Using Models to Assess Impact of Defective Software

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Abstract. We explain a new strategy to model a set of java classes and to abstract a model from them that can be use to study the impact that defects affecting specific classes can have on the whole system. We use the models of the software implementation as an abstraction of the software that can be used for experimentation We used simulation and verification in SPIN but the idea can be applied to implementations in other languages than java and the analysis of defects impact can be done with other verification tools as well.

1 Introduction

As software applications become more complex, software reliability is becoming more important. To improve reliability we need to reduce probability of software failure, which can be catastrophic. An important approach to reducing probability of failure is by predicting and then removing software defects. Existing methods for defect prediction include those proposed in (Gaffney, 1984), (Khoshgoftaar and Seliya, 2004) and (Subramanyam and Krishnan, 2003). These methods are all machine learning based using some program features or program complexity measures such as the CK measure (Chidamber and Kemerer, 1994). The performance of these methods is still not desirable and much research is needed. For example, Menzies et al. studied defect prediction for 5 software systems of NASA (Menzies et al., 2004), with an average prediction accuracy of 36%. (Wilkie and Hylands, 1998) reported that approximately 48% of defects could be predicted by the combination of two of the CK measures.

This paper presents part of the findings in our BEACON project which focuses on the creation of strategies that facilitates the study of software defects and their consequences. We present here one of the strategies we investigated which as a qualitative analysis focus and relies on the abstraction of models out of a collection of java classes in such a way that the fundamental entities and interactions are represented. This model is then used to assess the effect of injecting defects in different parts of the system and through simulation and verification (Berard et al., 2001) study how they affect the system. The content of the paper is organized as follows. Section 2 explains the step of model creation then in section 3 we explain how we used it and we finalize in Section 4 with conclusions and proposed further work.

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2 Model Extraction

Part of the project was devoted to qualitative analysis of defects and how features from the code can be used to create models which inform the analysis of impact of those defects in the overall system. Starting with a java implementation of a system we looked at the classes and their interaction to extract a model. We choose to model systems using a well known and respected system: Promela/SPIN (Holzmann, 2003).

The analysis of a java implementation used for didactical purposes led the identification of classes/methods which suggested the network of possible interactions amongst the classes depicted in Figure 1. These classes are used to illustrate basic concepts in java to the students of one of our modules.



Fig. 1. Network representing the interactions of classes.

We use this network as a basis to define the processes and their interactions in the Promela model. A basic transformation algorithm can be easily implemented to translate a network into a Promela model:

- 1) for each node create a process with that name
- 2) for each transition create a channel with the names of both
- 3) the direction of the arrow in the graph dictates who send/receives the message
- 4) the interaction can be normal or defective

5) once a process is contacted by a defective module it sends a defective message to everyone it interacts with and then blocks

A more subtle analysis can be carried of in the java code to identify the modality of that interaction:

• for any interaction: is it mandatory or is only a possibility?

• when more than one call, are they independent or related (all have to be done)

• if more than one incoming transition: are both required as a precondition for the receiving node to operate or just some of them are enough?

In our example we considered all the arrows in figure 1 as "and's", that is that a node with two outgoing arrows call methods from both destination nodes and that a node with two incoming arrows is called by both classes in the origin nodes (notice this still omit details like in which order the methods are called) then we can refine the above models to obtain (a subset of the first model of this doc).

The following model is a first approach to representing possible behaviour in the

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set of java classes under examination.
    mtype = \{ok\};
     chan car_status = [0] of {mtype};
     chan driver_car = [0] of {mtype};
     chan driver_manufacturer = [0] of {mtype};
     chan manufacturer_car = [0] of {mtype};
     chan roadusers_driver = [0] of {mtype};
     chan roadusers_manufacturer= [0] of {mtype};
     active proctype roadUsers ()
      ł
                  if
                                 /* i.e., methods in the classes of
                                both destination nodes are invoked
                               but we do not know in which order
so we need to provide SPIN with the
option to choose any of the
possible orders */
                  :: roadusers_manufacturer!ok ;
                           roadusers_driver!ok
roadusers_driver!ok
                            roadusers_manufacturer!ok
                   fi
      active proctype driver ()
                                                                                                                                                     IGY PUBLIC
                                                                                                                 INOL
                                                                                                                                                                                                                                  ATIONS
{
                 if
                   11
                           roadusers_driver?ok
                                                                                              -->
                            if
                            :: driver_car!ok ;
    driver_manufacturer!ok
                            :: driver_manufacturer!ok;
                                      driver_car!ok
                            fi
             fi
}
    active proctype manufacturer ()
{ if /* it has to be contacted by both in
    any order to proceed contacting the
    next class */
    :: driver_manufacturer?ok ;
    roadusers_manufacturer?ok -->
    manufacturer_car!ok
    contaction is in the product of the product o
                      roadusers_manufacturer?ok ;
                ::
                         driver_manufacturer?ok -->
                        manufacturer_car!ok
               fi
     }
     active proctype car () { if /* it has to be contacted by both in
                      any order to proceed contacting the
next class */
driver_car?ok; manufacturer_car?ok-->
                       car_status!ok; manufacturer_car?ok; driver_car?ok-->
                ::
                        car_status!ok
               fi
     }
     active proctype status ()
              if
      £
                        car_status?ok
               fi
     }
```

3 Exploring the Model

Once we have a model we can take advantage of SPIN's multiple features to examine a model and inform us of the potential implications of a system built following the strategy represented by that model.

On first inspection the model created in the previous section does not seem to uncover any problems. Figure 2 shows a simulation.

We can now inject faults into the system. This was done manually but a recommendation of our project is that this can be incorporated to a tool that first distils a model from a group of classes and then allows the user to select in which areas of the model to inject faulty interactions which can then be analyzed using SPIN.

The following model has been injected faulty interactions in a couple of arbitrary places. The message types are now offering SPIN two types of interactions. A line "manufacturer_car!defect" has been added to proctype 'manufacturer' and another line "manufacturer_car?defect --> skip" has been added to proctype 'car'. We have also modified the model to allow the checking of some behavioural properties which we can use to force SPIN to conduct a more rigorous analysis. We have added to the model introducing a boolean variable 'end_reached' which we use to create a "guarantee" formula checking if the end of 'status' can be reached (see that 'status' is

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the final class in according to the network of Figure 1).
mtype = {ok, defect};
chan car_status = [0] of {mtype};
chan driver_car = [0] of {mtype};
chan driver_manufacturer = [0] of {mtype};
chan manufacturer_car = [0] of {mtype};
chan roadusers_driver = [0] of {mtype};
chan roadusers_manufacturer = [0] of {mtype};
bool end_reached=false;
active proctype roadUsers ()
     if
          /* i.e., methods in the classes of
         both destination nodes are invoked
but we do not know in which order
so we need to provide SPIN with the
         option to choose any of the
         possible orders *,
        roadusers_manufacturer!ok ;
        roadusers_driver!ok
     ::
        roadusers_driver!ok;
        roadusers_manufacturer!ok
    fi
}
active proctype driver ()
    if
       roadusers_driver?ok -->
if
    ::
        :: driver_car!ok
           driver_manufacturer!ok
           driver_manufacturer!ok;
           driver_car!ok
        fi
    fi
```

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If we run a verification of that property, SPIN detects a problem as indicated in Figure 3.

File	PIN CONTROL 4.3.0 22 June 2007 Edit View Run Help	SPIN DESIGN VERIFICATION	Line#: 50	Find:	Sequence Chart
* f ad { i : : : : : : : : : : : : : : : :	the proctype driver () if roaduses_driver?ok → driver_manufacturer(ok driver_cantok driver_cantok driver_cantok in we proctype manufacturer() // if has to be contacted by both in any order (proceed contacting by any order (proceed contacting by any order) for proceed contacting by the proctype manufacturer()				
n d n fi add c n c fi add (ii	We proctype car () // if has to be contacted by both in any order to proceed contacting th next class */ ifver_car?ok manufacturer_car?ol ar_statuslok nanufacturer_car?ok; driver_car?ol ar_statuslok the proctype status ()	€			Sundation Output Starting foreadUsers with pid 0 O proc -(root) creates proc 0 (roadUsers) Starting offered with pid 1 O proc -(root) creates proc 0 (roadUsers) Starting offered with pid 1 O proc -(root) creates proc 0 (roadUsers) O O proc -(root) creates proc 0 (treates 1) foodusers_manufacturer(o) O T proc -(root) creates proc 4 (fattas) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (fooduser_motherCoot) O proc -(root) creates proc 4 (fattas) (fattas 4 (food
3	starting simulation> pin -X -p -v -Y -g -s - pin -Z pan_in ;# preprod done preprocess>				3. proc 1 (driver) line 30 pan. In' (state 4) (driver, manufacturerio) 3. proc 2 (manufacturer) line 31 pan. In' (state 5) (driver, manufacturerio) 3. proc 2 (manufacturer) line 43 pan. In' (state 5) (driver, annufacturerio) 4. proc 2 (manufacturer) line 43 pan. In' (state 5) (driver, annol) 4. proc 3 (car) line 51 pan. In' (state 5) (driver, carlos) 5. sngle Stege Stegend 5. Sngle Stege Stegend

Fig. 2. A simple simulation in SPIN showing a possible way of interaction amongst classes.

If we follow the Guided Simulation offered by SPIN we get a counterexample as feedback which we show in Figure 4 and shows the system gets stalled when 'car' expects an OK from 'manufacturer' and gets a defective interaction instead.

File., Edit., View., Run., Help SPIN DESIGN VERIFICATION Lin	ine#: 37 Find:	Search for: Find
bool end_reached=false; adiwe prodype roadUsers () { f # r i.e., methods in the classes of both destination nodes are invoked but which orders owe need to provide SPIN with the option to choose an of		
:: roadusers_manufactureriok; roadusers_driveriok :: roadusers_driveriok; roadusers_menufactureriot; fi schve proctype driver ()	never claim - (none specified) assertion violations - (disabled by -A flag) cycle checks - (disabled by -DSAFETY) invalid end states +	
{	State-vector 56 byte, depth reached 18, errors: 1 20 states, stored 2 states, matched 22 states, matched 22 states, matched 0 atomic steps hash condicts: 0 (resolved)	
f active proceptor manufacturer() if /* than so be contacted by both in any order to proceed contacting the n dower, manufacturer() if or than a to be contacted by both in any order to proceed contacting the n dower, manufacturer() if order and the set of the se	er_carldefect State-vector as stored = 15128 byte + 8 byte overh er_carlok 2.097 memory used for hash table (-w19)	ompression: 23650.00%)
ache prochpe car () (if n' fi has be contacted by both in any order to proceed contacting the n charter car 20%; manufacture_car 20% -> car_stalualok manufacture_car 20% every c_coto +> car_stalualok manufacture_car 20% effect => stop)	line 17, fpan_in_state 4, "roadusers_manufactur (1 of 7 states) unreached in proctype driver (0 of 10 states) unreached in proctype manufacturer	
active proctype status () { if :: car_status?ok fi:	line 32, "pan_in", state 5, "roadusers_manufactur (1 of 12 states) unreached in proctype car (0 of 11 states)	rer?ok"
end_reached=true	unreached in prodype status (0 of 5 states) Suggested Action	
<pre> gcc -w -o pan -D_POSIX_SOURCE -DMEMLIM=128 -DSAFETY - pan.c pan.c pan.v -X -m10000 -w19 -A -c1</pre>	-DXUSANS -DNOFAR Optionally, repeat the run with a different search depth to fin Or, perform a GUIDED simulation to retrace the error found in of steps if the error was found at a depth greater than about	this run, and skip the first series

Fig. 3. Problems detected by SPIN after injection of faults in the model.



Fig. 4. Simulation with SPIN after faults in the model were detected.

4 Conclusions

The most important finding on this aspect of the project is that:

• there is a way to automatically translate java into Promela so that the impact of defects can be studied.

• a tool can be built which can perform the translation and also allow the injection of Promela code in the model representing faulty classes. These classes may be those which have been identified by statistical methods and this will enrich the assessment of the effect of defects within a system.

The strategy we explored in this paper is one of several approaches. The focus in this strategy was qualitative. Other strategies considered tried to quantify which classes may be more prone to faults. There is no single strategy or group of strategies that can provide definitive answers to these problems but a number of them may help to identify were a team should focus their efforts.

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