# A VECTORIZED TRAVERSAL ALGORITHM FOR RAY TRACING

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Keywords: Raytracing, SIMD.

Abstract: This article presents an optimized ray tracing algorithm which improves standard existing algorithms by processing simultaneously a large set of rays and carrying out a single traversal of the spatial indexing of the scene with all rays. This allows hardware SIMD functionality to be used efficiently and produces coherent memory accesses. Furthermore, during the single traversal, rays are grouped in such a way that these advantages are maintained for non coherent ray sets. The algorithm was observed to reduce the computation costs with respect to other standard solutions, especially for non coherent ray sets. It is worth noting that its characteristics make it especially suitable for graphics hardware.

# **1 INTRODUCTION**

Ray tracing is the dominating technique in the realistic image synthesis field, when the computation time is a secundary issue. The large processing times are due to the high computational costs of ray tracing. For each pixel in the image one or more rays must be sent through the scene and intersected with the objects it contains, the colour of the pixel must be calculated, and finally the pixels must be shown in the screen.

Ray tracing is not only useful in visualization. A multitude of algorithms, especially in the realistic image synthesis field, are based on ray tracing, such as *photon mapping* (Jensen, 1996). Furthermore, these algorithms usually work with non coherent rays. The availability of efficient ray tracing algorithms would allow a substantial decrease in the execution times of these algorithms. Due to all these reasons, the development of new techniques which bring together ray tracing and real time applications is interesting.

It should be taken into account that using this algorithm for ray tracing or more generally for Global Illumination, implies that the application must generate large ray sets in order to calculate their intersections simultaneously. This means that the algorithm must be organized in a breadth first scheme (Hanrahan, 1986); the speedup comes at the cost of increased memory use. Furthermore, this scheme allows an efficient vectorized implementation of other parts of the algorithm.

This article is organized as follows: section 2 provides a general description of the design ideas of the algorithm, followed by an in depth description of the algorithm. Lastly, section 3 offers an efficiency comparison to other methods.

## 1.1 Previous Work

Even though ray tracing is a very old technique (Appel, 1968), its use for interactive applications is relatively new.

(Wald, 2004; Wald et al., 2001) presents a CPU based algorithm which exploits the existing coherence in primary rays. It is based in the idea that coherent rays will very probably make the same traversal through the spatial indexing structure which contains the scene, and will intersect the same objects. His proposal thus consists in traversing the spatial index and making the ray-triangle intersections with packets made of four coherent rays at the same time. Using SIMD instructions of current processors, the computation time of the algorithm is reduced. In addition, Wald presents a careful design oriented towards minimizing code complexity and cache faults.

Noguera J., Ureña C. and Jesús García R. (2009). A VECTORIZED TRAVERSAL ALGORITHM FOR RAY TRACING. In Proceedings of the Fourth International Conference on Computer Graphics Theory and Applications, pages 58-63 DOI: 10.2220/001767500580063 Copyright © SciTePress This idea is valid for primary rays since it is easy to locate rays which a priori have high probabilities of coherency. But it fails for secundary rays where this assumption does not hold. This article presents an extension to this idea, which permits the parallelization of calculations for different rays in any situation, even when these are not coherent. A preliminary version of this work (in Spanish) can be seen in (Noguera and Ureña, 2007).

Recently (Wald et al., 2007) have begun the design of stream and filter based raytracers. Rays are processed into a series of streams, and the operations on the rays are modelled as filters. After the filters, rays to be processed further are compressed into another stream. The filters can process the rays in parallel using SIMD instructions.

A modelling software has been used to test the SIMD efficiency of this design for different SIMD widths (4 to 16), and for different sizes of the ray packets (2x2 to 64x64), for 1k rays. The efficiency for large SIMD width decreases noticeably, but still obtains higher efficency than previous approaches due to the compacting. However, no ray tracer has been implemented to test the results in real hardware. Our approach, on the other hand, should not be negatively affected by large SIMD widths. (Gribble and Ramani, 2008) simulates an approach similar to ours including larger SIMD width, and report good scalability.

Another approach (Boulos et al., 2008) reorders and compacts rays only when SIMD efficiency is less than 50 %, missing some of the ray coherence. Out of core extensions (Pharr et al., 1997) and low memory data structures (Mahovsky and Wyvill, 2006) have also been published.

There have been different attempts to use GPU processors for ray tracing acceleration (Carr et al., 2002; Purcell et al., 2002; Foley and Sugerman, 2005; Christen, 2005). Our approach is suitable for GPU porting as well.

# **2** IMPLEMENTATION

The following sections describe our ray tracer implementation. Firstly a general description of the design ideas used is shown, and their application to the algorithm. The following sections describe the implementation details.

## 2.1 Motivation

Given a spatial indexing tree which contains the scene, the classic ray tracing algorithm travels recursively across the tree once for each ray. Wald originally proposed travelling across the tree once for each four ray packet, and later furthered the idea by traversing the tree with a stream of up to 64 rays. Our proposal continues this line of work, and proposes travelling across the tree once for all rays at the same time.

This approach provides the following advantages: As long as two or more rays which must traverse the same tree node exist, they will do so simultaneously, allowing the parallelization of the calculations by means of the hardware SIMD instructions. This translates to much higher probabilities of having some rays to process in parallel in each tree node, especially in the higher tree nodes.

Another advantage is that since the tree nodes are visited only once, and not tens of thousands of times as is the case of previous algorithms, the number of accesses made to the spatial indexing structure of the scene is greatly reduced.

## 2.2 Parallelism due to SIMD

Current processors include SIMD extensions which allow computing many floating point operations (usually four) with one instruction, therefore increasing the performance of these calculations. Our implementation uses SSE, the Intel SIMD instruction set.

A new data structure which contains the rays to be processed is needed. This structure does not mean more memory accesses with respect to classic ray tracing algorithms, since the same information which also must be stored and accessed in classic algorithms is simply reorganized in groups.

#### 2.3 Traversing the Kd-tree

Our implementation uses a kd-tree in which the plane divides the voxel through the center for the spatial indexing of the scene. The algorithm can be easily changed to use division planes in other positions. The main reason of this choice instead of other structures is the simplicity of the traversal algorithm. For each node and ray only two decisions must be made: whether to follow or not each of the two child nodes. Additionally, this structure is well adapted to scene complexity, and is easy to parallelize.

In order to parallelize the calculations, at each node the list of all the rays traversing it must be available. Firstly the workings of the algorithm will be shown for one ray, and then the algorithm will be expanded for the case of all rays at the same time.

#### 2.3.1 Traversal Algorithm for One Ray

An adaptation to kd-trees (with dividing planes in the center of the voxels) of the incremental parametric

traversing algorithm for octrees presented in (Revelles et al., 2000) was used. Let r = p + td be the equation of a ray, p being the origin vector, d the unit director vector and  $t \ge 0$  the parameter which determines a point along the ray.

For each node,  $t_{x0}, t_{x1}$  are the values of the ray parameter which intersect with the two planes perpendicular to the X axis delimiting the node voxel (ditto  $t_{y0}, t_{y1}$  for Y, and  $t_{z0}, t_{z1}$  for Z). We assume  $t_{x0} < t_{x1}$ . These values are calculated for the root node and incrementally recalculated in each step of the tree traversal.

Let  $t_{xm}$  the parameter for which the ray crosses the plane (perpendicular to X) which divides the voxel in two equal subvoxels. It may be calculated as:  $t_{xm} = (t_{x0} + t_{x1})/2$ . For a voxel of the kd-tree divided by a plane perpendicular to X, the values of the intersections for the first sub-node visited are  $t_{x0}, t_{xm}$ , and for the second  $t_{xm}, t_{x1}$ . The values  $t_{ym}, t_{zm}$  are defined, calculated and used analogously for the voxels divided by planes perpendicular to Y or Z. In general, the calculation of the values of the child nodes of a node only require adding and halving. (Revelles et al., 2000) shows that the ray intersects a node if the following holds:

$$max(t_{x0}, t_{y0}, t_{z0}) < min(t_{x1}, t_{y1}, t_{z1})$$
 (1)

The algorithms works recursively. For each inner node, the ray parameters are updated for its first child, and intersection is determined. If an intersection exists, the node is processed. Later, the process is repeated for the second child.

If the tree is traversed selecting first the first node in the ray direction, once an intersection in the visited node is found the recursivity can be ended, since any other potential intersection will always be farther.

#### 2.3.2 Traversal Algorithm for Multiple Rays

The kd-tree traversal algorithm for multiple rays works similarly to the previous one. The idea is processing at the same time all available rays, making only one traversal of the spatial indexing for all rays.

The pseudocode can be seen in Figure 1. Each node gets a list of intersecting rays, with its parameters adjusted for that node. If the list is empty, we return. If the node is a leaf, the function TrgIntersec calculates if the rays intersect the triangles in the node. Otherwise, we are in an inner node. The function calcNode uses SSE to update the parameters of all the rays for the first child and to determine which rays intersect it (These are copied to a new list). Then, recursively, the first child is traversed with the new list

```
Traverse ( node, RayList )
if nRays = 0
  return;
if isLeaf(node)
  TrgIntersec( node, RayList );
else foreach child h<sub>i</sub> of node
  RayListC<sub>i</sub>=calcNode(h<sub>i</sub>,RayList);
  Traverse( h<sub>i</sub>, RayListC<sub>i</sub> );
```

Figure 1: Traversal for multiple rays algorithm.

of rays just calculated. Once the traversal of this child is finished, the second child is traversed analogously.

These lists do not include those rays for which an intersection was already found, nor those rays farther than the maximum allowed ray parameter, avoiding unnecesary calculations. The total number of operations required to determine if four rays intersect a node are: one addition, one multiplication, one maximum, one minimum and one comparison, all SSE.

When working with multiple rays it may happen that different rays require the spatial indexing tree in a different order. (The order depends on the signs of the three components of the director vector of the ray). Therefore, this happens if rays with director vectors of different signs exist. In these cases, it is not possible to traverse the tree only once for all rays simultaneously without losing the desirable characteristic that the first intersection found, is the nearest one.

The solution implemented to solve this problem consists in preclassifying all rays in eight groups by the signs of the components of their director vectors. Since all the rays in a group will traverse the indexing tree exactly in the same order, the traversal algorithm can be applied with no problems for each group. With this idea, the tree must be traversed a maximum of eight times, much lower than conventional algorithms. In practice, for primary rays the traversal is done at most four times, as long as the normal vector to the vision plane is parallel to one of the edges.

#### 2.3.3 Ray Memory Layout

The data memory layout chosen to store the rays takes into account the fact that the implementation uses SIMD instructions in order to maximize efficiency and time (Intel Corporation, 2008).

Two data structures are used to manage the rays. A global one stores all rays and attributes; a per node structure stores some data needed for the traversal, plus pointers to the global structure.

It is necessary to store globally for each ray its director and origin vector, the maximum allowed ray parameter and whether an intersection was already



Figure 2: SIMD data structure.

found (and a triangle identifier, ray parameter and parametric coordinates at the intersection if so).

As was stated in the algorithm description, during tree traversal each node receives a list of intersecting rays. This list is local for each node, and stores for each ray the entry and exit values of the parameters of the ray for the voxel ( $t_{0x}, t_{0y}, t_{0z}, t_{1x}, t_{1y}, t_{1z}$ ), since they depend on the current node. A pointer to the global structure which stores the rest of the ray information is also stored per ray.

This data will be processed by SIMD operations; therefore for maximum performance they should be stored with a well known SIMD data structure created by taking four objects and substituting each field of the original structure with a vector with the corresponding components of the objects. An array of these structures substitutes the original array of objects (Figure 2). In this figure,  $t_{xi}^{(a.b)}$  represents the values of the *x* component of the *t* parameter for the four rays in positions *a* to *b*. *i* is 0 for entry and 1 for exit parameters. The same applies to *y* and *z*. If the number of rays is not a multiple of four, the last packet will be incomplete and will be filled with bogus data which will not be taken into account in the final results.

Each node gets a ray list, and must generate a new list for each child. A ray stack is used to prevent dynamic memory overhead. On tree construction, memory is allocated for a large enough stack. For each traversal, that stack is initialized with all the rays to be intersected and sent to the root node. The ray lists for the following nodes visited will be pushed and poped into the stack as needed, avoiding the need to allocate memory in each step.

Each node receives on the top of the stack the rays which intersect it and which should be processed. It calculates which of these intersect its first child, and pushes them into the stack. The child gets these new rays. Upon finishing processing the first child node, the rays sent to the child are removed from the stack, and the same process is done for the second child. This procedure ensures that only the rays which intersect one of the children need to be stored in the stack.

#### 2.3.4 Ray-Triangle Computation

Two variants have been used for ray-triangle intersections:

• The standard Moller-Trumbore algorithm (Möller

and Trumbore, 1997), which tests intersection of one ray with one triangle.

• A variation on the previous one, optimized, with normals and various intermediate results precalculated, with the data of four triangles packed, and making at the same time the intersection of one ray with four triangles with SSE.

The results are better (both in time and memory use) with the second option, (which requires lower depth trees), and therefore this option was used in tests.

# **3 PERFORMANCE**

To study the efficiency of this algorithm, three different scenes were used (see figure 3): *Stanford Bunny* (69451 triangles), *Stanford Dragon* (0,8 million triangles) and *Happy Buddha* (1,1 million triangles). The scenes used were the same for all algorithms.

To compare the algorithm proposed, Wald's algorithms (Section 1.1) were simulated using our own algorithm but limiting the maximum number of rays which can be traversed together to 4 and 64x64. The algorithm proposed is also compared to a classic recursive algorithm which traverses the rays through the structure one at a time.

The results were obtained for a maximum allowed depth of the spatial indexing tree of 16, 18, 20, 22 and 24, and the best results were selected. For most of the cases, depth 20 was optimal. All measurements were taken on a 2.6 GHz Intel Core2 Duo with 2 GB of RAM. Only a single core was used.

## 3.1 Coherent Ray Comparison

For each case, a 100 image animation of the camera rotating around three different scenes was generated, and the mean results were taken. Since the rays have a common origin and cross consecutive pixels in the image plane, they present high coherence.

Graphs 4 (left), (center) and (right) show the results obtained for the different scenes and algorithms, for both coherent and non coherent rays. It can be seen in these graphs that the classic algorithm is less efficient in the coherent ray test for all the scenes. The other three algorithms are noticeably better, since they parallelize calculations. Among these three algorithms, the version which traverses the structure with 4-ray packets obtains slightly worse results. The two remaining algorithms obtain the best time results, and are very similar to each other.

The four-ray packet version almost always parallelizes the computation of the rays, since they are co-



Figure 3: The three scenes used for our experiments.



Figure 4: Average performance in MegaRays per second when rendering the Stanford Bunny (left), Standord Dragon (center) and Happy Buddha (right) at a resolution of 1280x1024 pixels.

herent. However, this is not guaranteed to always happen, and sometimes the algorithms needs to break the packet. working with a larger number of rays simultaneously mean a larger quantity of data to process, so that the possibilities of having to work with packages of size smaller than four are smaller.

## **3.2** Comparison for Non Coherent Rays

In this article it was said that the main advantage obtained from travelling the tree with all the rays simultaneously appears when non-coherent rays are used. These are classified automatically when descending the indexing tree and the possibilities of having some rays in the same node (which means being able to parallelize the calculations) are much higher, especially in nodes near the root.

To prove this, a uniform distribution has been generated which throws random rays in such a way that the density of the rays is the same for all regions, in the manner described by (Santalo, 1976; Sbert, 1993). This distribution of rays is the worst possible scenario, since ray coherence is lowest.

Graphs 4 left, center and right contrast the results obtained in this scenario to the results for coherent rays. It can be clearly seen that the results obtained by the algorithm which traverses the structure with four ray packets have been notably reduced with respect to those obtained for coherent rays. The results are very similar to the clasical recursive algorithm,



Figure 5: Average performance in MegaRays/s for different ray packet sizes for the dragon scene at 1280x1024.

because when taking non-coherent four-ray packets, the possibilities that the packet is broken in the higher levels of the tree are very high, and the algorithm degenerates to the conventional non parallel algorithm.

The algorithm which uses streams of 64x64 rays obtains better results since it can reorganize the rays, getting rid of invalid ones and increasing coherence.

In spite of this, our algorithm is able to achieve a much higher gain, near 4. Since it works with the full set of rays at the same time (1280x1024 rays), the reordering of the rays when descending the tree is the best possible one. This fact implies an optimal hardware utilization. Additionally, reducing the number of traversals to the minimum means that the overhead of function calls and access to the nodes of the tree are drastically reduced, which also increases the gain.

Graph 5 shows the mean performance for different sizes of the ray packet. For coherent rays, performance increases rapidly even with a small ray packet. This is due to the higher probability of finding coherent rays. In contrast, non coherent rays show a slight ascending trend as the packet size increases. A larger packet size implies higher probabilities of being able to parallelize the calculations; the maximum possible efficiency is obtained with our algorithm since the packet contains all the available rays.

# 4 CONCLUSIONS

Until now, acceleration techniques for ray-tracing were useful only when it was easy to find coherent rays to process in parallel. This article introduces a ray tracing algorithm which instead of traversing the spatial indexing of the scene once per ray or per ray packet, it traverses the tree only once with all rays simultaneously (or a maximum of eight times if the rays have director vectors of different sign).

This technique has been shown to present the following advantages: It can extract a higher level of parallelism in ray processing, even for non coherent rays, and a node will be visited at most eight times, resulting in minimization of memory accesses.

Future work includes disabling the ray stack scheme when the number of rays in a node is small, to increase efficiency, studying wider SIMD arquitectures, since speedup should be higher (Benthin et al., 2006), out of core extensions and implementation on GPUs. A battery of tests with large complete scenes will validate applicability in real-world scenarios.

# ACKNOWLEDGEMENTS

This article has been partially financed by projects TIN2004-07672-C03-02 of the Spanish Ministry of Education and Science, and P06-TIC-01403 of the Junta de Andalucía and the European Union (FEDER).

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