Keywords: Intertransaction association rule mining, MFI.

Abstract: In this paper we propose an efficient method for mining frequent intertransaction itemsets. Our approach consists in mining maximal frequent itemsets (MFI) by extending the SmartMiner algorithm for the intertransaction case. We have called the new algorithm InterTraSM (Inter Transaction Smart Miner). Because it uses depth first search the memory needed by the algorithm is reduced; a strategy for passing tail information for a node combined with a dynamic reordering heuristic lead to improved speed. Experiments comparing InterTraSM to other existing algorithms for mining frequent intertransaction itemsets have revealed a significant gain in performance. Further development ideas are also discussed.

1 INTRODUCTION

Association rule mining is a field of the data mining domain that has developed extensively in the last years. After the problem was introduced in (Agrawal, Imielinski and Swami, 1993) and the A-Priori algorithm was introduced in (Agrawal and Srikant, 1994), the research expanded into a vast number of directions. There have been proposed other algorithms for the original problem, either Apriori-like or with a new structure. Also new algorithms have appeared for mining episodes and sequential patterns, mining correlations, mining generalized, multilevel or quantitative association rules. Our algorithm contributes to another of these new directions, mining intertransaction association rules.

The initial association rule mining problem ignored any correlation between the transactions and searched for associations only between items inside a transaction – we call this case intratransaction analysis. To search for associations between items across several transactions ordered on a dimension (usually time or space), intertransaction association rule mining has been used.

We use the stock market database example to differentiate between intra- and inter-transaction analysis. If the database contains the price for each stock at the end of the trading day, an intratransaction association rule might be “If stock prices for companies A and B go up for one day, there is a probability of over c% that the price for company C will also go up the same day”. However, analysts might be more interested in rules like “If stock prices for companies A and B go up for one day, there is a probability of over c% that the price for company C will go up two days later.” This rule describes a relationship between items from different transactions, and it can be discovered only by using intertransaction analysis.

Several algorithms for intertransactional association rule mining have been introduced.

The E-Apriori (Extended Apriori) and EH-Apriori (Extended Hash Apriori) algorithms have been proposed by (Lu, Feng and Han, 2000). They use the Apriori algorithm for mining frequent intertransaction itemsets. The EH-Apriori algorithm also uses a hash in order to reduce the number of candidate intertransaction itemsets with two elements.

(Tung et al., 2003) developed FITI – an algorithm that discovers first the frequent intratransaction itemsets and then uses them to generate the frequent inter-transaction itemsets.

The ITP-Miner algorithm from (Lee and Wang, 2007) uses a structure called dat-list to store item information and an ITP-tree to store discovered frequent inter-transaction patterns.

The EFP-Tree algorithm presented in (Luhr, West and Venkatesh, 2007) is an extension of the FP-Tree (Frequent Pattern Tree) algorithm for the
intertransaction case. It uses a divide-and-conquer approach to avoid candidate generation.

In this paper we propose a new method for mining frequent intertransaction itemsets called InterTraSM. This algorithm is an extension for the intertransactional case of the SmartMiner algorithm presented in (Zou, Chu and Lu, 2002).

As the authors remark in the above mentioned paper, mining frequent itemsets is infeasible when the frequent patterns are long because of the exponential number of frequent itemsets. An alternative is mining maximal frequent itemsets (MFI) – itemsets that are not a subset of any other frequent itemset. Once we have obtained the MFI we can easily obtain all the frequent itemsets, who can then be counted for support in a single scan of the database.

Like SmartMiner, InterTraSM uses a depth first search (DFS) to determine maximal frequent itemsets. It also uses a strategy for passing tail information for a node combined with a dynamic reordering heuristic that improve the speed of execution.

The remainder of this paper is organized as follows: in Section 2 we will give a formal definition for the problem of intertransaction association rules mining. In Section 3 we will describe our proposed algorithm, InterTraSM. In Section 4 we will present the experimental results we have obtained so far, and we will present the conclusions and plans for future work in Section 5.

2 PROBLEM DESCRIPTION

In this section we introduce some notations and we present a formal definition for the problem of mining intertransaction association rules.

Definition 2.1. Let \( I = \{e_1, e_2, \ldots, e_n\} \) be a set of items and let \( D \) be an attribute with values within the domain \( \text{Dom}(D) \). We call a transactional database a database with transactions (records) of the form \( (d, S) \), where \( d \in \text{dom}(D) \) and \( S \subseteq I \).

We restrict our search for associations to a maximum span of transactions, given as an input parameter.

Definition 2.2. A sliding window \( W \) in a transactional database \( T \) represents a set of continuous intervals from the domain \( D \), such that there exists in \( T \) a transaction associated to the first interval from \( W \). Each interval from \( W \) is called a subwindow of \( W \), and they are numbered corresponding to their temporal order \( d_0, d_1, \ldots, d_m \). We also use the notation \( W[0], W[1], \ldots, W[m] \).

Definition 2.3. Let \( T \) be a transactional database, let \( I \) be the set of items with \( n = |I| \) and let \( W \) be a sliding window with \( w \) intervals. A megatransaction \( M \) associated with \( W \) is the set
\[
M = \{e_j(i) | e_i \in W[j], 1 \leq i \leq n, 0 \leq j \leq w-1\}.
\]

Items from a megatransaction will be called from now on extended items.

Let \( E \) be the set of all possible extended items
\[
E = \{e_1(0), e_1(1), \ldots, e_1(w-1), e_2(0), \ldots, e_n(w-1)\}
\]

We call an intratransaction itemset a set of items \( A \subseteq I \). We call an intertransaction itemset a set of extended items \( B \subseteq E \) that contains at least an extended item \( e_i(0) \), with \( 1 \leq i \leq n \).

Definition 2.4. An inter-transaction association rule has the form \( X \rightarrow Y \), where:

i) \( X, Y \subseteq E \)
ii) There is at least one element \( e_i(0) \) in \( X, 1 \leq i \leq n \)
iii) There is at least one element \( e_i(j) \) in \( Y, 1 \leq i \leq n, 1 \leq j \leq w-1 \)
iv) \( X \) and \( Y \) are disjoint

Let \( T_{XY} \) be the set of megatransactions that contain the set \( X \cup Y \), let \( T_X \) be the set of megatransactions that contain the set \( X \) and let \( N \) be the total number of transactions.

Then \( S = |T_{XY}| / N \) and \( C = |T_{XY}| / |T_X| \) are the support and confidence for the intertransaction association rule.

As in the classical case, the problem of mining intertransaction association rules can be divided in two parts:

- finding the frequent itemsets
- generating the association rules.

The second problem takes much less computational time than the first one, so it presents little interest for research. A solution has been discussed for example in (Tung et al., 2003). Our algorithm (like the algorithms mentioned before) will therefore focus on a solution for the first problem.
3 ALGORITHM DESCRIPTION

We now introduce the InterTraSM algorithm. As we mentioned before, InterTraSM is an adaptation of the SmartMiner algorithm described in (Zou, Chu and Lu, 2002) for intertransaction association rule mining. Therefore the theoretical foundations of the two algorithms are very similar. Our main contribution has been to identify intertransaction mining as a domain where searching for MFI would lead to an improvement in performance, and to apply and customize the existing algorithm to the intertransaction analysis case; we also provided a new implementation.

InterTraSM finds maximal frequent itemsets (MFI) of extended items from a transactional database. The algorithm uses depth first search and for performance optimization it uses a dynamic reordering to eliminate infrequent items from the tail of a current node. A hash table is also used to save the itemsets discovered as frequent at node-level, in order not to go down a tree path that was already investigated while exploring a maximal frequent itemset.

As we mentioned the algorithm performs a depth first search, so at any step it works on a node from a search tree. We describe below the data managed at the level of a node used by the algorithm and how the data is processed.

A node \( N \) is identified as \( X:Y \), where \( X \) (the head) is the set of items that have been discovered to be part of a frequent itemset, and \( Y \) (the tail) is the set of items that still have to be explored. The purpose of the node is to find maximal frequent itemsets in the transaction set \( T(X) \) – all the transactions that include \( X \).

The starting node is \( \Phi:E \) (the empty set and the set of all the possible extended items).

The entry data for a node are:
- the transaction set \( T(X) \)
- the tail \( Y \)
- the global data information \( G_{inf} \), which is the tail information for the node known so far (this contains the itemsets that have been discovered in a previous step to be frequent in \( T(X) \)).

The exit data for a node are:
- the updated \( G_{inf} \)
- the discovered maximum frequent itemsets \( M_{fi} \).

The data processing at a node \( N \) is described below:

- count the support for each item from \( Y \) in the transaction set \( T(X) \)
- remove the infrequent items
- while (\( Y \) has at least one element)
  - select from \( Y \) an item \( a_i \) to be the head of the next state \( S_1 \)
  - \( Y_{i+1} = Y - a_i \) will be the tail for \( S_{i+1} \)
  - obtain the auxiliary tail information for \( S_{i+1} \) by projecting on \( Y_1 \) the itemsets that contain \( a_0 \) from the tail information \( G_{inf} \).
  - recursively call the algorithm for the node \( N_{i+1} = X_{a_i}:Y_{i+1} \). The returned values will be \( M_{fi} \) and the updated tail information.
  - \( Y = Y_{i+1} \)
- end (while)

The processing of the node returns the maximal frequent itemsets to be \( M_{fi} = \bigcup a_i M_{fi} \), and the updated \( G_{inf} \) as the itemsets in the original \( G_{inf} \) that have not been marked as deleted.

As we mentioned, InterTraSM uses extended items instead of the intratransaction items used by SmartMiner. A customization for this case is that the first node we select while searching in depth from the root of the tree corresponds only with items from the first interval (interval 0). This was done because each frequent itemset has to have at least an item from interval 0.

We created our own implementation of the algorithm using C, with a structure similar to the one for the SmartMiner algorithm described in (Zou, Chu and Lu, 2002) – but with modifications for the intertransaction case (The SmartMiner algorithm was implemented in Java). We felt that writing the algorithm in C enabled us to better control the memory use of the algorithm.

4 PERFORMANCE STUDY

We used both synthetic data and real data to evaluate the performance of the algorithm.

To generate the synthetic data we used the same generator as the one described in (Luhr, West and Venkatesh, 2007) to evaluate EFP-Tree, gracefully provided to us by the authors. It uses the same method as the one used to evaluate FITI and ITP-Miner.

The real data consists of two datasets, WINNER and LOSER, similar to those used in (Lee and
They have been obtained from the values of ten stock exchange indexes from January 1, 1991 to December 31, 2005. In the WINNER set a transaction for a trading day contains the stock indices that rise for the day, and in the LOSER set we have the stock indices that fall. The stock indices used are the ASX All Ordinaries Index (ASX), CAC40 Index (CAC), DAX Index (DAX), Dow Jones Index (DOW), FT-SE 100 Index (FTS), Hang Seng Index (HSI), NASDAQ Index (NDQ), Nikkei 225 Index (NKY), Swiss Market Index (SMI), and Singapore ST Index (STI).

Two synthetic data sets representing sparse and dense data were generated, with parameters identical to those used in the evaluation of EFP-Tree.

Table 1 lists the parameters used to create the data sets used in the experimentation.

Table 1: Parameters used in the generation of the synthetic data sets.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Sparse</th>
<th>Dense</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of intratransactions</td>
<td>500</td>
<td>200</td>
</tr>
<tr>
<td>Size of the intertransaction pool</td>
<td>50</td>
<td>200</td>
</tr>
<tr>
<td>Average length of intratransactions</td>
<td>5</td>
<td>25</td>
</tr>
<tr>
<td>Maximum length of intratransactions</td>
<td>10</td>
<td>50</td>
</tr>
<tr>
<td>Average length of intertransactions</td>
<td>5</td>
<td>8</td>
</tr>
<tr>
<td>Maximum length of intertransactions</td>
<td>10</td>
<td>20</td>
</tr>
<tr>
<td>Maximum number of unique items</td>
<td>500</td>
<td>100</td>
</tr>
<tr>
<td>Maximum interval span of intertransactions</td>
<td>4</td>
<td>6</td>
</tr>
</tbody>
</table>

The program was benchmarked under a Microsoft Windows XP OS, on a PC with Intel Pentium IV CPU with speed of 3GHz and main memory of 1GB. The code has been written and compiled using Microsoft Visual Studio 2003.

Since FITI has been shown to be more computationally efficient than EH-Apriori and both EFP-Tree and ITP-Miner have been shown to outperform FITI, we have only made comparisons with EFP-Tree and ITP-Miner.

We have compared the execution times of InterTraSM on the synthetic data sets with the execution times reported for EFP-Tree in (Luhr, West and Venkatesh, 2007) for synthetic data generated with the same parameters by the same generator (probably different actual data though, since it is a random generator).

For the synthetic sparse data set we have gradually lowered the support threshold from 1.6% to 0.6%, using a fixed intertransaction window size of 4. For the synthetic dense data set we have gradually lowered the support threshold from 13% to 8%, using a fixed intertransaction window size of 6.

The execution times for InterTraSM are displayed in Figure 1 (sparse data) and Figure 2 (dense data).

We have then incremented the intertransaction window size from w=0 to w=10 with fixed minimum supports of 1% for the sparse data and 10% for the dense data. The results are displayed in Figure 3 (sparse data) and Figure 4 (dense data).
We can see that all the execution times displayed in these charts are under one second, while the execution times reported for the EFP-Tree in (Luhr, West and Venkatesh, 2007) for synthetic data generated with the same parameters have values of tens, even hundreds of seconds. Since the processors used have similar performances, even accounting for the different implementation languages (C versus Ruby), InterTraSM seems to perform at least an order of magnitude better than EFP-Tree.

We have compared the execution times of InterTraSM on the real data sets with the execution times reported for ITP-Miner on transactions obtained from the same stock indices values. For both the WINNER and LOSER datasets we have used an intertransaction window size of 4 and we have varied the minimum support threshold from 4% to 12%. The results are displayed in Figure 5 (WINNER data set) and Figure 6 (LOSER data set).

Comparing the results with those reported for running ITP-Miner on the same data sets in (Lee and Wang, 2007), on a processor with similar performances using Microsoft Visual C++ 6.0, it seems that there is an order of magnitude difference in favor of InterTraSM, especially when the minimum support value decreases. For example on the LOSER data set with the support threshold at 4% InterTraSM takes less than 6 seconds, while the authors reported the ITP-Miner takes about 100 seconds.

5 CONCLUSIONS

In this paper we have proposed a new algorithm for mining intertransaction association rules called InterTraSM, an extension of the SmartMiner algorithm for the intertransaction case. InterTraSM uses depth first search and mines for maximal frequent itemsets (from which all the frequent itemsets can be easily generated). Previous algorithms have mined for all the frequent itemsets and they thus have had to count support for an exponentially large number of frequent itemsets compared to our algorithm. Experiments with similar data and on similar machines to those used for evaluating EFP-Tree and ITP-Miner have shown that InterTraSM outperforms them by at least an order of magnitude, especially when the minimum support threshold is reduced – generating longer maximal frequent itemsets.

In our future research we want to apply the algorithm to some more real data sets and see how it performs. We also want to extend the algorithm from 1-dimensional to n-dimensional transactional databases.
REFERENCES


