Propagating Model Refactorings to Graph Transformation Rules

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Abstract: Model-driven software engineering reduces the effort of developing software by replacing low-level programming with the construction of high-level executable models. Refactoring improves the structure of software artifacts without changing external behavior. Originally, refactoring was developed for and applied to (object-oriented) programs. In the context of model-driven software engineering, refactoring has to be applied to both structural and behavioral models. In this paper, we present tool support for model refactoring in ModGraph, a tool which employs Ecore class diagrams for structural modeling and graph transformation rules for behavioral modeling. In particular, we focus on the propagation of refactorings of the structural model into the behavioral model.

1 INTRODUCTION

Model-driven software engineering reduces the effort of developing software by replacing low-level programming with the construction of high-level executable models. To this end, both structural and behavioral models have to be developed. In the context of object-oriented modeling, there seems to be a general consensus to employ some variant of class diagrams for structural modeling, e.g., EMF, MOF, or UML class diagrams. In contrast, there is a wide spectrum of languages for behavioral modeling which are based on different computational paradigms (e.g., state machines, activity diagrams, or rule-based transformation languages). In this paper, we will focus on behavioral modeling with graph transformation rules: Models are considered as graphs, and transformations of these graphs are specified declaratively by graph transformation rules.

Software evolution is a discipline which provides concepts, methods, and tools for evolving software in response to changing requirements, platforms, technologies, etc. In the context of object-oriented software development, evolution support is provided in a variety of different ways, including e.g. design patterns and refactorings, both of which facilitate changes to the software. According to (Fowler, 1999), refactoring denotes the activity of restructuring software by applying a series of transformations without affecting its externally observable behavior.

The transformations proposed in (Fowler, 1999) were developed for and applied to (object-oriented) programs. In the context of model-driven software engineering, refactoring has to be applied to both structural and behavioral models. Previous work on model refactoring focused on structural models (Biermann et al., 2006; Mens, 2005; Mens et al., 2007). However, when the structural model is refactored, the respective changes have to be propagated into the behavioral model; otherwise, the behavioral model is no longer consistent with the structural model.

In this paper, we close the gap identified above, resulting in comprehensive support for model refactoring. We present tool support for model refactoring in ModGraph, a tool which employs Ecore class diagrams for structural modeling and graph transformation rules for behavioral modeling. For the structural model, ModGraph offers a set of refactoring transformations along the lines of Fowler’s work (Fowler, 1999). The tool support goes beyond previous work since the refactorings are propagated from the structural model into the behavioral model. Thus, graph transformation rules are updated in response to the changes of the underlying Ecore model.

Altogether, our work provides an important contribution to model evolution, which is an essential prerequisite for putting model-driven software engineering to work. The term model evolution is used in a variety of different contexts. Frequently, model evolution is concerned with the changes of model instances...
in response to changes of the underlying metamodel (Rose et al., 2010) (analogously to schema evolution in databases (Banerjee et al., 1987)). In contrast, the work presented in this paper has a different focus: It deals with the consistent refactoring of a set of interdependent models, which requires propagation of the changes of the structural model into the behavioral model. The migration of model instances is not considered here.

The rest of this paper is structured as follows: Section 2 explains the context in which this work was performed. Section 3 provides an overview of our approach to refactoring both structural and behavioral models. Section 4 describes the implementation. Section 5 illustrates this approach by several examples. Section 6 discusses related work. Section 7 concludes the paper.

2 BACKGROUND

The work reported in this paper was carried out in the context of the ModGraph project (Winetzhammer, 2012). ModGraph is a tool for model-driven software engineering which is based on the Eclipse Modeling Framework (EMF, see (Steinberg et al., 2009)). Structural modeling is performed with Ecore. Thus, a structural model consists of a set of classes owning structural and behavioral features. Structural features are partitioned into attributes, which have (collections of) simple values, and references to target classes. By default, references are uni-directional (from instances of the source class to instances of the target class); however, a pair of uni-directional references may be grouped into a bi-directional reference. Behavioral features are modeled by operations.

The structural model merely describes the signature of operations, but not their behavior. In ModGraph, the behavior of an operation may be specified declaratively by a graph transformation rule. Model instances are considered as graphs whose nodes and edges correspond to objects and links, respectively.

A graph transformation rule describes an in-place model transformation and essentially consists of a left-hand side and a right-hand side, like the one shown on the left-hand side of Figure 1 to move an attribute from one class to another. The left-hand side describes the subgraph to be searched (the match for the rule). Here the classes and the attribute that should be moved are located. The right-hand side specifies the subgraph which replaces the left-hand side. In case of Figure 1 the attribute references the other class. The left-hand side and the right-hand side may be merged into a single diagram, as shown in Figure 1 on the right-hand side. Elements belonging only to one side need to be marked: only left-hand side elements are colored red and marked --, only right-hand side elements are colored green and marked ++.

ModGraph uses the merged view. An example is given in Figure 4. In ModGraph single and double boxes represent single objects and sets of objects, respectively. The current object on which the operation implemented by the graph transformation rule is invoked is designated by the inscription this. Arrows represent links, i.e., instances of references. Conditions on and modifications of attribute values are shown in the Constraints and Changes compartments of the boxes for objects. Nodes designated as optional are not required for the rule to match. According to the merged view, insertions and deletions of objects and links are represented by coloring the respective graphical elements in green and red, respectively, as well as markers ++ and -- are placed inside the boxes for objects and near to the arrows for links. The meaning of the graph transformation rule shown in Figure 4 will be described later.

Altogether, the ModGraph tool has to deal with different types of models, as illustrated in Figure 2. The blue boxes represent the models which are created by the ModGraph user. The structural model is defined by an Ecore class diagram. The behavioral model consists of a set of graph transformation rules specifying the behavior of operations introduced in the structural model. At runtime, graph transformation rules will be applied to model instances by matching and replacing subgraphs. Internally, a graph transformation rule is represented as a model, too (namely an instance of the graph transformation metamodel, which in turn is an instance of the Ecore metamodel).
According to (Fowler, 1999), refactoring denotes the activity of restructuring software by applying a series of transformations without affecting its externally observable behavior. Fowler’s book contains a bunch of refactoring transformations, all of which refer to the source code. For each refactoring, the following information is supplied: the precondition under which the refactoring may be applied, the elementary steps of the core refactoring, the impact of changes on dependent model elements, and the postconditions which hold after the refactoring. The refactoring transformations vary in their complexity and range from simple ones such as renaming an attribute to complex refactorings such as pulling up an attribute from all subclasses to a common superclass and extracting a class from a given class by moving a part of its features to the extracted class.

In this paper, we discuss refactorings of models rather than refactoring of source code. More specifically, we focus on models created by ModGraph users (blue boxes in Figure 2). Thus, we consider the following problem: Let us assume that a consistent model is given, being composed of a class diagram and a set of graph transformation rules. The graph transformation rules are typed with the elements of the structural model. ModGraph ensures that typing is consistent: The class of an object must be defined in the class diagram, the attributes occurring in the object’s compartments must be owned by that class, etc. Now, the ModGraph user would like to apply a refactoring transformation to the overall model. To this end, the structural model has to be changed and the changes must be propagated into the graph transformation rules such that consistency is re-established.

3 PRELIMINARY CONSIDERATIONS

As the last section described EMF and ModGraph graph transformation rules as well as their interplay, this one focuses on problems concerning the graph transformation rules co-occurring within refactoring the EMF model. Figure 2 shows the meta-relationships between the models (blue, rounded arrows), as well as the interplay between EMF and ModGraph (black, bold arrows). ModGraph is deeply rooted in the EMF world, taking a look at the meta-relationships between the models. Each graph transformation rule (Figure 2, blue box on the right) is an instance of the graph transformation meta model (orange box on the right) which was built with EMF using the Ecore metamodel (green box on the left). This metamodel also serves as metamodel for Ecore class diagrams (blue box on the left). Those class diagrams may be instantiated (black box on the left).

The graph transformation rule interacts with the model and the instance. The first interaction is caused by the fact that each graph transformation rule implements an operation in the Ecore class diagram. Therefore it references this operation as well as its containing class. The second interaction takes place inside the rule: A ModGraph rule changes instances of the Ecore model by matching structures and applying changes. The components inside the rule reference elements of the structural model, e.g. an object in a rule must reference an EClass inside the Ecore model, a link has to reference an EReference.

These two dependencies are shown as black, bold arrows named "references" and "matches / transforms" in Figure 2. A full overview of the dependen-
ties between a rule and the structural model is given in (Buchmann et al., 2012) with an explanation of the ModGraph metamodel for graph transformation rules.

Within a static context, this kind of model execution will probably never cause any problems. But many models evolve in time: elements are renamed, references are re-targeted and classes are inlined or extracted in order to create or collapse inheritance hierarchies or just to improve the model - in short: the model is refactored. Regarding only Ecore models, applying these refactorings is well understood and implemented in several ways, as we will show in Section 6. The problem is caused by the strong interconnection between a rule and the Ecore model. If any element is just referencing the refactored model element, after refactoring, this reference is unset. Hence, the rule is invalid because of inconsistencies between the rule and the model, which will be eliminated by our approach. We propose to propagate refactorings to rules, with the help of rules, as shown in Figure 3. We use two steps: First we call the modeled refactoring rules, second we propagate the refactorings performed by these rules to the graph transformation rules belonging to the refactored static model by using mostly meta-rules. That way we provide refactoring on graph transformation rules executed by other graph transformation rules, we call meta-rules.

Our approach starts with Fowler’s definition of each refactoring. As these definitions were introduced for the refactoring of (object oriented) code, we adapt the definitions for Ecore model refactorings. Subsequently we investigate how the changes affect ModGraph rules. Here we consider different aspects concerning the structure of the rule. Depending on the element the changes may cause minor or major changes to be performed inside the rule. Table 1 shows some. One may classify the affected rule elements as shown in the first column: the whole rule, nodes (representing single- and multi-valued objects and parameters), attributes, links (instances of references), and textual conditions (e.g. pre- and postconditions, attribute conditions), as well as fields (e.g. operation calls). The result of this step is taken to propagate the changes to the rules. The change may be a simple setter call on the graph transformation rule or restructuring of the rule itself, as Section 5 will show.

The rule as a whole references an operation and its containing class. Hence changes performed here need to be propagated. A rule is superseded if the operation is deleted. If the class is deleted, the rule may be assigned to another class where a similar operation shall take place now or is also deleted. A node references a class. It needs to be updated whenever the class changes. If for example the class is renamed, the node needs to be retyped. If the class is made abstract, the node will be deleted if its status is CREATE. For any other status it will be retyped. Attributes are an essential part of a rule. They may be used to add conditions to nodes as well as for assignments. Each attribute in a rule references an attribute in the Ecore model. Therefore an attribute in a rule needs to follow the attribute in the class. Renaming and moving may be executed without problems, but retyping is quite interesting: If an attribute is retyped in the Ecore model and used inside a graph transformation rule the user has to assign its new value manually. Furthermore attributes may be moved and links be re-targeted. In both cases the refactoring tries to locate an instance of the new container / target of the attribute / link and to move / re-target the attribute / link, respectively. If this fails the attribute / link is deleted. For the textual parts of the rules exists no other way than parsing them to find the elements referred and to adapt them using naming conventions.
Table 1: Effects of model changes on different classes of rule elements.

<table>
<thead>
<tr>
<th>Affected rule element</th>
<th>Change in Ecore model</th>
<th>Change modeled in meta-rule</th>
</tr>
</thead>
<tbody>
<tr>
<td>Whole rule</td>
<td>Class rename</td>
<td>adapt references to class</td>
</tr>
<tr>
<td></td>
<td>delete</td>
<td>adapt references to class</td>
</tr>
<tr>
<td></td>
<td>make abstract</td>
<td>adapt references to class</td>
</tr>
<tr>
<td></td>
<td>Operation rename</td>
<td>adapt references to operation</td>
</tr>
<tr>
<td></td>
<td>delete</td>
<td>delete</td>
</tr>
<tr>
<td>Nodes</td>
<td>Class rename</td>
<td>retype</td>
</tr>
<tr>
<td></td>
<td>delete</td>
<td>delete</td>
</tr>
<tr>
<td></td>
<td>make abstract</td>
<td>retype or delete</td>
</tr>
<tr>
<td>Attributes</td>
<td>Attribute in Class rename</td>
<td>adapt name</td>
</tr>
<tr>
<td></td>
<td>move</td>
<td>try to move, if not successful: delete</td>
</tr>
<tr>
<td></td>
<td>retype</td>
<td>retype and assign new value delete</td>
</tr>
<tr>
<td></td>
<td>delete</td>
<td>delete</td>
</tr>
<tr>
<td>Links</td>
<td>Reference rename</td>
<td>adapt name</td>
</tr>
<tr>
<td></td>
<td>retarget</td>
<td>try to retarget, if not successful: delete</td>
</tr>
<tr>
<td></td>
<td>delete</td>
<td>delete</td>
</tr>
<tr>
<td>Textual conditions,</td>
<td>any change</td>
<td>parse condition and propagate re-factor in the model. They also work for Ecore models outside the ModGraph context.</td>
</tr>
<tr>
<td>paths and fields</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

4 IMPLEMENTATION

The refactoring is implemented using ModGraph/Xcore\(^2\) and Java to implement connectors using extension points. For the static structure, we mainly use two classes in an Xcore model: Refactoring, to specify the refactorings on the model (as started in (Dümmel, 2013)), and Propagation, to propagate the changes caused by them.

As we want to use the added value of rules only, each refactoring on the model is specified inside a graph transformation rule or an Xcore operation, depending on its complexity. We will see in detail in Section 5, that the change of a bi- to a unidirectional reference is modeled very well with rules, while renaming an element in the model can be written easily in a few lines inside an Xcore operation. Please note, that these refactorings are universal. They also work for Ecore models outside the ModGraph context.

The propagation of changes to ModGraph rules follows the same scheme. A complex propagation is modeled as a rule, using Xcore as control flow language, if necessary. A simple propagation is modeled as an Xcore operation. In Section 5, examples for propagations are shown.

As a final step of propagation, all rules changed are validated in order to secure their navigability. Navigability means, that each single- or multi-valued object needs to be reachable by navigating links starting from the current object or a single- or multi-valued parameter of the operation. If the validation fails the user needs to adjust the rule manually. That way, rules which are no more navigable after applying the refactoring will be marked for further manual changes.

The rest is implemented in Java: We extend a popup-menu in order to be able to call the refactoring in the model and call the refactoring rule implementation (written in ModGraph/Xcore). After successfully executing this one, we call the propagation implementation (written in ModGraph/Xcore), that propagates the changes to the rules. The propagation locates the rules, parses and fixes them as shown in the next paragraph. The result is a complete model evolution.

We implemented the following refactorings and their propagations. A validation of the rules is performed after each refactoring.

**Rename Element:** Renames an ENamedElement. Propagates the new name to the corresponding el-

\(^2\)The approach to use ModGraph and Xcore is described in (Winetzhammer and Westfechtel, 2013) and (Winetzhammer and Westfechtel, 2014).
Add Parameter: Adds an EParameter to an EOperation. Propagate to implemented operation and textual fields, especially to operation calls, in a rule. Both need to be adapted.

Remove Parameter: Removes an EParameter from an EOperation. Propagation: Removes the parameter from the rules, updates textual fields, especially operation calls.

Change Value to Reference: Updates the containment hierarchy. Propagation: Validate rules in order to check exclusiveness.

Change Reference to Value: Updates the containment hierarchy. No propagation.

Change Bi- to Unidirectional Reference: Removes an EReference. Propagation: Exchanges each link targeting the deleted reference by a link targeting its former opposite reference. (See Section 5)

Change Uni- to Bidirectional Reference: Adds an EReference and sets its EOpposite. No propagation.

Extract Class: Extracts an EClass from a given EClass and relates them using an EReference. Propagation: If a node references the given class, check if it uses attributes, operations or references moved to the extracted class. If yes, create a new node referencing the extracted class and relate both using a link which is an instance of the EReference relating the classes. Move the elements referencing the extracted class to the newly created node.

Extract Subclass: Extracts a subclass from a given class. Propagation: If operations, attributes or references are moved, retype the nodes referencing the class and using one of the operations, attributes or references to be typed over the subclass.

Extract Superclass: Extracts a superclass from a given class. No propagation, because of inheritance.

Inline Class: Inlines a class into another class. Propagation: Retype all nodes typed over the inlined class to the class in which it will be inlined.

Move Attribute / Operation: Moves the attribute / operation from a source to a target class. Propagation: If the rule contains two nodes referencing the source and target class move the usage of the attribute / operation. In all other cases delete the usage.

Pull Up Attribute / Operation: Moves the attribute / operation to the super class. No propagation, because of inheritance.

Push Down Attribute / Operation: Moves the attribute / operation to the subclasses. Propagation: If there is a node typed over the super class using the pulled down attribute / operation, ask user to which subclass it shall be retyped and perform this.

Replace Subclasses with Fields: Creates an enumeration with one literal for each subclass and an attribute typed over the enum in the superclass. Propagation: Replace the nodes typed over the subclasses by nodes typed over the superclass with the corresponding literal assigned.

Hide Delegate: Removes the direct EReference between two EClasses. Propagation: Delete corresponding links in rules.

Remove Middleman: Creates a direct EReference between two EClasses. Propagation: Create the corresponding link, if the rule contains two nodes referencing the classes.

Replace Inheritance with Delegation: Replace the existing generalization with an EReference. Propagation: If both classes are referenced by nodes inside a rule create a link referencing the EReference.

Replace Delegation with Inheritance: Replace an existing EReference with a generalization. Propagation: Remove the link referencing the EReference.

5 EXAMPLES

The general idea of propagating model refactorings on interdependent models is now concretely shown considering the refactoring change a bi-directional to a uni-directional reference in detail. The refactoring to rename an element is sketched by the special case of renaming a class.

5.1 Change a Bi- to a Uni-directional Reference

The refactoring we consider here in detail is changing a bi- to a uni-directional reference. Starting with Fowler’s definition (Fowler, 1999), we provide the implementation of the model change as well as the meta-rule to propagate the changes.

Fowler defines this refactoring as follows: "You have a two-way association but one class no longer
needs features from the other.” (Fowler, 1999) For EMF this means, that a reference with an opposite set is deleted. The rule to perform the operation in the Ecore model is shown in Figure 4. Here a reference called reference2 (red box, marked with --) is deleted within all its links (red arrows, marked with --) to classes as well as its opposite relation to the other reference, called formerOpposite.

To propagate this refactoring, we use an Xcore operation and two meta-rules. The meta-rule shown in Figure 5 shows the application of the refactoring propagation to links in the Graph Pattern. The
Figure 6: Graph transformation rule before and after in-voking the refactoring change bi- to uni-directional refer-
tence on the model. toS and toP are instances of the equally
named references in the model, which are opposed.

The meta-rule for the negative application conditions is
structured analogously and therefore not shown here.
These two meta-rules are applied to all rules belonging
to the refactored model using an Xcore implemented operation. First, the operation loads the whole
rule, second it calls both meta-rules, third it updates
the diagram for the rules and finally it saves the mod-
ified resource.

In the meta-rule the deleted reference is located.
Each link, called delLink in Figure 5, is referencing
the deleted reference and is therefore deleted. A new
link, called link in Figure 5, is created as an instance
of the former opposite of the deleted reference. The
numbers next to the paths in the rule represent OCL
path expressions, shown in Listing 1. They ensure
consistency for pattern matching. No. 1 and 2 en-
sure that the nodes node1 and node2 are referencing
the classes between which the references has been
deleted. No. 3 and 4 ensure that a link, referencing
the deleted reference exists in the rule.

The refactoring is applied to a test project. The
model contains a bidirectional reference modeled as
two EReferences called toS and toP, which are set as
EOpposite to each other. The project also contains a
set of graph transformation rules. One of the rules is
shown in the upper part of Figure 6, where an instance
of a parameter p shall be deleted. This instance is con-
ected using a link, an instance of the toP reference, to
the current object. Now let us assume that the user de-
cides toP is no more needed in the class diagram and
therefore calls our refactoring rule to change the bi-
to a uni-directional reference. toP is removed from the
structural model. Also our propagation mechanism changes the rules. The result for our example rule is
shown on the lower part of Figure 6.

5.2 Rename Element
Renaming an element is a quite common task in pro-
gramming and modeling ever since. This is an inter-
esting example, as the refactoring for the structural
model can be written as a few lines in an Xcore oper-
ation, but the changes to be propagated to the rules
are not as trivial. As this refactoring may affect all
components of an Ecore model, all components of
a rule referencing those may be concerned. Due to
the fact, that elements of a rule reference model el-
ements differently, there are a number of renaming
propagations. Starting with the simple part, Listing 2
shows the implementation for the refactoring of the
structural model. The renamed element as well as
its new name are given as a parameter to the opera-
tion. First our implementation checks if the element’s
container already contains any element with the given
new name. If not, the element is renamed.

The rather complex part is the propagation that is
decomposed into several meta-rules and Xcore opera-
tions. For the meta-rules we distinguish between each
element that is referred in the class diagram. Hence,
we need at least one rule or Xcore operation to prop-
agate the renaming of a class, an attribute, an opera-
tion, or a reference.

As an example we consider the propagation of re-
naming a class in the Ecore model. The first step is to
rename the class in the model itself. This is performed
by a call for a setter. In the next step all graph trans-
formation rules are searched using the criteria shown
in Table 2 in the left column. If one of these ques-
tions is answered positively, the rule needs to be ad-
justed. Depending on which answers were positives,
the changes in the right column are applied to the rule
in order to get a consistent set of models. Assume
the answer to the first question is yes. The next ques-
tion is also answered positively, if the rule contains
a current object. Figure 7 shows the meta-rule for
the propagation. The corresponding class of the rule
is set to the renamed one, the current object this is
retyped and possible unbound nodes referring to the
class are also retyped directly. Links as instances of
references typed over the renamed class are adapted
in an extra meta-rule, which is called by an operation
call in the current object of the meta-rule shown here.
The same holds for the operations resetImplmented-
doOperation and handleInnerNodes. The first one cor-
rects the implemented operation’s path, the latter han-
dles attributes in unbound nodes. Textual elements
are handled separately. Please note that retyping by
renewing the type / corresponding class is necessary
here, as these are inter-model dependencies which are
lost by changing the corresponding object.
Listing 1: OCL path expressions for change bi- to unidirectional reference.

1: \[
\text{exists}(e \mid e.oclIsTypeOf(GTThisObject) \land\ e.oclAsType(GTThisObject) >
\text{exists}(e \mid e.type.oclAsType(ecore::EClass).name =\n\text{formerOpposite.eReferenceType.oclAsType(ecore::EClass).name})
\lor (\text{self}.oclIsTypeOf(GPUnboundNode) \land\ self.oclAsType(GPUnboundNode) >
\text{exists}(e \mid e.type.oclAsType(ecore::EClass).name =\n\text{formerOpposite.eReferenceType.oclAsType(ecore::EClass).name}))
\lor (self.oclIsTypeOf(GPUnboundNode) \land\ self.oclAsType(GPUnboundNode) >
\text{exists}(e \mid e.type.oclAsType(ecore::EClass).name =
\text{formerOpposite.eReferenceType.oclAsType(ecore::EClass).name}))
\]

2: \[
\text{exists}(e \mid e.oclIsTypeOf(GTThisObject) \land\ e.oclAsType(GTThisObject) >
\text{exists}(e \mid e.type.oclAsType(ecore::EClass).name =\n\text{formerOpposite.eContainer.oclAsType(ecore::EClass).name})
\lor (\text{self}.oclIsTypeOf(GPUnboundNode) \land\ self.oclAsType(GPUnboundNode) >
\text{exists}(e \mid e.type.oclAsType(ecore::EClass).name =\n\text{formerOpposite.eContainer.oclAsType(ecore::EClass).name}))
\lor (self.oclIsTypeOf(GPUnboundNode) \land\ self.oclAsType(GPUnboundNode) >
\text{exists}(e \mid e.type.oclAsType(ecore::EClass).name =
\text{formerOpposite.eContainer.oclAsType(ecore::EClass).name}))
\]

3: \[
\text{select}(n \mid n.oclIsKindOf(GTNode) \land\ n.oclAsType(GTNode).outgoingEdges >
\text{exists}(e \mid e.oclIsKindOf/GPLink) \land\ e.oclAsType/GPLink.eReference.oclIsProxy()))
\]

4: \[
\text{select}(n \mid n.oclIsKindOf(GTNode) \land\ n.oclAsType(GTNode).incomingEdges >
\text{exists}(e \mid e.oclIsKindOf/GPLink) \land\ e.oclAsType/GPLink.eReference.oclIsProxy()))
\]

Listing 2: Xcore implementation of refactoring rename element.

```xcore
op void renameENamedElement(ENamedElement element, String newName) {
    if (element.eContainer.eContents.filter(e | e.eClass == EClass) .exists(e | 
        e.eClass == element.eClass && (e as EClass).name == newName) ) {
        return /*and show message */
    } else {
        element.name = newName
    }
}
```

Table 2: Strategy for the rename element refactoring when renaming a class in the structural model.

<table>
<thead>
<tr>
<th>Does the rule reference this class?</th>
<th>Change the reference to the class and retype the current object.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Does the rule contain nodes referencing this class?</td>
<td>Retype the node.</td>
</tr>
<tr>
<td>Do those nodes contain inner nodes referencing contained elements of this class?</td>
<td>Retype them.</td>
</tr>
<tr>
<td>Does the rule contain links referencing EReferences typed over the class?</td>
<td>Retype links.</td>
</tr>
<tr>
<td>Does the rule contain textual elements making use of this class?</td>
<td>Get the new name and replace the old one.</td>
</tr>
</tbody>
</table>

6 RELATED WORK

Mens (Mens and Tourwé, 2004) provides a comprehensive survey of software refactoring. Fowler’s book probably constitutes the most cited reference in this domain (Fowler, 1999). The refactorings presented in this book all apply to object-oriented programs, i.e., they are applied to source code rather than models. Furthermore, refactorings are described in an informal way. A formalization which is based on program graphs and graph transformation rules is presented in (Mens et al., 2005). A major restriction of this work consists in the fact that only individual rules are considered. In the general case, programmed graph transformation rules are required for specifying refactoring transformations. This is demonstrated e.g. by (Geiger, 2008), in which a refactoring case prepared for the GraBaTs 2008 workshop was realized in Fujaba (Zündorf, 2001). The refactoring case included three refactorings on program graphs.
Research on model refactoring primarily focuses on the structural model (class diagrams). For example, in (Biermann et al., 2006) refactoring of Ecore models is specified with graph transformation rules in the AGG environment (Taentzer, 2003). Mens (Mens, 2005) demonstrates how refactoring transformations on UML class diagrams may be specified in different graph transformation languages (AGG and Fujaba). In (Mens et al., 2007), critical pair analysis is applied to AGG rules for refactoring Ecore models in order to detect dependencies between refactoring transformations. Based on this analysis, the user is guided in the application of these transformations.

Bottoni (Bottoni et al., 2003) goes beyond these approaches by providing integrated refactoring transformations: A refactoring transformation is applied not only to a UML class diagram, but also to the source code implementing the structural model. Our work differs from this approach inasmuch as we consider behavioral models rather than source code.

We are aware of only a few approaches dealing with the propagation of changes from the structural model into the behavioral model. Rosner and Bauer (Roser and Bauer, 2008) propose an approach to update model transformations in response to metamodel changes. The approach requires an ontology mapping between metamodel versions and is applied to evolve QVT-R (OMG, 2011) transformations. The evolution of model transformations constitutes an example of a higher order transformation (Tisi et al., 2009).

To the best of our knowledge, the only approach dealing with the propagation of changes into graph transformation rules is presented in (Levendovszky et al., 2009), which refers to the GReAT language and environment (Agrawal et al., 2006). However, this approach suffers from several limitations. First, it considers only elementary changes to the metamodel, e.g., changing the name of a class. Second, changes are propagated only in a semi-automatic way. Third, the resulting graph transformation rules may contain syntactic and semantic errors.

Altogether, the work presented in this paper is unique inasmuch as it does not only support refactoring with, but also refactoring for graph transformation rules. Refactoring on the structural and the behavioral model are supported in an integrated way, and changes are propagated to graph transformation rules such that consistency is preserved.
7 CONCLUSION

Model-driven software engineering reduces development effort by replacing low-level programming with the construction of high-level models. To make these models executable, structural modeling has to be complemented with behavioral modeling. During their lifetime, models undergo many changes for a variety of different reasons. Thus, it is crucial to support model evolution. Refactoring of models provides an important contribution to model evolution since it aims at restructuring models such that future changes are facilitated.

In this paper, we presented tool support for model refactoring in the ModGraph environment. ModGraph employs Ecore models for structural modeling and graph transformation rules for behavioral modeling. Refactorings are supported in an integrated way: Each refactoring transformation on the structural model is consistently propagated into the behavioral model. In this way, our work goes considerably beyond previous work on model refactoring which was confined to the refactoring of the structural model. Furthermore, while several other approaches use graph transformations for refactoring (i.e., refactoring with graph transformations), our work is unique inasmuch as it addresses refactoring for graph transformations, as well.

REFERENCES


