the agents are seen to cross the environment between time steps of 290 to 390. In Figure 7b, c, d and e the number of agents are 25,000, 50,000, 75,000 and 100,000 respectively.

In Figure 7b, when the number of agents is 25,000 a peak is observed between the time steps of 935 and 945 with a maximum throughput of approximately 1,100. When the number of agents increased to 50,000 the peak is observed between the time steps of 2020 and 2030 with a maximum throughput of approximately 1,250.

From Figure 7e, when the number of agents increased to 100,000 the throughput takes an almost steady rise with the peak at the time step of 3,990 and then a sudden fall. There are 918 agents who are not able to cross the environment within the given time step of 6,000.

Figure 8a shows the mean of the throughput of the agents. The x-axis on the graph is the increase in the number of agents starting from 1,000 up to 100,000 with an increment of 1,000 in each step and y-axis the average time taken. Figure 8b shows a...
heat map surface plot. From the surface plot, the mean plot of Figure 8a is clearly seen. From the top view it is clear that when the density is low then the throughput of the agents is concentrated in a small region which spreads in a non-Gaussian manner as the density increases.

The periodic behaviour of the throughput is believed to be an artefact of the computation on the GPU. The cause of these phenomena is still being investigated. From Figure 8b it is observed that the distribution spreads along x-axis, with the increase in the number of agents. This spread however is heavily skewed. With the increase of the number of agents (measured along y-axis) the peak throughput values in some cases occurs at lower time (measure on the x-axis). This causes the graph of Figures 8a and 8b to take an unexpected and artificial periodic shape. One of such instance is shown in Figure 9. In Figure 9 peak of the throughput of 43,000 agents is less than that of 35,000 agents illustrating the periodic nature of Figure 8a and 8b.
These anomalies are systematic computation artefacts and would not seem to be real phenomena associated with crowd movement. We are again attempting to resolve this issue. Even in light of the anomalies and model discrepancies agent based modelling of pedestrians appears to be well suited to GPGPUs.

![Figure 9: Comparing throughput of 35,000 and 43,000 agents](image)

7 FUTURE WORK

In future work, in addition to concentrating on model matching and validation, it would be desirable to incorporate some psychological state in each agent to achieve a more realistic world scenario. There would also be some changes in the implementation of the CPU where we would utilize all the cores available on the CPU. The latest release of CUDA 5 also has dynamic parallelism that can be implemented on the KEPLER architecture. The GPU that has been used here has the FERMI architecture. As such, there are plans to use a GPU card with the KEPLER architecture. In addition, the obvious asymmetries in the simulations require further investigation.

REFERENCES


