Combining Supervisory Control, Object-oriented Petri-Nets and 3D Simulation for Hybrid Simulation Systems using a Flexible Meta Data Approach

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Abstract: The idea of Supervisory Control is to regard a control component (Supervisor) as a discreet event simulation. State Oriented Modeling combines the ideas of supervisory control and object-oriented Petri-nets. With these concepts as a starting point, our goal was to realize a truly hybrid simulation system, which allows the simultaneous use of discreet event simulation and continuous 3D-simulation on a unified database. The key component is an active real-time simulation database, which is an object-oriented, self-reflecting graph database, with a powerful meta-information system. All nodes are derived from a common base class and data is stored in properties with standardized getter and setter functions. The object-oriented Petri-nets are formally described in the State Oriented Modeling Language, which is itself an extension scheme of the simulation database.

1 INTRODUCTION

State Oriented Modeling (Schluse, 2002) combines the ideas of supervisory control introduced by (Ramadge and Wonham, 1984) and object-oriented Petri-nets (OPN) (Bastide, 1995). It has already been used for a large variety of different applications in the field of simulation (e. g. to simulate robot programs as described in (Baldini et al., 2005)), but also for the real-time control of physical systems using simulation technology as described in (Rossmann et al., 2008). To realize these applications, Supervisory Control provides the methods necessary to link the control algorithms with simulations or physical devices. To implement the controllers, Petri-nets are known to be able to map almost all of the most important state oriented description languages and even modern programming paradigms, to model complex scenarios.

Figure 1 illustrates the integration of Supervisory Control concepts, Petri-Nets and 3D simulation used so far. Here the simulation system and the supervisory control are separate entities, working on disjoint sets of data. Communication between these two sets has to be established via a mediator.

With this concept as a starting point, our goal was to realize a truly hybrid simulation system, which allows the simultaneous use of discreet event simulation and continuous 3D simulation on a unified database. It is then up to the simulation developer how to use these paradigms in parallel to realize convincing simulation applications in a wide range of application areas from "classical" simulation applications (driving simulators, virtual production, etc.), to new application areas like user interface design or Virtual Testbeds providing simulation-based development frameworks for complex systems, a key technology in the emerging field of eRobotics (Rossmann and Schluse, 2011).

This paper will detail the progress we have made in making State Oriented Modeling not only an add-on to a 3D simulation system, but incorporating the
principles directly into our real time simulation system database.

The key component is an active real-time simulation database, which is an object-oriented, self-reflecting graph database (as shown in Figure 2).

Figure 2: A simulation system based on an active database with an exterior supervisory control.

To reach the integration goal outlined above, the database has to fulfill the following demands:

- The database must support the integration of data (e.g. 3D simulation data) and algorithms (e.g. Petri Nets) in one single - now active simulation - database, supporting interface definition and providing means for state oriented as well as event based communication.
- The database must be able to flexibly adopt new data schemata (e.g. for the representation of various kinds of OPN) for its internal database, without additional alteration to the core programming.
- For the (real-time) simulation performance it is important how the data can be accessed and manipulated by the simulation algorithms (e.g. to implement even complex controllers using OPN). Ideally, database management itself should be time efficient, thus leaving computing power available to the simulation routines.
- In addition to this, it must be possible to easily add new simulation algorithms or enhance existing methods, while guaranteeing stability and performance of the overall system.
- The database itself must be independent from the type of simulation, to be able to incorporate quasi continuous as well as discrete event simulation paradigms into one single integrated simulation framework as depicted in Figure 3.

The rest of the paper is organized as follows: After detailing the architecture of the real-time simulation database in section 2 we discuss an important system extension, the plugin providing the basis for 3D simulation. In section 3 we then introduce our State Oriented Modeling Language, its syntax and how it is integrated into the graph database. To demonstrate the generality and versatility of our approach we show a selection of applications in section 4, based on the database and incorporating State Oriented Modeling.

2 THE REAL-TIME SIMULATION DATABASE

To fulfill the requirements mentioned in the introduction, and to eliminate unnecessary dependencies and provide a sustainable basis for various and diverse simulation applications, we developed a new architecture for 3D simulation systems, which is based on a small (micro-) kernel. This kernel is the Versatile Simulation Database (VSD), a real-time database containing all the data and algorithms needed for simulation applications. Fully implemented in C++, it provides the central building blocks for data management, meta-information, communication, persistence and user interface. The design of the simulation system as shown in Figure 4 is inspired by the Object Management Group (OMG) meta model hierarchy (Kurtev and van den Berg, 2005).

2.1 Meta Information

The uppermost layer (labeled M2 in Figure 4) is the meta-information system. It is essential for the flexibility, as well as the developer and end user friendliness of the database and the simulation system. The meta-information system is the basis for persistence, user interface, parallel and distributed simula-
tion, scripting and communication. It mainly consists of the following classes:

- **MetaTypeVal.** Describes all data types that can be used as values (e.g., int, double, string, simple structs, enumerations, flags).
- **MetaProperty.** Describes a property (see 2.3) with its getter and setter functions, its data type and a number of additional flags. These flags describe the behavior of the property as exposed to the user (editable, savable, etc.) as well as the properties ability to be used in parallel and distributed simulation.
- **MetaMethod.** Describes a method (member function) of an instance or a type.
- **MetaInstance.** Describes an instance including its class hierarchy. Each non-abstract meta-instance is able to create corresponding instances. Each meta-instance holds a list of the corresponding meta-methods and meta-properties, and furthermore provides a central point for executing member functions called "vsdMetaCall". This allows the corresponding instance to be used in generic scripting or parallel and distributed simulation. Messages are sent to instances by calling predefined methods via the vsdMetaCall function of the receiving instance.

In addition to "build-in" classes, it is also possible to generate meta-instances with the corresponding meta-properties and meta-methods during runtime (for example for object oriented scripting or new data models). Such "run-time meta-instances" are treated in exactly the same way as the build in meta-instances. There is no performance overhead in the data management.

The middle layer (labeled M1 in Figure 4) describes the data model of the simulation. In order to be able to retain semantic information and integrate data and algorithms into one single database, the VSD data model is an object oriented graph database (Gyssens et al., 1994), whose structure is detailed in this section. A simplified class hierarchy of the VSD core is shown in Figure 5.

2.2 Instances

The simulation model (labeled M0 in Figure 4) is an instantiation of the data model. All nodes in the graph database, the database itself and even the simulation environment are derived from a single base class called "Instance". This base class provides mechanisms for inter-instance communication, as well as access to the meta-information system, which allows introspection of class hierarchy, properties and methods (see 2.1).

The simulation model (labeled M0 in Figure 4) is an instantiation of the data model.

2.3 Properties

Derived from the instance class is the "SimStateInstance". Besides providing a reference to its simulation state (see 2.4) it may contain so called "Properties" and encapsulates the access to them. Properties are standardized getter and setter functions that encapsulate the data itself. All data in the simulation system is stored as properties. Properties can encapsulate any single value or value containers (lists, vectors, etc.), whose data types are known to the meta-information system. Properties can also hold refer-
ences or lists of references. References come in two different varieties, composite aggregation (with reference counting as described in (Levanoni and Petrank, 2006)), and shared aggregation. All parent child relations within the database are implemented as composite aggregation references. Shared aggregation references do not change the reference counter of the instance, but are informed if the instance gets deleted.

2.4 Database Structure

As shown in Figure 5 all nodes in the graph database, as well as the database itself are derived from the instance base class.

- **Environment and Simulation States.** The complete simulation is described by an “Environment”, which contains at least one “Simulation State” (SimState). Simulation states provide mechanisms for copying content from one state to another and are thus the basis for data partitioning in distributed or parallel simulation. For this, a simulation state can keep a list of all transactions, which can then be used to apply all bundled state changes to another simulation state on the same or on other computers. For example, when streaming data from an external database a separate thread with its own simulation state does handle the database interface. When the data has been loaded it will be transferred into the main simulation state.

- **Container and Element Index.** The database graph itself is kept by a “Container” class, a collection of graph nodes. A special container is the “Database” class, which acts as the spanning tree of the database. Other containers can be constructed, offering a different “View” onto the database, by rearranging all nodes or a subset thereof in different order. An example for this is the spatial view, which shows nodes in their spatial arrangements and gets updated when objects are grabbed or moved by other objects.

Furthermore the database offers convenience access to specific instance types via the “ElementIndex” class, which for every meta-instance provides a list of all instances of that type. By this mechanism it is possible to view the graph database in a traditional table based manner without performance restrictions.

- **NamedModelInstances.** This derivation from SimStateInstance provides a name property, as well as a list of extensions.

- **Nodes.** Most commonly used is the ”Node” class, which adds a child reference list property to the "NamedModelInstances" class.

- **Extensions.** “Extensions” are used to add data and functionality to a variety of nodes. Extensions can not only be attached to nodes, but also to other extensions.

2.5 Active Database

As mentioned above the VSD is not only a static data container but also contains the simulation algorithms itself. The environment, as well as all containers and element indexes actively inform interested listeners about new instance creation or deletion, as well as property modifications. Furthermore each instance sends a signal when one of its properties has been changed. Thus interested parts of the simulation systems can be informed about state changes in the simulation, eliminating the need to continuously poll the database content. With the active messaging system and the availability of element index lists we have minimized the need for data polling and tree traversal. By creating derived classes (like “Actuator”, “Sensor”, “Robot”, etc.) from the Node or Extension base class, the simulation algorithms (actuator control, sensor simulation, robot controller, etc.) itself are integrated into VSD. Simulation algorithms that need a complete overview over the simulation state (like rigid body simulations) are integrated on the database level but still manage their data on the node and extensions level as illustrated before.

That’s why we call the VSD active. To achieve a complete decoupling of the different system components with well defined interfaces (introspectable using the meta-information system), methods are provided for event as well as state based communication. These methods can be used to let the components exchange information as defined by the algorithm developer or the simulation expert.

2.6 3D Simulation

As already stated above, even the 3D simulation capabilities are an extension of the core database. An excerpt of class hierarchy for these classes is shown in Figure 6.

The data for 3D simulation may be interpreted by a collision detection system, a kinematic animation system, the physics simulation, the renderer or other application specific simulation algorithms. The most important classes are:

- **3D Node.** A node which contains a frame property describing its position and orientation relative to the parent node.
Figure 6: Add-ons to the core database (white): components for 3D simulation (blue).

- **Hull Node.** A 3D node that holds a reference to geometry. It also has a reference list of map nodes and an associated material node.

- **Geometries.** This instance holds all data necessary to describe a geometry. Properties include vertex, facet and texture coordinate lists. It can be referenced (shared) from many hulls.

- **Maps and Materials:** A map is a node with an url to a texture. Furthermore the function of the texture is given (for example: color map, normal map or diffuse map). A material node describes physical properties of a geometry, like color or electric conductivity. Maps and material are not only used for rendering. For example it is possible to attach a radar reflection intensity map to a hull node, in order to enable a more realistic radar simulation.

2.7 Simulation System

For real-world applications the database must be extended by new data schemes and simulation algorithms like 3D simulation (VSD3D, see section 2.6) or State Oriented Modeling (VSDNet, see section 3). This concept is illustrated in Figure 7. Further functionality like rendering, data processing, file loading, hardware interfaces or simulation scheduling, is handled by plugins, which may also add database enhancements like kinematics, dynamics (detailed in (Jung, 2011)), process simulation or GIS (Geo Information Systems, see (Longley et al., 2005)).

3 STATE ORIENTED MODELING LANGUAGE

In this section we introduce the features of the "State Oriented Modeling language" (SOML++) and show how object oriented Petri-nets are integrated into our simulation system. As mentioned before, the goal was to make State Oriented Modeling and Petri-nets an integral part of the simulation system.

3.1 Supervisory Control

This basic concept to integrate Petri-nets with 3D simulation systems is based on the Supervisory Control approach. The idea of Supervisory Control describes a technique to regard a control component (Supervisor) as a discrete-event simulation (DES), with its own state space and transitions caused by events that reflect changes at observed outputs of the controlled system (Plant). Figure 8 shows the basic structure of Supervisory Control: selected, re-fed state transitions \( \sigma \) of the plant generate events, which trigger further state transitions in the supervisor. Based on those state transitions, the supervisor can react, in order to adjust the plant with control commands \( \gamma \).

Figure 8: The basic structure of Supervisory Control, where one DES (Supervisor) commands another DES (Plant).

With respect to the focus of this paper, the integration of Petri-Nets and 3D Simulation, Petri Nets are used to implement the control component (Super-
visor) as a DES supervising and controlling the 3D simulation (Plant).

### 3.2 Language Features

SOML++ is a language that describes objects that can contain Petri-nets (see (Bastide, 1995)). In an object-oriented fashion these objects can be derived from other objects and can encapsulate data.

In the global name-space all the language constructs described below are allowed.

- **Object-classes.** An object-class must be instantiated as an object in order to be used. It is possible to subclass any object-class, regardless whether it has been defined in the SOML++ code or is a build-in class like VSD3D::Node.

- **Objects.** An object can be created from scratch or it can be derived from any object-class. Objects can be constructed with arguments, which will get passed to the constructor function.

- **Properties.** A value of all data types, that are known to the meta-information system of the simulation. Properties can also hold references to SOML++ objects.

- **Functions.** A block of code to be executed. Like in C++, a function has any number of arguments and a defined return type. Functions can be called from other functions or transitions anywhere within the SOML++ script.

- **Places.** A place as defined for a traditional Petri-net. Places and transitions are special objects, thus it is possible to define properties or functions within them.

- **Transitions.** Transitions may contain conditions and actions. If the conditions are met a marker may pass the transition and the actions are executed. Both, conditions and actions, are defined by user defined code which may in turn call other functions. During code execution, the "mark" context is known, representing the mark object currently testing or traversing the transition.

- **Arrows.** Link places to transitions and transitions to places.

- **Start-place.** Declares the object containing this statement to be a token, which moves through the Petri-net as defined by places, arrows and transitions (in the example: each DEMO instance becomes a token accessible by the mark statement in transitions).

This short summary outlines only a few language elements. In addition to this, the Petri-net implementation of State Oriented Modeling provides arrow conditions, different arrow types (normal, inhibitor, communication), Petri-net substitution and invocation hierarchy with arguments and return values, to name only the most important features.

### 3.3 Interpreting and Executing SOML++ Code

For the representation of the SOML++ code within the simulation database, a new data schema "VSDNet" as described in Figure 9 has been developed, providing classes for all SOML++ language elements.

![Figure 9: Add-ons to the core database (white): components for state oriented modeling (yellow and green).](image)

When a piece of SOML++ code is loaded by the simulation system, equivalent VSDNet "SourceComponent"-instances are created for classes, objects, arrows, etc. After that, the source representation in the simulation database is traversed and meta-instances (see Section 2.1) are generated for each object-class and each object. Properties are mirrored as meta-properties and functions as meta-methods. Afterwards new instances (see Section 2.2) are created from this meta-data and added to the database. Contained places and transitions are sub-nodes of the object nodes, arrows are modeled as sub-nodes of the originating place and transition nodes.

Since every SOML++ object has now become a
"normal" database node, net functions can use all the functionality provided by the database or the meta-information system, respectively. Net functions can interact with the rest of the database by creating new instances (using not only those defined in the SOML++ code, but all classes known to the simulation system), obtain references to other database nodes and call functions of these nodes.

Of course other database nodes can interact in the same way with the net objects, for example signals are sent when a net object is created or a property within this object has changed. There is no additional interface layer between the core database and the Petri-nets - the Petri-nets are an integral part of the simulation database.

Our simulation system also provides a comprehensive set of SOML++ debugging tools. All tokens can be traced and it is possible to set breakpoints within functions. Upon reaching a breakpoint a call stack and a list of local variables is available.

3.4 Realizing Supervisory Control

At this point, all components necessary for realizing Supervisory Control with State Oriented Modeling technique in 3D simulation systems are available. Using the same concepts, Petri-nets can observe any simulation state. State changes cause transitions to "fire" (representing events $\sigma$) which produce control commands $\gamma$.

4 APPLICATIONS

The hybrid simulation approach presented in this paper greatly simplifies the realization of new 3D simulation applications. Nearly all the applications realized so far benefit this approach. They use quasi continuous simulation technology to simulate kinematics, physics, actuators, sensors, various processes, etc. and State Oriented Modeling for supervisor and controller implementation, interfaces, user interaction, and supervisory control of the overall simulation. In this section we will focus on two application areas illustrating the application range of the concept presented above.

4.1 Virtual Testbeds

Using the concepts illustrated above we are now able to simulate complex systems with all relevant system components and their interdependencies. The result is a comprehensive development and testing environment based on simulation technology, a Virtual Testbed (see Figure 10).

The Virtual Testbed concept is a key technology
in the emerging field of eRobotics, because Virtual Testbeds can act as a central focal point in multidisciplinary development projects. For this reason, the first application areas of Virtual Testbeds are in the field of robotics, i.e. for the development of exploration robots, production plants or other complex systems.

### 4.2 Simulation based Control

The use of supervisory control and state oriented modeling is not confined to the virtual world. The very same concept can also be applied to control real hardware with the same software. We are using simulation technology to directly control physical systems, which we call "Simulation-Based Control" (Rossmann et al., 2012). This way, the same simulation and algorithms which were prototyped in a Virtual Testbed, control the actual hardware afterwards. An example is the control system of the multi-robot workcell depicted in Figure 11 consisting of two redundant 8-axis robots (linear axis plus a 7-axis robot).

**Figure 11: Simulation Based Control of a robot work cell.**

For robot control the database of our simulation system is extended with new types of node extensions, able to model and control kinematic chains and kinematic trees. The extension supports rotary and prismatic joints, as well as universal joints and joints directly defined via their Denavit-Hartenberg parameters.

The control concept of the multi robot system is based on the Intelligent Robot Control System (IRCS) structure, developed and introduced in the 1990s by (Freund and Rossmann, 1995). The IRCS addresses the main aspects of multi robot control by breaking up given tasks into smaller, manageable pieces in a "divide and conquer" fashion, delegating control over several layers of abstraction and responsibility. Figure 12 shows a simplified structure for the robot controller.

**Figure 12: The Intelligent Robot Control System.**

Here, the 3D simulation control (named "Multi Robot Control") acts as a coordinator of the vendor specific robot control units by implementing the simulation-based control concept. To communicate with the physical devices, the Ethernet based Fast-Research-Interface (FRI) is used for the KUKA Light-Weight-Robots (Bischoff et al., 2010); while a Profibus-Interface is used for the linear axes. The user interaction via a "User Interface" and the real-time robot coordination "Multi Robot Control" is performed on different computers, both running the same simulation system, though with different configurations on different operating systems (Windows and QNX). Both simulations use the same model, which is kept in sync between the computers by distribution methods provided by the core database.

The "Meta Control" layer (action generation and distributed planning using algorithms from the field of artificial intelligence), as well as the "Multi Robot Control" layer are Supervisory Controls implemented in SOML++.

### 5 CONCLUSIONS AND FUTURE WORK

In this paper we presented a new structure for an object oriented graph database for versatile 3D simulation systems. Due to their meta-information management, such systems can adapt to new data schemes even at run-time of the simulation, without the need for further programming. This approach allows us to integrate Petri-net objects as modeled in the State Oriented Modeling language. These Petri-net objects become an integral part of the simulation database and have full access to the 3D simulation data and algorithms, which enables supervisory control of quasi continuous simulation applications using dis-
crete event simulations. The result is a hybrid simulation system which has proven its applicability in large variety of applications, "classical" simulation applications like driving or production simulators, but also new fields of applications like GUI modeling or Virtual Testbeds.

Although performance of the interpreted SOML++ code segments never was a problem so far, we plan to introduce a compiler, which will transfer the generated meta-instances into native C++ code. In addition to this we plan the integration of further Petri-net approaches like hybrid and continuous petri nets (Alla and David, 1998) to widen the methodical base of the overall concept.

REFERENCES


