AUTOMATIC TRANSFORMATION OF SQL RELATIONAL DATABASES TO OWL ONTOLOGIES

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Abstract: This paper proposes a novel approach to automatic transformation of relational databases (written in SQL) to ontologies (written in OWL), where domains and constraints CHECK are also considered. The proposed approach can identify (inverse) functional, symmetric and transitive properties, cardinality and value restrictions, and enumerated classes and data types.

1 INTRODUCTION

Today it is common to get data from relational databases over the Web. These databases are generally separate and not easily used as merged data sources. The W3C vision sees ways to unify the description and retrieval of the data using ontologies, thus allowing much of the Web to be part of a large interoperable database.

Thus, there is a need to transform relational databases to ontologies. However, manual transformation is hard to do and often takes a lot of time. Thus, there is also a need to automate the transformation.

2 RELATED WORK

While there are several approaches to automatic transformation relational databases to ontologies - e.g. (Buccella et al., 2004; Li et al., 2005; Shen et al., 2006; Astrova and Kalja, 2006; Sequeda et al., 2007), many situations are too complex or require more flexibility than the existing approaches enable.

E.g. a company may wish to trace the skills of its employees in order to assign the employees to the projects. Since an employee may have many skills, skill becomes multivalued. It is possible to represent skill as a data type property in the ontology. But it is not possible to represent skill as a column in the relational database, because the column may have at most one value for each row in the table (atomicity). One solution to this problem is to create a separate table (McFadden et al., 1999). This table could map to a data type property during the transformation. However, the existing approaches cannot recognize such a situation. Rather, they map the table to a class. Moreover, the existing approaches cannot identify inverse functional, symmetric and transitive properties, value restrictions, and enumerated classes.

As an attempt to resolve these problems, this paper proposes a novel approach to automatic transformation of relational databases to ontologies. The main objective of the proposed approach is to preserve as many semantics as possible during the transformation. The proposed approach assumes that a relational database is written in SQL (SQL, 2002) and that an ontology is written in OWL (OWL, 2004).

3 APPROACH

The proposed approach maps constructs of a relational database (i.e. tables, domains, columns, constraints, and rows) to an ontology using the names of constructs in the relational database as the names of constructs in the ontology. Next this mapping will be illustrated by example. An example is the relational database of a company.

3.1 Mapping Tables

A table can be mapped to three different constructs in the ontology: a class, a data type property, and an object property and its inverse.
A table \textit{Project} in Figure 1 has its own primary key. Therefore, this table maps to a class \textit{Project}.

\begin{verbatim}
CREATE TABLE Project(
    ProjectID INTEGER PRIMARY KEY)
\end{verbatim}

\begin{verbatim}
/owl:Class rdf:ID="Project"/>
\end{verbatim}

Figure 1: Table maps to class.

The primary key of a table \textit{SoftwareProject} in Figure 2 is a foreign key to another table \textit{Project}. Therefore, this table maps to a class \textit{SoftwareProject}.

\begin{verbatim}
CREATE TABLE SoftwareProject(
    ProjectID INTEGER PRIMARY KEY,
    FOREIGN KEY (ProjectID) REFERENCES Project)
\end{verbatim}

\begin{verbatim}
/owl:Class rdf:ID="SoftwareProject"/>
\end{verbatim}

Figure 2: Table maps to class (contd.).

The primary key of a table \textit{Involvement} in Figure 3 is composed of foreign keys to two other tables \textit{Project} and \textit{Employee}, indicating a binary (many-to-many) relationship. Since there are no other columns in the table \textit{Involvement}, it maps to two object properties: \textit{EmployeeID} (that uses classes \textit{Project} and \textit{Employee} as its domain and range, respectively) and \textit{ProjectID}. The latter is an inverse of the former, meaning that the relationship is bidirectional (i.e. a project involves employees and an employee is involved in projects). If the table \textit{Involvement} had an additional column say \textit{hours}, it would be mapped to a class \textit{Involvement}.

\begin{verbatim}
CREATE TABLE Involvement(
    EmployeeID INTEGER REFERENCES Employee,
    ProjectID INTEGER REFERENCES Project,
    PRIMARY KEY (EmployeeID, ProjectID))
\end{verbatim}

\begin{verbatim}
/owl:ObjectProperty rdf:ID="EmployeeID">
  <rdfs:domain rdf:resource="#Project"/>
  <rdfs:range rdf:resource="#Employee"/>
</owl:ObjectProperty>

/owl:ObjectProperty rdf:ID="ProjectID">
  <owl:inverseOf rdf:resource="#EmployeeID"/>
</owl:ObjectProperty>
\end{verbatim}

Figure 3: Table maps to object property and its inverse.

The primary key of a table \textit{SkillValue} in Figure 4 is composed of a column \textit{skill} and a foreign key to another table \textit{Employee}, meaning that the column \textit{skill} is multivalued (i.e. an employee may have zero or more skills). Since there are no other columns in the table \textit{SkillValue}, it maps to a data type property \textit{skill} that uses a class \textit{Employee} as its domain. If the table \textit{SkillValue} had an additional column say \textit{level}, it would be mapped to a class \textit{SkillValue}.

\begin{verbatim}
CREATE TABLE SkillValue(
    skill VARCHAR,
    EmployeeID INTEGER REFERENCES Employee,
    PRIMARY KEY (skill, EmployeeID))
\end{verbatim}

\begin{verbatim}
/owl:DatatypeProperty rdf:ID="skill">
  <rdfs:domain rdf:resource="#Employee"/>
  <rdfs:range rdf:resource="&xsd:string"/>
</owl:DatatypeProperty>
\end{verbatim}

Figure 4: Table maps to data type property.

The primary key of a table \textit{Involvement} in Figure 5 is composed of foreign keys to three other tables \textit{Employee}, \textit{Project} and \textit{Skill}, indicating a ternary relationship. Since only binary relationships can be represented through object properties, this table maps to a class \textit{Involvement}.

\begin{verbatim}
CREATE TABLE Involvement(
    EmployeeID INTEGER REFERENCES Employee,
    ProjectID INTEGER REFERENCES Project,
    SkillID INTEGER REFERENCES Skill,
    PRIMARY KEY (EmployeeID, ProjectID, SkillID))
\end{verbatim}

\begin{verbatim}
/owl:Class rdf:ID="Involvement"/>
\end{verbatim}

Figure 5: Table maps to class (contd.).

3.2 Mapping Domains

A domain maps to a class unless there is a constraint \texttt{CHECK} with enumeration on it. Then it maps to an enumerated class.

A domain \textit{ProjectType} in Figure 6 is defined as the data type of all strings. Therefore, this domain
maps to a class ProjectType, with a data type property say type that uses string as its range.

CREATE DOMAIN ProjectType AS VARCHAR

[owl:Class rdf:ID="ProjectType"]
  [owl:DatatypeProperty rdf:ID="type"]
    [rdfs:range rdf:resource="&xsd;string"/>
  [owl:DatatypeProperty>
</owl:Class>

Figure 6: Domain maps to class.

A domain ProjectType in Figure 7 is defined as the data type of all strings, again. However, there is now a constraint CHECK on it. This constraint specifies the domain ProjectType through a list of values Software and Hardware (also known as enumeration). Therefore, it maps to an enumerated class ProjectType, with individuals for each value in the list.

CREATE DOMAIN ProjectType AS VARCHAR
CONSTRAINT ProjectType_Constraint
CHECK IN ('Software', 'Hardware')

[owl:Class rdf:ID="ProjectType"]
  [owl:oneOf rdf:parseType="Collection">
    [owl:Thing rdf:about="#Software"/>
    [owl:Thing rdf:about="#Hardware"/>
  [owl:oneOf>
</owl:Class>

Figure 7: Domain maps to enumerated class.

### 3.3 Mapping Columns

A column that is not (part of) a foreign key maps to a data type property unless it uses a domain as its data type. Then it maps to an object property. This is because the domain maps itself to either a class or an enumerated class (see Section 3.2).

A column ssn in a table Employee in Figure 8 is not a foreign key. Therefore, this column maps to a data type property ssn that uses a class Employee as its domain. This property has a maximum cardinality of 1, because the column ssn may have at most one value for each row in the table Employee (atomicity). Alternatively, the property ssn could be defined as functional, which is the same as saying that the maximum cardinality is 1. It should be noted that if the column ssn were a surrogate key, it would be ignored. A surrogate key is internally generated by the relational database management system using an automatic sequence number generator or its equivalence; e.g. an IDENTITY in SQL Server and Sybase, a SEQUENCE in Oracle and an AUTO_INCREMENT in MySQL.

CREATE TABLE Employee(
  ssn INTEGER CHECK (ssn > 0))

[owl:DatatypeProperty rdf:ID="ssn"]
  [rdfs:domain rdf:resource="#Employee"/>
  [rdfs:range rdf:resource="&xsd;positiveInteger"/>
</owl:DatatypeProperty>

[owl:Class rdf:ID="Employee"]
  [rdfs:subClassOf]
    [owl:Restriction]
      [owl:onProperty rdf:resource="#ssn"/>
      [owl:maxCardinality rdf:datatype="&xsd;nonNegativeInteger"1/>
  [rdfs:subClassOf>
</owl:Class>

Figure 8: Column maps to data type property.

Most of the mapping of columns has to do with the mapping of data types from SQL to XSD. Unlike SQL, OWL does not have any built-in data types. Instead, it uses XSD (XML Schema Data types).

A column ssn in Figure 8 uses INTEGER as its data type. Therefore, a data type property ssn could use integer as its range. However, there is a constraint CHECK on the column ssn. This constraint further restricts the range of values for the column ssn to all integers greater than 0 (i.e. all positive integers). Therefore, the data type property ssn uses positiveInteger as its range.

### 3.4 Mapping Constraints

SQL supports constraints UNIQUE, NOT NULL, REFERENCES, FOREIGN KEY, PRIMARY KEY, CHECK, and DEFAULT. However, not all the constraints can be mapped to OWL. E.g. a constraint DEFAULT (that defines a default value for a given column) has no correspondence in OWL. Therefore, it is ignored.

#### 3.4.1 Mapping Constraints UNIQUE

UNIQUE is a column constraint. It maps to an inverse functional property.

A constraint UNIQUE in Figure 9 specifies that a column ssn in a table Employee is unique,
meaning that no two rows in the table Employee have the same value for the column ssn (i.e. social security numbers uniquely identify employees). Therefore, this constraint maps to an inverse functional property.

CREATE TABLE Employee(
  ssn INTEGER UNIQUE
)


Figure 9: Constraint UNIQUE maps to inverse functional property.

3.4.2 Mapping Constraints NOT NULL

NOT NULL is a column constraint. It maps to a minimum cardinality of 1.

A constraint NOT NULL in Figure 10 specifies that a column ssn in a table Employee is not null, meaning that all rows in the table Employee have values for the column ssn (i.e. all employees are assigned social security numbers). Therefore, this constraint maps to a minimum cardinality of 1.

CREATE TABLE Employee(
  ssn INTEGER NOT NULL
)


Figure 10: Constraint NOT NULL maps to minimum cardinality of 1.

3.4.3 Mapping Constraints REFERENCES and FOREIGN KEY

REFERENCES is a column constraint (to refer to a single column), whereas FOREIGN KEY is a table constraint (to refer to multiple columns). Both constraints are used for specifying foreign keys. A foreign key can be mapped to four different constructs in the ontology: an object property, class inheritance, a symmetric property, and a transitive property.

A constraint REFERENCES in Figure 11 specifies that a column ProjectID in a table Task is a foreign key to another table Project, indicating a binary (one-to-zero-or-one, one-to-one or many-to-one) relationship. Since the foreign key is not the primary key, it maps to an object property ProjectID that uses classes Task and Project as its domain and range, respectively. This property has a maximum cardinality of 1 (atomicity). In addition, the property ProjectID is restricted to all values from the class Project, because the foreign key implies that for each (non-null) value of the column ProjectID there is the same value in the table Project.

CREATE TABLE TASK(
  TaskID INTEGER PRIMARY KEY,
  ProjectID INTEGER REFERENCES Project
)


Figure 11: Foreign key maps to object property.

A constraint FOREIGN KEY in Figure 12 specifies that a column ProjectID in a table SoftwareProject is a foreign key to another table Project, indicating a binary relationship, again. However, since the foreign key is now the primary key, it maps to class inheritance: SoftwareProject is a subclass of Project (i.e. a software project is a project).

CREATE TABLE SoftwareProject(
  ProjectID INTEGER PRIMARY KEY,
  FOREIGN KEY (ProjectID) REFERENCES Project
)


Figure 12: Foreign key maps to class inheritance.

A constraint REFERENCES in Figure 13 specifies that a column spouse in a table

indicating a binary (one-to-zero-or-one, one-to-one or many-to-one) relationship. Since the foreign key is not the primary key, it maps to an object property ProjectID that uses classes Task and Project as its domain and range, respectively. This property has a maximum cardinality of 1 (atomicity). In addition, the property ProjectID is restricted to all values from the class Project, because the foreign key implies that for each (non-null) value of the column ProjectID there is the same value in the table Project.

CREATE TABLE TASK(
  TaskID INTEGER PRIMARY KEY,
  ProjectID INTEGER REFERENCES Project
)


Figure 11: Foreign key maps to object property.

A constraint FOREIGN KEY in Figure 12 specifies that a column ProjectID in a table SoftwareProject is a foreign key to another table Project, indicating a binary relationship, again. However, since the foreign key is now the primary key, it maps to class inheritance: SoftwareProject is a subclass of Project (i.e. a software project is a project).

CREATE TABLE SoftwareProject(
  ProjectID INTEGER PRIMARY KEY,
  FOREIGN KEY (ProjectID) REFERENCES Project
)


Figure 12: Foreign key maps to class inheritance.

A constraint REFERENCES in Figure 13 indicates a binary (one-to-zero-or-one, one-to-one or many-to-one) relationship. Since the foreign key is not the primary key, it maps to an object property ProjectID that uses classes Task and Project as its domain and range, respectively. This property has a maximum cardinality of 1 (atomicity). In addition, the property ProjectID is restricted to all values from the class Project, because the foreign key implies that for each (non-null) value of the column ProjectID there is the same value in the table Project.
Employee is a foreign key to the same table, indicating a unary relationship. Therefore, the foreign key maps to a symmetric property `spouse` that uses a class `Employee` as both its domain and range (i.e. if one employee is a spouse of another employee, then the second employee is a spouse of the first employee).

CREATE TABLE Employee(
    EmployeeID INTEGER PRIMARY KEY,
    spouse INTEGER REFERENCES Employee
)

↓

<owl:SymmetricProperty rdf:ID="spouse">
    <rdfs:domain rdf:resource="#Employee"/>
    <rdfs:range rdf:resource="#Employee"/>
</owl:SymmetricProperty>

Figure 13: Foreign key maps to symmetric property.

A constraint `REFERENCES` in Figure 14 specifies that a column `subtask` in a table `Task` is a foreign key to the same table, indicating a unary relationship, again. However, since the foreign key is now accompanied by a trigger `ON DELETE CASCADE`, this relationship consists of a whole and a part, where the part cannot exist without the whole (i.e. if a task is deleted, then all its subtasks must also be deleted). Therefore, the foreign key maps to a transitive property `subtask` that uses a class `Task` as both its domain and range (i.e. if one task is a subtask of another task and the other task is a subtask of yet another task, then the first task is a subtask of the third task).

CREATE TABLE Task(
    TaskID INTEGER PRIMARY KEY,
    subtask INTEGER REFERENCES Task ON DELETE CASCADE
)

↓

<owl:TransitiveProperty rdf:ID="subtask">
    <rdfs:domain rdf:resource="#Task"/>
    <rdfs:range rdf:resource="#Task"/>
</owl:TransitiveProperty>

Figure 14: Foreign key maps to transitive property.

3.4.4 Mapping Constraints PRIMARY KEY

There are two forms of constraint `PRIMARY KEY`: using it as a column constraint (to refer to a single column) and using it as a table constraint (to refer to multiple columns). A constraint `CHECK` maps to a value restriction unless it has enumeration. Then it maps to an enumerated data type. It should be noted that OWL is not powerful enough to express all the value restrictions that can be imposed by a constraint `CHECK` (e.g. an employee’s age as an integer between 18 and 65).

A constraint `CHECK` in Figure 16 specifies that a column `type` in a table `Project` may have only a value `Software`. Therefore, a data type property `type` is restricted to have the same value for all instances in a class `Project`.

3.4.5 Mapping Constraints CHECK
CREATE TABLE Project(
    type VARCHAR
)

↓

<owl:Class rdf:ID="Project">
    <rdfs:subClassOf>
        <owl:Restriction>
            <owl:onProperty rdf:resource="#type"/>
            <owl:hasValue rdf:datatype="&xsd;string">Software</owl:hasValue>
        </owl:Restriction>
    </rdfs:subClassOf>
</owl:Class>

Figure 16: Constraint CHECK maps to value restriction.

A constraint `CHECK` in Figure 17 specifies the range for a column `type` in a table `Project` through a list of values `Software` and `Hardware`. Therefore, this constraint maps to an enumerated data type `Project`, with one element for each value in the list.

CREATE TABLE Project (
    type VARCHAR
)

↓

<owl:DatatypeProperty rdf:ID="type">
    <rdfs:domain rdf:resource="#Project"/>
    <rdfs:range>
        <owl:DataRange>
            <owl:oneOf>
                <rdf:List>
                    <rdf:first rdf:datatype="&xsd;string">Software</rdf:first>
                    <rdf:rest rdf:resource="&rdf;nil"/>
                </rdf:List>
            </owl:oneOf>
        </owl:DataRange>
    </rdfs:range>
</owl:DatatypeProperty>

Figure 17: Constraint CHECK maps to enumerated data type.

3.5 Mapping Rows

A row maps to an instance.

A row in a table `Project` in Figure 18 has a value `Software` for a column `type`. Therefore, this row maps to an (anonymous) instance of a class `Project` that has the same value for a data type property `type`.

INSERT INTO Project (type) VALUE ('Software')

↓

<Project>
    <type>
        <rdf:datatype="&xsd;string">Software</rdf:datatype>
    </type>
</Project>

Figure 18: Row in table maps to instance of class.

4 CONCLUSIONS

This paper has proposed a novel approach to automatic transformation of relational databases to ontologies, where domains and constraints `CHECK` are also considered. The proposed approach can map all constructs of a relational database to an ontology, with the exception of those constructs that have no correspondences in the ontology (e.g. constraints `DEFAULT`).

REFERENCES


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