CHAOTIC ITERATIONS FOR STEGANOGRAPHY
Stego-security and Chaos-security

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Keywords: Steganography, Chaos, Security, Chaotic iterations.

Abstract: In this paper is proposed a novel steganographic scheme based on chaotic iterations. This research work takes place into the information hiding security fields. We show that the proposed scheme is stego-secure, which is the highest level of security in a well defined and studied category of attack called "watermark-only attack". Additionally, we prove that this scheme presents topological properties so that it is one of the firsts able to face, at least partially, an adversary when considering the others categories of attacks defined in the literature.

1 INTRODUCTION

Robustness and security are two major concerns in information hiding (Katzenbeisser and Dittmann, 2004; Domingo-Ferrer and Bras-Amorós, 2008). These two concerns have been defined in (Kalker, 2001) as follows. “Robust watermarking is a mechanism to create a communication channel that is multiplexed into original content [...] . It is required that, firstly, the perceptual degradation of the marked content [...] is minimal and, secondly, that the capacity of the watermark channel degrades as a smooth function of the degradation of the marked content. [...] . Watermarking security refers to the inability by unauthorized users to have access to the raw watermarking channel. [...] to remove, detect and estimate, write or modify the raw watermarking bits.” We will focus in this research work on security.

In the framework of watermarking and steganography, security has seen several important developments since the last decade (Barni et al., 2003; Cayre et al., 2005; Ker, 2006; Bras-Amorós and Domingo-Ferrer, 2008). The first fundamental work in security was made by Cachin in the context of steganography (Cachin, 1998). Cachin interprets the attempts of an attacker to distinguish between an innocent image and a stego-content as a hypothesis testing problem. In this document, the basic properties of a stegosystem are defined using the notions of entropy, mutual information, and relative entropy. Mittelholzer, inspired by the work of Cachin, proposed the first theoretical framework for analyzing the security of a watermarking scheme (Mittelholzer, 1999).

These efforts to bring a theoretical framework for security in steganography and watermarking have been followed up by Kalker, who tries to clarify the concepts (robustness vs. security), and the classifications of watermarking attacks (Kalker, 2001). This work has been deepened by Furon et al., who have translated Kerckhoffs’ principle (Alice and Bob shall only rely on some previously shared secret for privacy), from cryptography to data hiding (Furon, 2002). They used Diffie and Hellman methodology, and Shannon’s cryptographic framework (Shannon, 1949), to classify the watermarking attacks into categories, according to the type of information Eve has access to (Cayre et al., 2005; Perez-Freire et al., 2006), namely: Watermarked Only Attack (WOA), Known Message Attack (KMA), Known Original Attack (KOA), and Constant-Message Attack (CMA). Levels of security have been recently defined in these setups. The highest level of security in WOA is called stego-security (Cayre and Bas, 2008), whereas chaos-security tends to improve the ability to withstand attacks in KMA, KOA, and CMA setups (Guyeux et al., 2010).

To the best of our knowledge, there exist only two information hiding schemes that are both stego-secure and chaos-secure (Guyeux et al., 2010). The first one is based on a spread spectrum technique called Natural Watermarking. It is stego-secure when its parameter η is equal to 1 (Cayre and Bas, 2008). Unfortunately, this scheme is neither robust, nor able to face an attacker in KOA and KMA setups, due to its lack of a topological property called expansivity (Guyeux et al., 2010). The second scheme both
chaos-secure and stego-secure is based on chaotic iterations (Bahi and Guyeux, 2010b). However, it allows to embed securely only one bit per embedding parameters. The objective of this research work is to improve the scheme presented by authors of (Bahi and Guyeux, 2010b), in such a way that more than one bit can be embedded.

The remainder of this document is organized as follows. In Section 2, some basic recalls concerning both chaotic iterations and Devaney’s chaos are given. In Section 3 are presented results and information hiding scheme on which our work is based. Classes of attacks considered in this paper are detailed in Section 4. Stego-secure and chaos-secure are recalled too in this section. The new information hiding scheme is given in Section 5. Its stego-secure is studied in this section. The new information hiding scheme is given in Section 5. Its stego-secure is studied in this section. The topological framework making it possible to evaluate chaos-security is introduced in Section 7. Then the topological properties of chaotic iterations. As the next section, leading to the evaluation of its chaos-secure. This research work ends by a conclusion section where our contribution is summarized and intended future researches are presented.

2 BASIC RECALLS

2.1 Chaotic Iterations

In the sequel \(S^n\) denotes the \(n^{th}\) term of a sequence \(S\) and \(V_i\) is for the \(i^{th}\) component of a vector \(V\). Finally, the following notation is used: \([0;N] = \{0,1,\ldots,N\}\).

Let us consider a system of a finite number \(N\) of elements (or cells), so that each cell has a boolean state. A sequence of length \(N\) of boolean states of the cells corresponds to a particular state of the system. A sequence that elements belong into \([0;N-1]\) is called a strategy. The set of all strategies is denoted by \(\mathcal{S}\).

**Definition 1.** The set \(\mathcal{B}\) denoting \(\{0,1\}\), let \(f : \mathbb{B}^N \to \mathbb{B}^N\) be a function and \(S \in \mathcal{S}\) be a strategy. The so-called chaotic iterations are defined by \(x^n \in \mathbb{B}^N\) and \(\forall(n,i) \in \mathbb{N}^+ \times [0;N-1]\):

\[
x^n_i = \begin{cases} 
x^{n-1}_i & \text{if } S^n \neq i, \\
(f(x^{n-1}))_S & \text{if } S^n = i.
\end{cases}
\]

2.2 Devaney’s Chaotic Dynamical Systems

Some topological definitions and properties taken from the mathematical theory of chaos are recalled in this section.

Let \((x,d)\) be a metric space and \(f\) a continuous function on \((x,d)\).

**Definition 2.** \(f\) is said to be topologically transitive if, for any pair of open sets \(U,V \subset x\), there exists \(k > 0\) such that \(f^k(U) \cap V \neq \emptyset\).

**Definition 3.** \((x,f)\) is said to be regular if the set of periodic points is dense in \(x\).

**Definition 4.** \(f\) has sensitive dependence on initial conditions if there exists \(\delta > 0\) such that, for any \(x \in x\) and any neighborhood \(V\) of \(x\), there exist \(y \in V\) and \(n \geq 0\) such that \(|f^n(x) - f^n(y)| > \delta\).

\(\delta\) is called the constant of sensitivity of \(f\).

It is now possible to introduce the well-established mathematical definition of chaos (Devaney, 1989).

**Definition 5.** A function \(f : x \to x\) is said to be chaotic on \(x\) if:

1. \(f\) is regular,
2. \(f\) is topologically transitive,
3. \(f\) has sensitive dependence on initial conditions.

**Theorem 1** (Banks et al., 1992). If a function is regular and topologically transitive on a metric space, then the function is sensitive on initial conditions.

3 INFORMATION HIDING BASED ON CHAOTIC ITERATIONS

3.1 Topology of Chaotic Iterations

In this section, we give the outline proofs establishing the topological properties of chaotic iterations. As our scheme is inspired by the work of Guyeux et al. (Guyeux et al., 2010; Bahi and Guyeux, 2010b; Bahi and Guyeux, 2010a), the proofs detailed at the end of this document will follow a same canvas.

Let us firstly introduce some notations and terminologies.

**Definition 6.** Let \(k \in \mathbb{N}^+\). A strategy adapter is a sequence which elements belong into \([0,k-1]\). The set of all strategies with terms in \([0,k-1]\) is denoted by \(\mathcal{S}_k\).
Definition 7. The discrete boolean metric is the application \( \delta : B^2 \rightarrow B \) defined by \( \delta(x, y) = 0 \Leftrightarrow x = y \).

Definition 8. Let \( k \in \mathbb{N}^* \). The initial function is the map \( i_k \) defined by:
\[
i_k : S_k \rightarrow [0, k - 1] \quad (S^N)_{n \in \mathbb{N}} \mapsto S^0
\]

Definition 9. Let \( k \in \mathbb{N}^* \). The shift function is the map \( \sigma_k \) defined by:
\[
\sigma_k : S_k \rightarrow S_k \quad (S^N)_{n \in \mathbb{N}} \mapsto (S^{n+1})_{n \in \mathbb{N}}
\]

Definition 10. Given a function \( f : B^N \rightarrow B^N \), the function \( F_f \) is defined by:
\[
F_f : [0; N - 1] \times B^N \rightarrow B^N \quad (k, E) \mapsto (E, \delta(k, j) + f(E)_j, \delta(k, j))_{j \in [0; N - 1]}
\]

Definition 11. The phase space used for chaotic iterations is denoted by \( X \), and defined by \( X_1 = S_N \times B^N \).

Definition 12. Given a function \( f : B^N \rightarrow B^N \), the map \( G_f \) is defined by:
\[
G_f : X_1 \rightarrow X_1 \quad (S, E) \mapsto (\sigma_N(S), F_f(N(S), E))
\]

With these definitions, chaotic iterations can be described by the following iterations of the discrete dynamical system:
\[
\begin{align*}
X^0 & \in X_1 \\
\forall k \in \mathbb{N}^*, \ X^{k+1} = G_f(X^k)
\end{align*}
\]

Finally, a new distance \( d_1 \) between two points has been defined by:

Definition 13 (Distance \( d_1 \) on \( X_1 \)). \( \forall (S, E), (\bar{S}, \bar{E}) \in X_1, \ d_1((S, E); (\bar{S}, \bar{E})) = d_{BN}(E, \bar{E}) + d_{SN}(S, \bar{S}) \), where:
\[
\begin{align*}
& d_{BN}(E, \bar{E}) = \sum_{k=0}^{N-1} |S_k - \bar{S}_k| \in [0; N] \\
& d_{SN}(S, \bar{S}) = \frac{9}{N} \sum_{k=1}^{N} |S_k - \bar{S}_k|^2 \in [0; 1]
\end{align*}
\]

are respectively two distances on \( B^N \) and \( S_N \) (\( \forall N \in \mathbb{N}^* \)).

Remark 1. This new distance has been introduced by authors of (Bahi and Guyeux, 2010a) to satisfy the following requirements. When the number of different cells between two systems is increasing, then their distance should increase too. In addition, if two systems present the same cells and their respective starts with the same terms, then the distance between these two points must be small, because the evolution of the two systems will be the same for a while. The distance presented above follows these recommendations.

It is then proven that,

**Proposition 1.** \( G_f \) is a continuous function on \( (x_1, d_1) \), for all \( f : B^N \rightarrow B^N \).

Let us now recall the iteration function used by authors of (Bahi and Guyeux, 2010b).

**Definition 14.** The vectorial negation is the function defined by:
\[
f_0 : B^N \rightarrow B^N \quad (b_0, \ldots, b_{N-1}) \mapsto (\bar{b}_0, \ldots, \bar{b}_{N-1})
\]

In the metric space \( (x_1, d_1) \), \( G_{f_0} \) satisfies the three conditions for Devaney’s chaos: regularity, transitivity, and sensitivity. So,

**Theorem 2.** \( G_{f_0} \) is a chaotic map on \( (x_1, d_1) \) according to Devaney.

Finally, it has been stated in (Bahi and Guyeux, 2010a) that,

**Proposition 2.** The phase space \( X_1 \) has, at least, the cardinality of the continuum.

### 3.2 Chaotic Iterations for Data Hiding

To explain how to use chaotic iterations for information hiding, we must firstly define the significance of a given coefficient.

#### 3.2.1 Most and Least Significant Coefficients

We first notice that terms of the original content \( x \) that may be replaced by terms issued from the watermark \( y \) are less important than other: they could be changed without be perceived as such. More generally, a **signification function** attaches a weight to each term defining a digital media, depending on its position \( t \).

**Definition 15.** A signification function is a real sequence \( (u^k)_{k \in \mathbb{N}} \).

**Example 1.** Let us consider a set of grayscale images stored into portable graymap format (P3-PGM): each pixel ranges between 256 gray levels, i.e., is memorized with eight bits. In that context, we consider \( u^k = 8 - (k \mod 8) \) to be the \( k \)-th term of a signification function \( (u^k)_{k \in \mathbb{N}} \). Intuitively, in each group of eight bits (i.e., for each pixel) the first bit has an importance equal to 8, whereas the last bit has an importance equal to 1. This is compliant with the idea that changing the first bit affects more the image than changing the last one.

**Definition 16.** Let \( (u^k)_{k \in \mathbb{N}} \) be a signification function, \( m \) and \( M \) be two reals s.t. \( m < M \).

- The most significant coefficients (MSCs) of \( x \) is the finite vector
  \[
u_m = \left( k \mid k \in \mathbb{N} \text{ and } u^k \geq M \text{ and } k \leq \left| x \right| \right);
  \]
The least significant coefficients (LSCs) of $x$ is the finite vector

$$u_m = \{ k \mid k \in \mathbb{N} \text{ and } u^k \leq m \text{ and } k \leq |x| \}.$$ 

The passive coefficients of $x$ is the finite vector

$$u_p = \{ k \mid k \in \mathbb{N} \text{ and } u^k \in [m; M] \text{ and } k \leq |x| \}.$$ 

For a given host content $x$, MSCs are then ranks of $x$ that describe the relevant part of the image, whereas LSCs translate its less significant parts. These two definitions are illustrated on Figure 1, where the significance function ($u^k$) is defined as in Example 1, $M = 5$, and $m = 6$.

![Image](a) Original Lena. (b) MSCs of Lena. (c) LSCs of Lena ($\times 17$).

### 3.2.2 Presentation of the Scheme

Authors of (Bahi and Guyeux, 2010b) have proposed to use chaotic iterations as an information hiding scheme, as follows:

- $(K, N) \in [0; 1] \times \mathbb{N}$ be an embedding key,
- $X \in \mathbb{B}^N$ be the $N$ LSCs of a cover $C$,
- $(S^n)_{n \in \mathbb{N}} \in [0, N - 1]^\mathbb{N}$ be a strategy, which depends on the message to hide $M \in [0; 1]$ and $K$,
- $f_0 : \mathbb{B}^N \rightarrow \mathbb{B}^N$ be the vectorial logical negation.

So the watermarked media is $C$ whose LSCs are replaced by $Y_K = X^N$, where:

$$\begin{cases} X^0 = X \\ \forall n < N, X^{n+1} = G_{f_0}(X^n). \end{cases}$$

Two ways to generate $(S^n)_{n \in \mathbb{N}}$ are given by these authors, namely Chaotic Iterations with Independent Strategy (CIIS) and Chaotic Iterations with Dependent Strategy (CIDS). In CIIS, the strategy is independent from the cover media $C$, whereas in CIDS the strategy will be dependent on $C$. As we will use the CIDS strategy in this document, we recall it below. Finally, MSCs are not used here, as we do not consider the case of authenticated watermarking.

### 3.2.3 CIIS Strategy

Let us firstly give the definition of the Piecewise Linear Chaotic Map (PLCM, see (Shujun et al., 2001)):

$$F(x, p) = \begin{cases} x/p & \text{if } x \in [0; p], \\ (x - p)/(1 - p) & \text{if } x \in [p; \frac{1}{2}], \\ F(1 - x, p) & \text{else}, \end{cases}$$

where $p \in [0; \frac{1}{2}]$ is a “control parameter”.

Then, the general term of the strategy $(S^n)_n$ in CIIS setup is defined by the following expression:

$$S^n = [N \times K^n] + 1,$$

in which $\otimes$ denotes the bitwise exclusive or (XOR) between two floating part numbers (i.e., between their binary digits representation).

### 4 DATA HIDING SECURITY

#### 4.1 Classification of Attacks

In the steganography framework, attacks have been classified in (Cayre and Bas, 2008) as follows.

**Definition 17.** Watermark-Only Attack (WOA) occurs when an attacker has only access to several watermarked contents.

**Definition 18.** Known-Message Attack (KMA) occurs when an attacker has access to several pairs of watermarked contents and corresponding hidden messages.

**Definition 19.** Known-Original Attack (KOA) is when an attacker has access to several pairs of watermarked contents and their corresponding original versions.

**Definition 20.** Constant-Message Attack (CMA) occurs when the attacker observes several watermarked contents and only knows that the unknown hidden message is the same in all contents.
4.2 Stego-Security

In the prisoner problem of Simmons (Simmons, 1984; Bergmair and Katzenbeisser, 2006), Alice and Bob are in jail, and they want to, possibly, devise an escape plan by exchanging hidden messages in innocent-looking cover contents. These messages are to be conveyed to one another by a common warden, Eve, who over-drops all contents and can choose to interrupt the communication if they appear to be stego-contents.

The stego-security, defined in this framework, is the highest security level in WOA setup (Cayre and Bas, 2008). To recall it, we need the following notations:

- $K$ is the set of embedding keys,
- $p(X)$ is the probabilistic model of $N_0$ initial host contents,
- $p(Y|K_1)$ is the probabilistic model of $N_0$ watermarked contents.

Furthermore, it is supposed in this context that each host content has been watermarked with the same secret key $K_1$ and the same embedding function $e$.

It is now possible to define the notion of stego-security:

**Definition 21** (Stego-Security). The embedding function $e$ is stego-secure if and only if:

$$\forall K_1 \in K, p(Y|K_1) = p(X).$$

To the best of our knowledge, until now, only two schemes have been proven to be stego-secure. On the one hand, the authors of (Cayre and Bas, 2008) have established that the spread spectrum technique called Natural Watermarking is stego-secure when its distortion parameter $\eta$ is equal to 1. On the other hand, it has been proven in (Guyeux et al., 2010) that:

**Proposition 3.** Chaotic Iterations with Independent Strategy (CIBS) are stego-secure.

4.3 Chaos-Security

To check whether an information hiding scheme $S$ is chaos-secure or not, $S$ must be written as an iterative process $x^{n+1} = f(x^n)$ on a metric space $(X,d)$. This formulation is always possible (Bahi and Guyeux, 2010c). So,

**Definition 22** (Chaos-Security). An information hiding scheme $S$ is said to be chaos-secure on $(X,d)$ if its iterative process has a chaotic behavior according to Devaney.

In the approach presented by Guyeux et al., a data hiding scheme is secure if it is unpredictable. Its iterative process must satisfy the Devaney’s chaos property and its level of chaos-security increases with the number of chaotic properties satisfied by it.

This new concept of security for data hiding schemes has been proposed in (Bahi and Guyeux, 2010c) as a complementary approach to the existing framework. It contributes to the reinforcement of confidence into existing secure data hiding schemes. Additionally, the study of security in KMA, KOA, and CMA setups is realizable in this context. Finally, this framework can replace stego-security in situations that are not encompassed by it. In particular, this framework is more relevant to give evaluation of data hiding schemes claimed as chaotic.

5 THE IMPROVED ALGORITHM

In this section is introduced a new algorithm that generalize the scheme presented by authors of (Bahi and Guyeux, 2010b).

Let us firstly introduce the following notations:

- $x^0 \in B^N$ is the $N$ least significant coefficients of a given cover media $C$.
- $m^0 \in B^p$ is the watermark to embed into $x^0$.
- $S_1 \in S_N$ is a strategy called place strategy.
- $S_2 \in S_P$ is a strategy called choice strategy.
- Lastly, $S_3 \in S_P$ is a strategy called mixing strategy.

Our information hiding scheme called Steganography by Chaotic Iterations and Substitution with Mixing Message (SCISMM) is defined by $\forall(n, i, j) \in N^* \times \{0; N - 1\} \times \{0; P - 1\}$:

$$x^n_i = \begin{cases} x^{n-1}_i & \text{if } S^n_i \neq i \\ \frac{m^n_j}{m^{n-1}_j} & \text{if } S^n_i = i. \end{cases}$$

$$m^n_j = \begin{cases} m^{n-1}_j & \text{if } S^n_i \neq j \\ \frac{1}{m^{n-1}_j} & \text{if } S^n_i = j. \end{cases}$$

where $m^n_j$ is the boolean negation of $m^{n-1}_j$.

The stego-content is the boolean vector $y = x^P \in B^N$.  

6 STUDY OF STEGO-SECURITY

Let us prove that,
Proposition 4. SCISMM is stego-secure.

Proof. Let us suppose that \( x^n \sim U(B^N) \) and \( m^n \sim U(B^P) \) in a SCISMM setup. We will prove by a mathematical induction that \( \forall n \in \mathbb{N}, x^n \sim U(B^N) \). The base case is obvious according to the uniform repetition hypothesis.

Let us now suppose that the statement \( x^n \sim U(B^N) \) holds for some \( n \). For a given \( k \in B^N \), we denote by \( \tilde{k} = (k_0, k_1, \ldots, k_{N-1}, k_{N-2}, \ldots, k_0) \).

Let \( E_{i,j} \) be the following events:

\[
E_{i,j} = \{ i \land S_2^{n+1} = j \} \land m_j^{n+1} = k_i \land (x^n = k \land x^n = \tilde{k}) \land p = P(x^n+1 = k).
\]

These four events are independent in SCISMM setup, thus:

\[
p = \sum_{i \in [0,N-1], j \in [0,P-1]} P_1(i) P_2(j) P_3(i,j) P_4(i).
\]

According to Proposition 3, \( P(x_{j+1}^{n+1} = k_i) = \frac{1}{2} \). As the two events are incompatible:

\[
P(x^n = k \land x^n = \tilde{k}) = P(x^n = k) + P(x^n = \tilde{k}).
\]

Then, by using the inductive hypothesis:

\[
P(x^n = k) = \frac{1}{2^N}, \quad \text{and} \quad P(x^n = \tilde{k}) = \frac{1}{2^N}.
\]

Let \( S \) be defined by

\[
S = \sum_{i \in [0,N-1], j \in [0,P-1]} P_1(i) P_2(j).
\]

Then \( p = 2 \times \frac{1}{2^N} \times S = \frac{1}{2^N} \times S \).

\[S \]

\[= \sum_{i \in [0,N-1], j \in [0,P-1]} P_1(i) P_2(j) = \sum_{i \in [0,N-1]} P_1(i) \sum_{j \in [0,P-1]} P_2(j).
\]

The set of events \( \{ S_1^{n+1} = i \} \) for \( i \in [0,N-1] \) and the set of events \( \{ S_2^{n+1} = j \} \) for \( j \in [0,P-1] \) are both a partition of the universe of possible, so \( S = 1 \).

Finally, \( P(x_{j+1}^{n+1} = k) = \frac{1}{2^N} \), which leads to \( x^{n+1} \sim U(B^N) \). This result is true \( \forall n \in \mathbb{N} \), we thus have proven that the stego-content \( y \) is uniform in the set of possible stego-content, so \( y \sim U(B^N) \) when \( x \sim U(B^N) \).

\[\Box\]

7 TOPOLOGICAL MODEL

In this section, we prove that SCISMM can be modeled as a discrete dynamical system in a topological space. We will show in the next section that SCISMM is a case of topological chaos in the sense of Devaney.

7.1 Iteration Function and Phase Space

Let

\[
F : [0: N - 1] \times B^N \times [0: P - 1] \times B^P \longmapsto B^N
\]

\[
(k, x, \lambda, m) \longmapsto (\delta(k, j), x_j + \delta(k, j), m_j)
\]

where + and \cdot are the boolean addition and product operations.

Consider the phase space \( X_2 \) defined as follow:

\[
X_2 = S_N \times B^N \times S_P \times B^P \times S_P
\]

where \( S_N \) and \( S_P \) are the sets introduced in Section 5.

We define the map \( g_{f_0} : X_2 \rightarrow X_2 \) by:

\[
g_{f_0}(S_1, x, S_2, m, S_3) = (\sigma_N(S_1), F(x, S_2, m), \sigma_P(S_2), G_m(S_2, S_3), \sigma_S(S_3))
\]

Then SCISMM can be described by the iterations of the following discrete dynamical system:

\[
\begin{align*}
X_0 \in X_2 \\
X_{n+1} = g_{f_0}(X_n).
\end{align*}
\]

7.2 Cardinality of \( X_2 \)

By comparing \( X_2 \) and \( X_1 \), we have the following result.

Proposition 5. The phase space \( X_2 \) has, at least, the cardinality of the continuum.

Proof. Let \( \varphi \) be the map defined as follow:

\[
\varphi: X_1 \rightarrow X_2
\]

\[
(S, x) \mapsto (S, x, 0, 0, 0)
\]

\( \varphi \) is injective. So the cardinality of \( X_2 \) is greater than or equal to the cardinality of \( X_1 \). And consequently \( X_2 \) has at least the cardinality of the continuum. \( \Box \)

Remark 2. This result is independent on the number of cells of the system.

7.3 A New Distance on \( X_2 \)

We define a new distance on \( X_2 \) as follow: \( \forall X, \tilde{X} \in X_2 \), if \( X = (S_1, x, S_2, m, S_3) \) and \( \tilde{X} = (\tilde{S}_1, \tilde{x}, \tilde{S}_2, \tilde{m}, \tilde{S}_3) \), then:

\[
d_2(X, \tilde{X}) = d_{g_N}(x, \tilde{x}) + d_{g_P}(m, \tilde{m}) + d_{g_S}(S_1, \tilde{S}_1) + d_{g_S}(S_2, \tilde{S}_2) + d_{g_S}(S_3, \tilde{S}_3),
\]

where \( d_{g_N}, d_{g_P}, d_{g_S} \), and \( d_{g_S} \) are the same distances than in Definition 13.
7.4 Continuity of SCISMM

To prove that SCISMM is another example of topological chaos in the sense of Devaney, \( g_{f_0} \) must be continuous on the metric space \((x_2, d_2)\).

**Proposition 6.** \( g_{f_0} \) is a continuous function on \((x_2, d_2)\).

**Proof.** We use the sequential continuity.

Let \((S_1^n, x^n, (S_2^n), m^n, (S_3^n))_{n \in \mathbb{N}} \) be a sequence of the phase space \(x_2\), which converges to \((S_1, x, S_2, m, S_3)\). We will prove that \((g_{f_0}((S_1^n, x^n, (S_2^n), m^n, (S_3^n)))_{n \in \mathbb{N}}\) converges to \(g_{f_0}(S_1, x, S_2, m, S_3)\). Let us recall that for all \(n\), \((S_1^n, (S_2^n)\) and \((S_3^n)\) are strategies, thus we consider a sequence of strategies (i.e., a sequence of sequences).

As \(d_2((S_1^n, x^n, (S_2^n), m^n, (S_3^n)), (S_1, x, S_2, m, S_3))\) converges to 0, each distance \(d_{S_2}(x^n, x), d_{S_2}(m^n, m), d_{S_2}(S_3^n, S_3), d_{S_2}(S_3^n, S_3)\) converges to 0. But \(d_{S_2}(x^n, x)\) and \(d_{S_2}(m^n, m)\) are integers, so \(\exists \in \mathbb{N}, \forall n \geq n_0, d_{S_2}(x^n, x) = 0\) and \(\exists \in \mathbb{N}, \forall n \geq n_1, d_{S_2}(m^n, m) = 0\).

Let \(n_3 = \text{Max}(n_0, n_1)\). In other words, there exists a threshold \(n_3 \in \mathbb{N}\) after which no cell will change its state: \(\exists n_3 \in \mathbb{N}, n \geq n_3 \implies (x^n = x) \land (m^n = m)\).

In addition, \(d_{S_2}(S_2^n, S_2) \rightarrow 0\), \(d_{S_2}(S_3^n, S_3) \rightarrow 0\), and \(d_{S_2}(S_3^n, S_3) \rightarrow 0\), so \(\exists n_4, n_5, n_6 \in \mathbb{N}\):

- \(\forall n \geq n_4, d_{S_2}(S_1^n, S_1) < 10^{-1}\)
- \(\forall n \geq n_5, d_{S_2}(S_2^n, S_2) < 10^{-1}\)
- \(\forall n \geq n_6, d_{S_2}(S_3^n, S_3) < 10^{-1}\)

Let \(n_7 = \text{Max}(n_4, n_5, n_6)\). For \(n \geq n_7\), all the strategies \((S_1^n)\), \((S_2^n)\), and \((S_3^n)\) have the same first term, which are respectively \((S_1)_0\), \((S_2)_0\) and \((S_3)_0\): \(\forall n \geq n_7, (S_1^n)_0 = (S_1)_0 \land (S_2^n)_0 = (S_2)_0 \land (S_3^n)_0 = (S_3)_0\).

Let \(n_8 = \text{Max}(n_3, n_7)\). After the \(n_8\)-th term, states of \(x^n\) and \(x\) on the one hand, and \(m^n\) and \(m\) on the other hand, are identical. Additionally, strategies \((S_1^n)\) and \((S_1)_0\), \((S_2^n)\) and \((S_2)_0\), and \((S_3^n)\) and \((S_3)_0\) start with the same first term.

Consequently, states of \(g_{f_0}(S_1^n, x^n, (S_2^n), m^n, (S_3^n))\) and \(g_{f_0}(S_1, x, S_2, m, S_3)\) are equal, so, after the \(n_9\)-th term, the distance \(d_2\) between these two points is strictly smaller than \(3.10^{-2}\), so strictly smaller than 1.

We now prove that the distance between \((g_{f_0}((S_1^n, x^n, (S_2^n), m^n, (S_3^n))))\) and \((g_{f_0}(S_1, x, S_2, m, S_3))\) is convergent to 0. Let \(\varepsilon > 0\).

- If \(\varepsilon \geq 1\), we have seen that distance between \(d_{S_2}(S_1^n, x^n, (S_2^n), m^n, (S_3^n))\) and \(d_{S_2}(S_1, x, S_2, m, S_3)\) is strictly less than 1 after the \((n_9)\)-th term (same state).
- If \(\varepsilon < 1\), then \(\exists \in \mathbb{N}, n_0 \geq \frac{\varepsilon}{\varepsilon} \geq 10^{-(k+1)}\). As \(d_{S_2}(S_1^n, x^n, (S_2^n), m^n, (S_3^n))\) converges to 0, we have:
  - \(\exists n_9 \in \mathbb{N}, \forall n \geq n_9, d_{S_2}(S_1^n, x^n, (S_2^n), m^n, (S_3^n)) < 10^{-(k+2)}\)
  - \(\exists n_{10} \in \mathbb{N}, \forall n \geq n_{10}, d_{S_2}(S_1^n, x^n, (S_2^n), m^n, (S_3^n)) < 10^{-(k+2)}\)
  - \(\exists n_{11} \in \mathbb{N}, \forall n \geq n_{11}, d_{S_2}(S_1^n, x^n, (S_2^n), m^n, (S_3^n)) < 10^{-(k+2)}\).

Let \(n_{12} = \text{Max}(n_9, n_{10}, n_{11})\) thus after \(n_{12}\), the \(k + 1\) first terms of \((S_1^n)\) and \((S_1, S_2, m, S_3)\) are equal.

As a consequence, the \(k + 1\) first entries of the strategies of \(g_{f_0}(S_1^n, x^n, (S_2^n), m^n, (S_3^n))\) and \(g_{f_0}(S_1, x, S_2, m, S_3)\) are the same (due to the shift of strategies) and following the definition of \(d_{S_2}\) and \(d_{S_2}\):

\[d_2(g_{f_0}(S_1^n, x^n, (S_2^n), m^n, (S_3^n)); g_{f_0}(S_1, x, S_2, m, S_3))\]

is equal to:

\[d_{S_2}(S_1^n, S_1) + d_{S_2}(S_2^n, S_2) + d_{S_2}(S_3^n, S_3)\]

which is smaller than \(3.10^{-(k+1)} < 3, \frac{\varepsilon}{\varepsilon} = \varepsilon\).

Let \(n_0 = \text{max}(n_9, n_{12})\). We can claim that:

\[\forall \varepsilon > 0, \exists n_0 \in \mathbb{N}, \forall n \geq n_0, d_2(g_{f_0}(S_1^n, x^n, (S_2^n), m^n, (S_3^n)); g_{f_0}(S_1, x, S_2, m, S_3)) \leq \varepsilon\]


\(g_{f_0}\) is consequently continuous on \((x_2, d_2)\). \(\square\)

8 SCISMM IS CHAOTIC

To prove that we are in the framework of Devaney’s topological chaos, we have to check the regularity, transitivity, and sensitivity conditions.

8.1 Regularity

**Proposition 7.** Periodic points of \(g_{f_0}\) are dense in \(x_2\).

**Proof.** Let \((S_1, x, S_2, m, S_3) \in x_2\) and \(\varepsilon > 0\). We are looking for a periodic point \((S_1, x, S_2, m, S_3)\) satisfying \(d_2((S_1, x, S_2, m, S_3); (S_1, x, S_2, m, S_3)) < \varepsilon\).

As \(\varepsilon\) can be strictly less than 1, we must choose \(\tilde{x} = x\) and \(m = m\). Let us define \(k_0(\varepsilon) = \lfloor \log_{10} \frac{\varepsilon}{\varepsilon} \rfloor + 1\) and consider the set:

\[\{S \in S_1 \times S_2 \times S_3 | (S_1)^k = S_1^k \land ((S_2)^k = S_2^k) \land ((S_3)^k = S_3^k)\}, \forall k \leq k_0(\varepsilon)\].

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Then, \( \forall (S_1, S_2, S_3) \in S_{S_1, S_2, S_3, k_0}(\epsilon), \)
\( d_2((S_1, x, S_2, m, S_3); (S_1, x, S_2, m, \bar{S}_3)) < 3.\frac{\epsilon}{4} = \epsilon. \)
It remains to choose \((S_1, S_2, S_3) \in S_{S_1, S_2, S_3, k_0}(\epsilon)\) such that
\((S_1, \tilde{x}, S_2, m, \tilde{S}_3) = (S_1, \tilde{x}, S_2, m, \bar{S}_3)\) is a periodic point for \(g_{f_0}.\)

Let \( \gamma = \{ i \in [0; N - 1]/x_i \neq \bar{x} \}, \) where
\((S_1, x, S_2, m, S_3) = [g_{f_0}^{k_0}(S_1, \tilde{x}, S_2, m, \bar{S}_3)\}, \)
\( \lambda = \text{card}(\gamma), \) and \( f_0 < j_1 < \ldots < j_{\lambda - 1} \) the elements of \( \gamma. \)

1. Let us firstly build three strategies: \( S_1', S_2', \) and \( S_3' \), as follows.

(a) \((S_1')^k = S_1^k, (S_2')^k = S_2^k, \) and \((S_3')^k = S_3^k, \) if \( k < k_0(\epsilon). \)

(b) Let us now explain how to replace \( \tilde{x}_{j_t}, \forall q \in [0; \lambda - 1]; \) \( Ft. \)

First of all, we must replace \( \tilde{x}_{j_t}; \)

i. If \( \exists x_0 \in [0; N - 1]/x_{j_t} = m_{k_0}, \) then we can choose \((S_1')^{k_0 + 1} = f_0. (S_2')^{k_0 + 1} = k_0, (S_3')^{k_0 + 1} = k_0, \) and \( f_0 = 2. \)

ii. If such a \( k_0 \) does not exist, we choose:
\( (S_1')^{k_0 + 1} = f_0. (S_2')^{k_0 + 1} = 0, (S_3')^{k_0 + 1} = 0. \)
\( (S_1')^{k_0 + 2} = f_0. (S_2')^{k_0 + 2} = 0, (S_3')^{k_0 + 2} = 0, \)
and \( f_0 = 2. \)

All of the \( \tilde{x}_{j_t} \) are replaced similarly. The other terms of \( S_1', S_2', \) and \( S_3' \) are constructed identically, and the values of \( I_{Q_\gamma} \) are defined in the same way.

Let \( \gamma = \sum_{q=0}^{\lambda - 1} J_{q_\gamma}. \)

(c) Finally, let \((S_1')^{k_0 + \gamma} = (S_2')^{k_0 + \gamma} = (S_3')^{k_0 + \gamma}, \)
where \( j < k_0(\epsilon) + \gamma \) is satisfying \( j \equiv k [\text{mod } (k_0(\epsilon) + \gamma)], \) if \( k > k_0(\epsilon) + \gamma. \)

So, \( g_{f_0}^{k_0 + \gamma + \gamma}(S_1, \tilde{x}, S_2, m, S_3) = (S_1, \tilde{x}, S_2, m, \bar{S}_3) \)

Then, \((S_1, S_2, S_3) \in S_{S_1, S_2, S_3, k_0}(\epsilon)\) defined as previous is such that \((S_1, \tilde{x}, S_2, m, \bar{S}_3)\) is a periodic point, of period \( k_0(\epsilon) + \gamma + \alpha, \) which is \( \epsilon \)-close to \((S_1, \tilde{x}, S_2, m, S_3). \)

As a conclusion, \((x_2, g_{f_0}) \) is regular. \( \square \)

### 8.2 Transitivity

**Proposition 8.** \((x_2, g_{f_0}) \) is topologically transitive.

**Proof.** Let us define \( X : x_2 \rightarrow B^N, \) such that
\( x(S_1, x, S_2, m, S_3) = x \) and \( M : x_2 \rightarrow B^P, \) such that \( M(S_1, x, S_2, m, S_3) = m. \)

Let \( \mathcal{B}_A = B(X_A, ra) \) and \( \mathcal{B}_B = B(X_B, rb) \) be two open balls of \( x_2, \) with \( X_A = (S_1)^A_{x1}, (S_2)^A_{x2}, m, (S_3)^A_{x3} \) and \( X_B = (S_1)^B_{x1}, (S_2)^B_{x2}, m, (S_3)^B_{x3} \). We are looking for \( X = (S_1, \tilde{x}, S_2, m, \tilde{S}_3) \) in \( \mathcal{B}_A \) such that
\( \exists m_0 \in N, g_{f_0}^{m_0}([\tilde{x}]_A) \in \mathcal{B}_B. \)

\( \tilde{x} \) must be in \( \mathcal{B}_A \) and \( ra \) can be strictly lesser than 1, so \( \tilde{x} = x, \) and \( \tilde{m} = m. \)

Let \( k_0 = [\log_{10}(\beta^2) + 1]. \) Let us notice \( S_{x_2, x_3} = (S_1, S_2, S_3) \in \mathcal{S}_N \times (\beta^2)^2 / k \leq k_0, \)
\((S_1) = (S_2) = (S_3)^A_{x1}, (S_2) = (S_3) = (S_3) = (S_3)^B_{x3}. \)

Then \( \forall \tilde{x}(S_1, S_2, m, S_3) \in S_{x_2, x_3}, \)
\( \tilde{x} = (S_1, \tilde{x}, S_2, m, S_3) \in \mathcal{B}_A \).

Let \( \gamma = \{ i \in [0; N - 1]/x_i \neq x(X_B) \}, \) where
\((S_1, \tilde{x}, S_2, m, \tilde{S}_3) = g_{f_0}^{m_0}(X_B), \)
\( \lambda = \text{card}(\gamma), \)
and \( j_0 < j_1 < \ldots < j_{\lambda - 1} \) the elements of \( \gamma. \)

1. Let us firstly build three strategies: \( S_1', S_2', \) and \( S_3' \)

(a) \((S_1')^k = (S_1)^A_{x1}, (S_2')^k = (S_2)^B_{x2}, \) and \((S_3')^k = (S_3)^B_{x3}, \)
if \( k < k_0. \)

(b) Let us now explain how to replace \( X(X_B)_{j_0}, \)
\( \forall q \in [0; \lambda - 1]; \)

First of all, we must replace \( X(X_B)_{j_0}:
2. Let us secondly build three other strategies: $\tilde{S}_1$, $\tilde{S}_2$, and $\tilde{S}_3$ as follows.

(a) $\tilde{S}_1 = (S_1)^k$, $\tilde{S}_2 = (S_2)^k$, and $\tilde{S}_3 = (S_3)^k$ if $k \leq k_0 + \gamma$.

(b) Let us now explain how to replace $\mathcal{M}(X_B), \forall \mu \in [0; \mu - 1]$. First of all, we must replace $\mathcal{M}(X_B)$.

i. If $\exists \lambda_0 \in [0; P - 1]/X(\Xi_B)_{j_0} = \tilde{m}_{\lambda_0}$, then we can choose $(S_1)^{k_0+1} = j_0$, $(S_2)^{k_0+1} = \lambda_0$, $(S_3)^{k_0+1} = \lambda_0$, and so $I_{j_0}$ will be equal to 1.

ii. If such a $\lambda_0$ does not exist, we choose:

$(S_1)^{k_0+1} = j_0$, $(S_2)^{k_0+1} = 0$, $(S_3)^{k_0+1} = 0$, $(S_1)^{k_0+2} = j_0$, $(S_2)^{k_0+2} = 0$, $(S_3)^{k_0+2} = 0$ and so let us notice $I_{j_0} = 2$.

All of the $X(X_B)_{j_0}$ are replaced similarly. The other terms of $S_1$, $S_2$, and $S_3$ are constructed identically, and the values of $I_{j_0}$ are defined on the same way.

Let $\gamma = \sum_{m=0}^k f_{j_0}$.

(c) $(S_1)^k = (S_1^k)$, $(S_2^k) = (S_2^k)$ and $(S_3^k) = (S_3^k)$ where $j \leq k_0 + \gamma$ is satisfying $j \equiv k \mod (k_0 + \gamma)$, if $k > k_0 + \gamma$.

So, $g_{j_0}^{k_0+5}(S_1, S_2, m, S_3) = \tilde{S}_1 = (S_1, x_B, S_2, m, S_3)$.

Let $\mathcal{K} = \{i \in [0; P - 1]/m \neq \mathcal{M}(X_B)_i\}$, where $(S_1, x_B, S_2, m, S_3) = g_{j_0}^{k_0+5}(S_1, x_A, S_2, m, S_3)$.

Then we have proven the result.

\section{8.3 Sensitivity on Initial Conditions}

\begin{proposition}
\label{prop9}
If $f_{j_0}$ is regular and transitive. Due to Theorem 1, $\tilde{g}_{j_0}$ is sensitive.
\end{proposition}

\section{8.4 Devaney’s Chaos}

In conclusion, $(X_2, g_{j_0})$ is topologically transitive, regular, and has sensitive dependence on initial conditions. Then we have the result.

\begin{theorem}
\label{thm3}
$g_{j_0}$ is a chaotic map on $(X_2, d_2)$ in the sense of Devaney.
\end{theorem}

\section{9 CONCLUSIONS}

In this research work, a new information hiding scheme has been introduced. It is chaos-secure and stego-secure, and thus is able to withstand attacks in Watermark-Only Attack (WOA) and Constant-Message Attack (CMA) setups. These results have been obtained after having studied the topological behavior of this data hiding scheme. To the best of our knowledge, this algorithm is the third scheme that has been proven to be secure, according to the information hiding security field.

In future work, we intend to study the robustness of this scheme, and to compare it with the two other secure algorithms. Additionally, we will investigate the topological properties of our scheme, to see whether it is secure in KOA and KMA setups.

\section{REFERENCES}


