A Relation-Algebra Language to Specify Declarative Business Rules

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Abstract: Business rules that apply within a business context must be formulated in a comprehensible way to allow validation by their stakeholders, but at the same time they must be specified with enough precision to assure their correct implementation in computer applications. These opposing demands of business rule modeling are not easily reconciled. Formal rule modeling languages may be exact but they are often lacking in understandability, whereas controlled natural languages are more easily understood but generally fall short in exactness. We use Relation Algebra as the foundation to set up a controlled language for declarative business rules that is compatible with practical demands, such as laid out in the Business Rules Manifesto. Our version of controlled language comprises just five language statements that are orthogonal by design, which makes for a language that is suited for use by novice business rule modelers. The language lets users set up a business vocabulary that stakeholders can understand, and it allows to specify business rules about the objects in the vocabulary in a comprehensible if-then syntax. Rules expressed in our language are precise enough to permit the automatic generation of a prototype information system which is guaranteed to comply with the rules. Stakeholders can explore this prototype to verify the vocabulary, and to check whether the specified rules are valid and match their original intent of the business context. We show how we can ascertain correctness of our language and metamodel, by adopting a reflective approach and subject our context to rule analysis and specification, just like any other business context. It provides us with a prototype system that lets us explore the rules about rules, and validate the rule compliance.

1 INTRODUCTION

Business rules play an important role in day-to-day business operations and supportive IT applications. Declarative rules restrict what states are permitted in the business, and which operations may be executed by employees and information systems of that business (Hay, Healy 2000).

There is consensus that the business rules should be validated by stakeholders in the organization to ensure their overall correctness and coherence (Business Rules Manifesto 2003). Therefore, rules must be expressed in a way that a target business audience clearly understands. But to use those very rules in software applications calls for exact specifications and computer precision. This poses contradictory demands: comprehensibility for lay users, but perfect exactness for programmers and applications.

The prime deliverable of rule-based design is a compliant database application. In practice, the informal rules of business behaviour are rephrased and transformed in a chain of handovers until their encapsulation in an enterprise information system (figure 1). At each point in the chain, requirements are translated into yet another language, a process which is prone to misinterpretations, loss of detail, and other problems, even in the presence of a validated vocabulary (Bajwa et al., 2011).

![Figure 1: Chain of handovers of business rules.](image)

To reduce the need for translation, we propose a simple language founded on proven theory to cover the major part, if not the entire chain of handovers.

The business rules considered in this paper are declarative: there is no procedural dependence or hidden sequencing. The rules are also invariant: they concern persistent states only, not some transient situations that exist for just a brief moment in time,
e.g. only while a data transaction lasts. This differs from IT-approaches like the Event-Condition-Action (Poulouvasilis et al., 2003) paradigm, or Communicating Sequential Processes (Hoare, 1985; Wedemeijer, 2012). From a business point of view, the ECA type of rules have a technical ring, and their relevance is experienced as vague, difficult to retrace, and hard to explain (Andreescu, Mircea 2014).

The paper outline is as follows. Section 2 discusses some contemporary languages for declarative business rules, and design considerations for our language. Section 3 describes the proposed language. The syntax of each basic statement is depicted as a railroad diagram, and we explain core ideas. Section 4 puts the language to work, by describing features of supportive design- or prototype environments. Such an environment can be regarded as a business context having its invariant rules captured. In section 5 we pursue this idea by developing a meta-model of the language. Section 6 presents conclusions.

2 RELATED WORK

Business rule languages must be comprehensible for business workers on the one hand, and faultlessly exact for computer applications on the other (Bjekovic, Proper 2013). Numerous languages to express declarative business rules exist (Kardasis, Loucopoulos 2004). Our discussion of languages is restricted due to lack of space.

2.1 Declarative Rule Languages


These approaches rely on business vocabularies, also called ‘fact models’, in order to capture the true meanings and definitions of business data. Hence, comprehensibility and business focus is a strong point. However, controlled languages still permit a large variety in phrasing, and lack uniformity. As a result, rules are not always concisely and clearly expressed, making validation difficult and leaving room for interpretation, two fatal shortcomings for IT implementation (Weigand et al., 2011).

Other standards based on SBVR are FBM (FBM Working Group 2011) and Object-Role Modelling (Halpin, 2011). Both standards depict conceptual models in the customary way, and then depict the constraints visually. As a result, the diagrams with constraints become quite confusing, and they are barely intelligible for lay users.

A middle field is languages that aim to describe enterprise architectures, stakeholder concerns, goals and business rules (Quartel et al., 2009). Generally, these languages are not geared to capture rules, and are too high-level to allow validation by business stakeholders, or implementation in IT-systems.

On the other side of the spectrum are languages with an IT-provenance, such as UML- and XML-based languages or DTD's. Many of these languages are 'rich', meaning that a business feature may be captured in a variety of ways (Lamrani et al., 2013). Hence, it requires a thorough knowledge of implementation details to disclose the business relevance of an implemented rule (Beckner, 2014). Andreescu and Mircea (2014) remark on the reluctance to use OCL in the early design phases, when IT specialists need to cooperate with business people.

RuleML is an evolving family of XML-based languages (Boley et al., 2004). Semantic Web Rule Language, SWRL for short, achieves an expressive power superior to our language in some areas, e.g. to specify derivations, numeric and time calculations (Horrocks et al., 2004). SWRL also includes the Horn-clause syntax for rules, a strong point that which we will employ in our language. Nonetheless, the IT-orientation and notational complexity of SWRL, and XML-based languages in general, make them unsuitable for an average business user or novice designer (Akbari et al., 2103).

We conclude that (controlled) natural languages may capture business rules in a comprehensible and validatable manner, but not precise enough for computer applications. Formal rule modeling languages or general IT languages may be exact enough, but they lack in understandability.

2.2 Language Considerations

With the above in mind, our language for business rules was devised with the intention to:

- ensure comprehensibility for business people by relying on business vocabulary (terms and phrases of the business context).
- ensure that business workers can understand and validate their rules, and so minimize the need for back-and-forth translation of rules.
- ensure orthogonality of the language, so that
features are always expressed in just one way, and so avoid the problems of ‘rich’ languages.

- ensure exactness of rules, by founding them on rigorous mathematical theory; we opt for binary Relation Algebra (Maddux, 2006).

To prevent trivial but cumbersome errors in data entry, we prefer names and identifiers to be case-insensitive. Also, leading and trailing spaces should be avoided as much as possible.

2.3 Way of Working

Rule design may be conducted in a progressive way of working (figure 2).

![Figure 2: Way of working in rule-based systems design.](image)

The approach starts at business behaviour, which in most cases is only informally understood.

In the analysis and design phase, a business model of concepts and relations is created capturing the relevant parts of the business vocabulary. And, very important, the declarative rules are captured.

In the test environment, data for concepts and relations is loaded incrementally to test whether the predicted rule violations emerge. Rule enforcements are specified, determining how workflow processes should deal with rule violations in practice.

3 PROPOSED LANGUAGE

Our language provides five statement types that a designer may use in the specification of a business context. A railroad diagram of the overall language set-up is shown in figure 3.

Statements in a script may appear in arbitrary order, to suit an incremental, step-by-step, top-down, big-bang, modular, or any other preferred approach of the designer. The statements are uniform in make-up: first a reserved language imperative, identifying name(s) next, and then the further particulars.

3.1 Model

The model statement defines the concepts and binary relations that are part of the business vocabulary (figure 4). It sets up the structure of concepts and relations that the designer considers to be important.

![Figure 4: Railroad diagram for the model statement.](image)

Concepts have unique names, enclosed in square brackets for clarity, and starting with a letter. Relations are uniquely identified by a colloquial name to call the relation by, plus the names of its domain and range concepts. In addition to the colloquial name for the relation itself, another name may be provided for the inverse relation, indicated by the ~ inversion symbol. Uniqueness requirements for the relation name also apply to the inverse name.

We prefer the infix style of notation for relations. Technically speaking, prefix or other styles are equivalent. A script may contain multiple model statements, so that concepts and relations can be incrementally introduced. And because concepts are easily deduced from relation domains and ranges, a designer may even forego the explicit modelling of concepts.

By definition, it is impossible to violate a model. All true facts observed in the business context, either fit perfectly in the structure, or they are irrelevant. If some business fact is relevant but still cannot be
expressed as atoms or tuples, then the structure is wrong: it is an inadequate model of the business context. Thus, a model may be regarded as a set of structural rules. As no other rules apply to it, we call this an Unconstrained Conceptual Model.

3.2 Rule

We are now in a position to specify 'behavioural' rules that business stakeholders ought to live by. These rules should always evaluate to being satisfied, but in a running business, they may temporarily be violated. The implication is not, that the model is wrong. Rather, the business stakeholders should take action to remedy the violations.

The combination of model and rule statements is sufficient for business rule analysis and design. The joint deliverable may be called a Conceptual Model, and a good designer will make sure that it meets the usual quality requirements, such as completeness, and consistency of its rules (Moody, 2005).

To emphasize the behavioural aspect of rules, a core position is given to the rule keyword 'must' in the rule statement. The idea, in accordance with the ideas of RuleSpeak (2014), is to help users grasp the rule intent: guiding the business behaviour and have people refrain from violating the rule.

Each rule comes with a unique rule identifier, starting with a digit 0.9. Other statements can refer to the rule by way of this identifier, and it also comes in handy when violations are to be reported.

3.2.1 Simple Rules: Cardinality Constraints

A single relation may already be subject to a simple rule, i.e. cardinality constraints may apply. For and understandability, our language provides keywords to express cardinalities and common combinations.

For instance, the keyword 'function' means that a relation must be univalent and total. In addition, we provided keywords for rule types of homogeneous relations. And although simple rule types usually apply to simple relations, a compound expression may also be subjected to this kind of rule.

Combining disparate cardinalities under one heading defies the idea of having a unique identifier for each distinct rule. The designer should decide whether or not to combine rules, depending on how the user community understands these rules and deals with possible violations.

Notice that simple rule types are syntactic sugar: all simple constraints are perfectly expressible in mathematically equivalent compound rules. We include the rule keywords in our language for the sake of simplicity. In practice, it makes little difference: a rule is referenced only through its identifier, independent of the mathematical formulation.

3.2.2 Compound Rules: If-then Phrases

The real benefit of Relation Algebra is found in its ability to formulate compound rules in the concise yet straightforward way of normalized Horn clause format (Horrocks et al., 2004):

antecedent ⇒ consequent

Both the antecedent and consequent are binary relations, either a plain relation of the Unconstrained Conceptual Model or a compound expression, and must have the same concepts for domain and range. This format is easily translated to a semi-formal if-then sentence (1), into which we like to include the important rule keyword 'must':

IF the antecedent is confirmed, THEN MUST the consequent be confirmed (1)

We define a rule violation as: a pair in the antecedent, but absent from the consequent expression. With this definition, the text becomes:

IF a pair is present in the antecedent, THEN MUST the pair be in the consequent

The Horn-clause format easily pronounced as 'if... then must...', but the mathematical expressions in a rule may be quite complex, as seen in the railroad diagram of figure 5. Both expressions in the Horn clause may be a relation of the Unconstrained Conceptual Model, or may combine several relations using unary and binary operations. Special relations and constants may also be included, as explained below. It requires skills and business knowledge to translate the complex expressions into terms that the user community can understand. Better still is to
avoid complex expressions altogether, and find easy-
to-explain, natural rule assertions to begin with.

### 3.2.3 Special Relations and Constants

Complex expressions in rules may call for special relations. Our language provides a number of them, such as entire Cartesian Product, the empty relation, and the identity relation on a concept.

Expressions may also contain constant values or literals. Such values act as atoms or tuples in the rule expressions, but they need not be on record as they do not necessarily represent true business facts. Constant values in rules may force certain tuples to be on record. For instance, the rule 'the president of the USA must be a citizen', implicitly assumes that a nation named 'USA' is recorded. If we eliminate the constant by rephrasing the rule to 'the president of a nation must be a citizen of that nation', then the empty database no longer violates it. In general, an empty database never violates a rule if no constants are involved in the rule (Decker, Martinenghi 2006).

### 3.3 Explain

True business relevance means that each node, edge and clause in the specifications, can be clearly explained for, to, or even by the business workers. To help the audience grasp the detailed meaning and structure, explanatory texts are helpful (figure 6).

### 3.4 Load

An important means to put a model and rules to the test is by loading data and check for rule violations. The ability to load data is also useful when a design is demonstrated to the business stakeholders.

![Figure 7: Railroad diagram for the load statement.](image)

The load statement places sets of atoms, delimited by curly brackets `{ .. }`, into a concept extension, or sets of tuples in relation extensions (figure 7). Loading of data is not obligatory, but if data is loaded, then entity integrity and referential integrity is required (Date, 1981). At design time however, there is no need to worry about this, because it can be automatically ensured at load- and runtime.

### 3.4.1 Specifying Data

Like concepts, the atoms in our language are self-identifying: an atom is fully specified by its name, which is merely a text string, plus (the name of) the concept it belongs to. We do not distinguish between atom, atom-name, atom-value, or identity, distinctions that are hard to explain to lay users. Moreover, the atoms and their distinguishing names, or id's, or values must be linked tightly. Links that, on a meta-level, prove to be isomorphisms, or almost so. We think that such intricacies are better avoided.

Writing lots of data for loading is boring, and prone to typing errors. Our language provides two shortcuts. First, several relations may be loaded at once, provided of course that the domain and range are identical for all of them. Second, instead of specifying one tuple at a time, a set of tuples can be specified in one go, by combining a set of atoms from the domain with a set of atoms of the range.

### 3.5 Enforce

Enforcement is how runtime violations of the rules
should be dealt with from the business point of view. In realistic business environments, enforcement may range from 'avoid violation at all cost' to 'comply or explain' or even 'ignore all violations'. From the IT perspective however, there are just three main strategies called 'projector', 'rejector', and 'producer' (Dietz, 2008).

![Figure 8: Railroad diagram for the enforce statement.](image)

Our language provides for variants for all three variants. Figure 8 depicts the railroad diagram for the enforce statement in our language.

Enforce is not an incremental statement: a rule is subjected to one enforcement strategy at most, and specifying multiple enforcements for one rule has no use. If no enforcement option is specified for a rule, then the 'report' strategy applies by default.

### 3.5.1 Report

We call 'report' what the literature is referred to as 'projector'. This strategy for a rule means that after some edited data is committed, the rule is assessed, meaning that all violations are 'projected' into a special database table. Next, that listing is reported to the stakeholders. Notice that all violations ought to be reported, not just the new ones caused by the latest data edit.

Generally speaking, the 'report' strategy is easy to understand and robust to implement, which is why it is the safe choice for any business rule. For this reason too, this strategy is the default at load time.

### 3.5.2 Reject

The 'reject' type of enforcement strategy means that the rule must be checked prior to committing a change of data. If the change would result in a new violation for this rule, then the change should be rejected offhand, and the data not recorded.

The assumption underlying reject is that data violating this particular rule cannot even be valid in the business context. Rejection may be bothersome for business workers because this assumption is sometimes wrong, so that perfectly valid data is rejected for a bad reason.

There is another loophole: data may actually be in conflict with the rule, but if by coincidence the violating tuple is already present for another reason, then the erroneous data can be recorded nonetheless.

### 3.5.3 Resolve by

We introduce 'resolve' as final enforcement strategy for rules, referred to in the literature as 'producer'. The idea here is that sometimes in the business context, there is only one viable way to resolve a violation. And if the solution is known, why not let the computer apply it automatically?

We recall our definition of violation as a pair in the antecedent expression of a Horn clause, but not in the consequent. Hence, adding the offending pair into the consequent is a straightforward solution, and this is exactly what the enforcement strategy 'resolve by addition' intends. The strategy called 'resolve by deletion' takes the opposite tack and bluntly deletes the pair from the antecedent. As expressions in general cannot be edited, a necessary restriction is that the expression being edited must be a relation of the Unconstrained Conceptual Model.

A data-edit transaction is produced and in fact, this may result in the violation of some other rule. The automatic transition may even be rejected by another rule, or a subsequent transaction may be produced, and so on, potentially creating deadlocks or interminable loops. While compliant with the theory of Relation Algebra, this strategy goes beyond our context of invariant, i.e. state-oriented business rules. In defining our language, we did not investigate these effects nor have we tools for the rule designer to control them.

### 4 LANGUAGE ENVIRONMENT

The prime deliverable of rule design is a working database application that assures rule compliancy. Putting our rule language to work requires a design- and runtime environment in which all five of our statements contribute to that deliverable (figure 9).

As we strictly adhered to the sound theory of Rela-
tion Algebra, the generated prototype is guaranteed to comply with the invariant rules in the script.

4.1 Design Time Interface

A design time interface should support a designer to create, expand, refine and correct her script, and also to save the script to continue work at a later time.

A graphic display of the Unconstrained Conceptual Model, with drag-and-drop and rearrange features to unclutter the diagram, is a wonderful help in composing and understanding. Still, it is the model, the rule statements and explanations that should be at the core of the design effort, not the diagram.

The rule statement calls for a smart formula editor, with an option to link each rule expression to corresponding nodes or edges in the model diagram. Various flags would also be desirable, such as flags for faulty expressions, rule inconsistencies, potential simplifications of rules, or unconstrained relations.

The explain statement can well be supported by providing text editing functions in a mouse-over of the diagram.

The load statement needs generous support to let a designer include a full load of initial data in the script, and integrity should be taken care of automatically, at load time at the latest. Copy and paste of realistic data acquired from the business arena would be greatly appreciated. Automatic generation of datasets in compliance or violation of a specific rule, would also be desirable.

Lastly, no special support appears to be required for the enforce statement at design time.

4.2 Load Time Interface

At load time, all data in the script should be loaded into the Unconstrained Conceptual Model. Only the 'report' enforcement strategy is feasible at this time, to prevent undesired outcomes or even deadlocks caused by 'reject' or 'resolve' types of enforcement.

But data integrity must be made to hold. Referential integrity holds that each tuple refers to atoms on record in the domain and range, respectively. One option is to apply 'cascading delete', i.e. ignore all tuples that refer to a unrecorded atom. The opposite option is more attractive, i.e. to automatically insert the domain and range atoms of all tuples. Entity integrity holds that duplicates of an atom or tuple already on record, should not be loaded.

Once data is loaded into the Unconstrained Conceptual Model, only then should the rules be checked and violations reported. The report should enable the designer to trace each violation: what rule is violated, which atoms and tuples play a part.

A smart designer selects her data for loading in such a way that each violation is clearly understood, explained, and repaired. If a violation cannot be understood, then either the loaded atoms and tuples make no sense in the business, or the rule itself is in doubt. Or if particular violations can only be repaired by rigorously deleting data, then apparently some rules are contradictory. In any case, business people should be consulted to clarify the issue.

4.3 Runtime Interface

The script captures the rules of the business context, and a computer interface serves to confront the business community with their rules. Conveniently, it makes little difference how a designer organizes her statements in the script, because the business users are not exposed to the script directly.

A minimal requirement is a browse-and-explain interface. This should help users understand their exact business rules and violations. It ought to depict uncluttered diagrams of the entire Conceptual Model or parts thereof. It should display relevant explanations for all the concepts, relations and rules in the diagram. Also for each rule a complete list of all persistent violations must be provided, with explanations and traces how each violating tuple is determined from the corresponding Horn-clause formula. From there, the interface should support drill-down features to scrutinize partial populations or even individual instances in the diagrams.

Second, a demonstration interface is appreciated to emulate a workflow case in a (partial) business process. A series of predefined datasets is loaded in sequence, showing the emergence and subsequent resolution of violations as new data is being entered.

Adjusting rule enforcements on the fly in the runtime interface is still better. This would allow experimentation how to deal with rule violations and to probe the effects on the workflow processing.

The real benefit of our language is to generate a rule-compliant runtime prototype application with full data edit capabilities. Business people can then put that prototype to the question by entering all conceivable kinds of business data, view the responses by the system, and come to understand the workflow processes for dealing with the violations of their business rules.

5 LANGUAGE METAMODEL

So far, we discussed the modelling of an arbitrary
business context, its vocabulary and rules. In this section, we change the perspective and select 'rule design' for our business context.

5.1 Metamodel

What is the business vocabulary of design? What are its rules? If we could specify this special context in perfect detail, and if a compliant tool environment would be available, then we could use these... to generate the prototype system for rules design, in a truly reflective approach.

The idea, although not new (Schön, 1992), is still worthwhile to pursue. Having outlined the relevant business context in section 3, we present a conjectured metamodel in figure 10. In the diagram, we use a freehand style, and we omitted the meta-relation names. Interestingly, the statements of our language can be associated with five distinct areas in the metamodel.

Evidently, the metamodel comes with its own set of rules constraints on the metamodel concepts, associations, and contents. Without being exhaustive we outline some major rules per area.

5.2 Meta-rules and Enforcement

Entity integrity is an intrinsic demand of relation algebra, and it applies in this metamodel as well. Duplicate atoms of a concept or duplicate tuples in a metamodel association are unacceptable at any time. This integrity demand can be enforced as 'reject', in other words: duplicate entries are simply ignored.

In the 'model' area, it is compulsory that each relation has exactly one domain concept, one range concept. Referential integrity proclaims that both concepts must be present in the extension of the Concept concept. This rule is easily enforced as 'resolve by addition', meaning that missing concept names are automatically inserted. Moreover, the domain concept, range concept, and colloquial name together must uniquely identify the relation. And if a relation's inverse name is given, then it must adhere to the same uniqueness demand.

In the 'rule' area, a simple rule is associated with one expression and one (or more) cardinality- or homogeneous-ruletype. The compound rules, i.e. those in Horn clause format, are associated with both an antecedent and a consequent expression that must have the same domain and range concepts. Regarding the 'expression' concept, it must first be noticed that any particular expression may well be associated with several rules. Second, an expression can involve many relations, or concepts, or even constants, which is why we depict an non-specific line from 'expression' to the 'model' area in the metamodel diagram. Finally, it must be realized that expressions in general will constitute derived relations. That is, except in the special case where the expression equals a relation defined in the Unconstrained Conceptual Model.

In the 'explain' area, explanatory texts are associated to concepts, relations and rules, but no particular restrictions apply. In the 'enforce' area of the metamodel, only a uniqueness restriction applies.

5.3 Example of a Meta-rule

In the 'load' area, referential integrity must again be made to hold. As an example how analysis may be conducted and how our language supports the rule designer in her analysis, let us show how to capture referential integrity in a rule enforcement.

Referential integrity holds that for any tuple in any relation, both of its atoms must be on record:

\[
\text{IF the tuple } T \text{ has the domain atom } A, \quad \text{THEN MUST that tuple } T \text{ is contained-in some relation } R \text{ which has domain some concept } C \text{ which contains that atom } A.
\]

This controlled-language sentence translates to a rule in our language as an Horn-clause implication, with symbol \( \vdash \) indicating composition of relations:

\[
\text{rule 2-referential-integrity-domain as } \quad [\text{Tuple}] \text{ has-domain } [\text{Atom}] \text{ must imply } \quad [\text{Tuple}] \text{ contained-in } [\text{Relation}] \quad ; \quad [\text{Relation}] \text{ has-domain } [\text{Concept}] \quad ; \quad [\text{Concept}] \text{ -contained-in } [\text{Atom}]
\]

The antecedent expression is a relation in the Unconstrained Conceptual Metamodel. Hence, this rule assertion (2) permits us to impose the 'resolve by deletion' enforcement strategy: any tuple refer-
ring to an unrecorded atom is immediately deleted. This strategy is known as 'cascading delete' in relational database technology. A similar meta-rule will do for the atoms in the range of the relation.

The rule of integrity must hold at runtime, but still its enforcement strategy can be made to vary. A 'reject' or 'report' strategy may be better in a running business environment. Moreover, at load time the 'resolve by addition' strategy is to be preferred.

This illustrates how there are unresolved issues in the analysis and design of the metamodel calling for further research. Another issue is how to deal with constant values in rule expressions.

6 CONCLUSIONS

We proposed a rule language to capture and express declarative business rules. The language, combining business vocabulary with precise mathematical features, is comprehensible for business users; and precise enough to generate rule-compliant IT applications. We outlined how the language may be employed in design- and runtime interfaces.

6.1 Expressive Power of the Language

We claim that the language has adequate expressive power for rule design and analysis.

Following the ideas of SBVR, our language is founded on business vocabulary, compelling the designer to use the plain business phrases for the relevant terms and facts, a major strength of our language. Declarative, invariant business rules are described in a comprehensible if-then syntax. In our experience, this is a great help for people reading a script. In particular, the 'must' keyword provides an immediate clue of what a rule intends to say, even when complicated expressions are involved.

Our language has a clear, uniform makeup. This, and the simple naming regime make for easy-to-read scripts that are straightforward to interpret by business people, even without supportive IT-tools. Each line in a script starts with an imperative keyword that clearly indicates the focus of that line, underpinning the orthogonality of our language. The reserved keywords of our language are concise and learnable, appealing for both skilled business workers and novice rule designers.

Statements of our language are orthogonal by design. Each statement addresses a single aspect of the business context. The language statements are loosely coupled, but as some statements necessarily depend upon previous ones, complete independence is not possible.

No restrictions apply to the order or sequence of statements. The designer may first specify all aspects of one business feature, or start a model with a few rules in one section of the script and add a section with load statements later, etc. Therefore, no particular design approach is forced upon the designer. Having said that, a strong point of our language is that it does force a designer to consider all business features of the relevant rules, and to capture its aspects in distinct statements. For business users, it makes little difference how the statements in the script are organized, because in theory, users do not browse the script but use a dedicated interface to explore the rule-based design. In practice however, users will probably read it, and even begin to add and amend the script.

Our statements are devoid of typical IT jargon such as primary-keys and attributes, functional dependencies, cascading deletes, etc. Imperative ECA-type rules cannot be formulated in our language, with one exception. The enforce statement variant called 'resolve' does initiate data editing operations in response to a rule violation. This is a digression from the strictly declarative and invariant nature of our language, the consequences of which need to be further researched.

In our opinion, rule design for a business context is a superior approach than the dual approach of creating on the one hand an implementation data model with objects, entities, keys, and an activity model with data transitions and processing features on the other. Business stakeholders have little affinity with such refined IT-models, and lack the ability to validate the correct implementation in computer applications.

6.2 Language Extensions

Several extensions to the language may be considered to enhance usability for stakeholders and compliancy of the implementations. Of course, expressive power and understandability for business users should not be affected.

Support for specialization/generalization relations among concepts is one possible extension. This is somewhat problematic for Relation Algebra theory because an atom might belong to more than one concept for some time, or even switch over time: a person is student at one time, and teacher at another. Specialization/generalization is relatively unimportant in business practice, where models often do not need it or can use a work-around.

Better support for the 'resolve' enforcement
strategy is needed. Fundamental research is called for to understand the coordination of rules, and to prevent contradictory enforcement strategies.

Support for the Role-based Access standard (Edward et al., 2011) is also suggested. Instances of a Role concept should be assigned the right to access all contents of certain concepts and relations in the Unconstrained Conceptual Model. It calls for a mix of model and metamodel capabilities, thus extending the ideas exposed in section 5.

A serious shortcoming of Relation Algebra is that it lacks arithmetic and temporal capabilities: it cannot express calculations such as 'add 18% VAT' or comparisons like 'if born before 1980'. Support for this kind of rules will greatly enhance usability, provided that the orthogonality of the language and most importantly, the clear and uniform expression of declarative rules in if-then syntax is safeguarded.

### 6.3 Future Research

We indicate some areas of ongoing research that may improve the applicability of our approach, methods, and tools for business rules design.

Currently, texts available for explanations in the user interface are only static. Ongoing research aims to determine what instructions or explanations in which interfaces are most helpful to achieve high-quality designs (Michels, 2011).

Integration of our language with the typical IT-domain of Semantic Web Rule Language is being researched (Grosof, 2013). The aim is to improve the expressive power without compromising orthogonality of the language and comprehensibility of the if-then syntax of rules.

Interface design is an ongoing area of research. In this paper, we proposed to compose and then compile scripts. But instead of compiling, an interpretative way of working might provide better support for the designer and business stakeholders.

Research is being conducted to develop the reflective meta-modelling approach, its vocabulary and the rules of rule design. The idea is to build a generator from this; a generator that is capable of converting any rule-based design into a fully functional and compliant prototype application.

The vision is that in the future, stakeholders may formulate and validate their own business rules, and do so in a language with enough precision to enable a straightforward implementation in computer applications, without the intervention of IT specialists.

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