

A TWO-STAGED APPROACH FOR ASSESSING FOR THE QUALITY OF INTERNET SURVEY DATA

Chun-Hung Cheng

Dept of Systems Engineering & Engineering Management, The Chinese University of Hong Kong, Hong Kong

Chon-Huat Goh

School of Business, Rutgers University, Camden, NJ 08102, USA

Anita Lee-Post

School of Management, University of Kentucky, Lexington, KY 40506, USA

Keywords: Internet Survey Data, Quality, TSP, genetic algorithm.

Abstract: In this work, we propose to develop a procedure to detect errors in data collected through Internet surveys. Although several approaches have been developed, they suffer many limitations. For instance, many approaches require prior knowledge of data and hence they need different procedures for different applications. Others have to test a large number of parameter values and hence they are not very efficient. To develop a procedure to overcome the limitations of existing approaches, we try to understand the nature of this quality problem and establish its linkage to travelling salesman problem (TSP). Based on the TSP problem structure, we propose to develop a two-staged approach based on a genetic algorithm to help ensure the quality of Internet survey data.

1 INTRODUCTION

Surveys have been widely used in various disciplines to understand public opinions and views. Typically, interview over telephone and postal mail surveys are often used. Interview over telephone is only effective for small and simple surveys while postal mail surveys are costly and slow. With the ubiquitous of personal computers, and the availability of high-speed broadband network, Internet survey has become a reality.

With the Internet, surveys can be conducted to reach out a large number of potential survey subjects at very low cost. However, the subjects enter their responses without any assistance. Although data-type and data-range checking may be implemented together with an electronic survey form, they are not always effective given a diversity of survey questions and responses. Hence, data collected through this means may contain erroneous values. These erroneous data must be identified to ensure

the survey quality. In this work, we shall focus on the detection of unsystematic errors. These errors are those that are caused by survey design faults.

There are two kinds of approaches to assessing the quality of data. Application-dependent approaches use knowledge of data and, therefore, require different quality-checking procedures for different survey applications. On the other hand, application-independent approaches do not need any prior knowledge of data and provide one general quality-checking procedure for all applications. The flexibility and generalisability of application-independent approaches make them more appealing than the application-dependent approaches. In this work, we shall also focus on application-independent approaches.

Currently, application-independent approaches use clustering analysis techniques. Clustering analysis gathers data records into groups or clusters based on their field values. Similar data records occupy the same group while dissimilar records do not coexist in the same group. Records, whose field values make

them significantly different from all others, may not find themselves related to any other group members at all (see Figure 1). They are called outliers (Storer and Eastman, 1990).

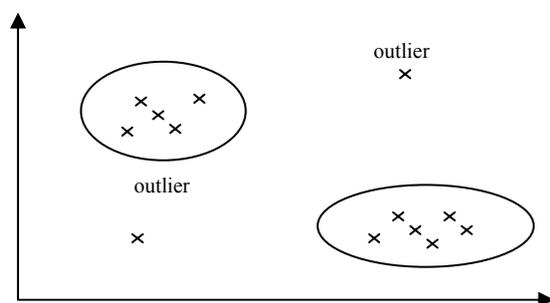


Figure 1: Examples of outliers.

In this work, we apply a quality-checking technique for data collected through Internet surveys. The new approach provides many advantages over current approaches.

- Unlike some current approaches, our approach does not require knowledge of data. Therefore, this approach is applicable to many different Internet survey applications.
- Many clustering approaches, including ours, require the user to test the clustering algorithms with different parameter values. However, we will show that our approach significantly cuts down on the search and evaluation of a number of parameter values and, therefore, is more efficient than other approaches.
- Many current approaches do not specify systematic ways to choose the best clustering result among many candidate solutions. Our approach applies a selection criterion to systematically evaluate the clustering of a sample data in order to best separate the outliers from similar records in a group.

2 LITERATURE REVIEW

Although the classical database literature considers errors in a database a serious problem (e.g., Felligi and Holt, 1976 and Naus et al., 1972), few studies propose ways to deal with the problem. There are two kinds of quality-checking approaches: application-dependent and application-independent approaches.

2.1 Application-Dependent Methods

Application-dependent approaches such as those by Freund and Hartley (1967), Naus et al. (1972), and Felligi and Holt (1976) are all statistical-based. In detecting errors in a database, these approaches require knowledge of the data. Using these approaches, software developers may have to develop different programs for different database applications.

2.2 Application-Independent Methods

All application-independent approaches use clustering analysis techniques. Lee et al. (1978) first applied a clustering approach. They defined a distance function to measure the difference between two records. Based on a distance matrix, they found the shortest path between a pair of records. Since the determination of the shortest path is an NP-complete problem (Storer and Eastman, 1990), the shortest spanning path algorithm (Slagle et al., 1975) is used to find an approximate solution. A link between two records that is longer than the pre-specified threshold value will be broken. Records whose distances are less than the threshold value are similar and are placed in the same group. A record with no similar partners is an outlier.

Storer and Eastman (1990) proposed three related clustering approaches. They used the same distance function as defined by Lee et al. (1978). The first approach is called the leader algorithm (Hartigan, 1975). The leader algorithm clusters M records into K groups, where M and K are positive integer values and $M \geq K$. It assumes that the distance function between two records and the threshold value for group membership are available. The first record is a leader for the first group. A record is assigned to an existing group if its distance from the group leader is less than the threshold value. It becomes a new leader for a new group if its distance from every existing leader is more than the threshold value.

The second approach is a modification of the leader algorithm that we refer to as an average record leader algorithm. This modified algorithm uses the average record instead of the first record as an initial leader. Therefore, the algorithm can generate a solution independent of record order. On each pass, a record that is furthest from its group leader becomes a leader for a new group. If the algorithm were to produce K groups, it requires K passes through the data.

The third approach is another modification of the leader algorithm. Storer and Eastman (1990) call it the greatest distance algorithm. The greatest distance algorithm uses a different criterion for selecting new

group leaders. First, Storer and Eastman (1990) define a non-deviant cluster as one that has more than one percent of all records. A new leader is the record that is furthest from a leader of a non-deviant cluster and is greater than the average record distance from its cluster leader.

Many existing approaches use a non-hierarchical approach. Cheng et al. (2006) proposed the use of hierarchical clustering. They demonstrated that the use of hierarchical reduces the number of parameter values to test for data quality.

2.3 Limitations of Current Methods

Table 1 summarizes the parameters that need to be pre-defined for all four approaches. The shortest spanning path algorithm requires pre-specifying a threshold value. Since there is no upper bound on the value of the distance function, there are many possible threshold values to test before a desirable value is found. Worse yet, there is no systematic way to find the desirable value. The search for the most desirable parameter value may be time-consuming. A smaller threshold value than the desirable may result in a larger number of groups and possibly a larger number of outliers that are actually error free. A larger threshold value than the desirable may lead to a fewer number of groups but it may group erroneous records into existing groups along with correct records.

Table 1: Required parameters for different approaches.

Approach	Pre-specified threshold value	Pre-specified number of groups
SSP	Yes	No
LA	Yes	Yes
ARLA	No	Yes
GDA	No	Yes
HC	Yes	No

Note

- SSP shortest spanning path algorithm
- LA leader algorithm
- ARLA average record leader algorithm
- GDA greatest distance algorithm
- HC hierarchical clustering

The leader algorithm also uses a pre-specified threshold value and the number of groups to be formed. It is difficult for the user to come up with the desirable values for both the threshold value and the number of groups to be formed as the number of possible combinations for each value is very large. The desirable values can be found only after an

extensive search. The remaining two algorithms need the pre-specified number of groups. A larger number of groups than desirable produces many outliers containing no errors while a smaller number misses some outliers containing errors since we do not have prior knowledge of the data, we have to test the algorithms with many parameter values.

Neither Lee et al. (1978) nor Storer and Eastman (1990) discuss how to find the desirable parameter values among many possible values. Moreover, they do not specify any systematic methods to choose the desirable clustering result among many candidates that are generated from a given set of parameter values.

To systematically generate parameter values, Cheng et al. (2006) propose the use of a dendrogram. Their method significantly reduces the number of threshold values to test. However, Cheng et al. (2006) do not provide any justification for the use of a hierarchical approach. In this work, we try to provide a justification for our approach.

3 TSP & DATA QUALITY

In this section, we try to establish the linkage between data quality problem and travelling salesman problem (TSP). Based TSP characteristics, we develop an algorithm to deal with data quality problem.

A data record, R_i , may be represented by a vector. That is, $R_i = (x_{i1}, x_{i2}, \dots, x_{iN})$, where x_{ip} is the value of the p th field of R_i , for $p = 1, 2, \dots, N$ and $i = 1, 2, \dots, M$. A record can be classified into one of the three types (Lee et al., 1978).

Type I records: All field values in this type of record are numerical. The distance between two records R_i and R_j is defined as:

$$d_{ij} = \sum_{p=1}^N c(x_{ip}, x_{jp}) / N,$$

where $c(x_{ip}, x_{jp}) = |x_{ip} - x_{jp}| / S_p$, and (1)

$$S_p = |\max_{1 \leq i \leq M} x_{ip} - \min_{1 \leq i \leq M} x_{ip}|$$

For example, if $R_i = (4.5, 3.1, 0.9, -2.1)$, $R_j = (4.1, 2.1, 0.3, -1.1)$, $S_1 = 5.0$, $S_2 = 4.0$, $S_3 = 2.0$, and $S_4 = 2.1$, then $d_{ij} = 0.2765$.

Type II records: All field values in this type of record are non-numerical. The distance between two records R_i and R_j is defined as:

$$d_{ij} = \sum_{p=1}^N c(x_{ip}, x_{jp}) / N$$

where

$$c(x_{ip}, x_{jp}) = \begin{cases} 1 & \text{if } x_{ip} \neq x_{jp} \\ 0 & \text{otherwise.} \end{cases} \quad (2)$$

Type III records: Fields in a type III record may assume either numerical or non-numerical values. The distance between two records R_i and R_j is defined as:

$$d_{ij} = \sum_{p=1}^N c(x_{ip}, x_{jp}) / N$$

where

for a numerical field p ,

$$c(x_{ip}, x_{jp}) = |x_{ip} - x_{jp}| / S_p, \text{ and} \\ S_p = |\max_{1 \leq i \leq M} x_{ip} - \min_{1 \leq i \leq M} x_{ip}| \quad (3)$$

or for a non-numerical field p ,

$$c(x_{ip}, x_{jp}) = \begin{cases} 1 & \text{if } x_{ip} \neq x_{jp} \\ 0 & \text{otherwise.} \end{cases}$$

For example, if $R_i = (\text{black, black, 3.1, 5.0})$, $R_j = (\text{black, white, 2.1, 5.1})$, $S_3 = 4.0$, and $S_4 = 5.5$, then $d_{ij} = 0.3170$.

Lee et al. (1978), and Storer and Eastman (1990) use Euclidean distances or city block distances for type I records, and hamming distances for type II records. There is no upper bound on the value of either distance function. Therefore there are a large number of possible threshold values.

To illustrate the new distance function, consider a simple example with type III records. Table 1 is a personnel database for a hypothetical company.

Matrix (4) shows the distance value between a pair of records. Note that in the matrix, $d_{ii} = 0$ and $d_{ij} = d_{ji}$. A small distance value between two records implies that they are similar, while a large distance value means that they are different.

An erroneous record, being so different from other records, has large distance values with other records. When records are clustered into groups, erroneous records (i.e., outliers) will not be associated with other records.

Table 2: Example.

Record	POS ¹	EDU ²	MON ³	SAL ⁴
1	0	0	15	20,000
2	1	1	10	20,000
3	0	0	11	20,000
4	1	1	35	60,000
5	1	0	17	30,000
6	0	1	17	30,000
7	0	0	16	20,000
8	1	1	33	65,000
9	1	0	16	46,000
10	0	0	50	80,000

Note:

1. POS = 1, when an employee has a middle management position; and POS = 0, when an employee has a supervisor position.
2. EDU = 1, when an employee has a college degree; and EDU = 0, when an employee does not have a degree.
3. MON is the number of months an employee has worked for the company.
4. SAL is the current salary of an employee.

	Records									
	1	2	3	4	5	6	7	8	9	10
1	.00	.52	.02	.73	.29	.29	.01	.73	.34	.36
2	.52	.00	.51	.25	.32	.32	.53	.26	.36	.89
R 3	.02	.51	.00	.74	.31	.31	.03	.75	.36	.89
e 4	.73	.25	.74	.00	.43	.43	.72	.03	.39	.64
c 5	.29	.32	.31	.43	.00	.50	.29	.44	.06	.57
o 6	.29	.32	.31	.43	.50	.00	.29	.44	.56	.57
r 7	.01	.53	.03	.72	.29	.29	.00	.73	.33	.36
d 8	.73	.26	.75	.03	.44	.44	.73	.00	.39	.63
s 9	.34	.36	.36	.39	.06	.56	.33	.39	.00	.53
10	.36	.89	.89	.64	.57	.57	.36	.63	.53	.00

	Records									
	1	3	7	4	8	5	9	2	6	10
1	.00	.02	.01	.73	.73	.29	.34	.52	.29	.36
R 3	.02	.00	.03	.74	.75	.31	.36	.51	.31	.89
e 7	.01	.03	.00	.72	.73	.29	.33	.53	.29	.36
c 4	.73	.74	.72	.00	.03	.43	.39	.25	.43	.64
o 8	.73	.75	.73	.03	.00	.44	.39	.26	.44	.63
r 5	.29	.31	.29	.43	.44	.00	.06	.32	.50	.57
d 9	.34	.36	.33	.39	.39	.06	.00	.36	.56	.53
s 2	.52	.51	.53	.25	.26	.32	.36	.00	.32	.89
6	.29	.31	.29	.43	.44	.50	.56	.32	.00	.57
10	.36	.89	.36	.64	.63	.57	.53	.89	.57	.00

When we rearrange rows and columns in Matrix (4) with the purpose of putting similar records together, we may get one possible solution shown in Matrix (5). It is not difficult to observe that there are three clusters: {1,3,7}, {4,8}, {5,9}. It is also

apparent that Records 2, 6, and 10 are not associated with other records in any way. Therefore, they are the outliers.

Suppose we let a record in the row of Matrix (4) be a city in a TSP. The distance value between two records is the distance between two cities. We need to find a sequence of the cities for the row such that cities closer to each other will be placed closer together. Hence, this sequencing problem can be formulated as a TSP (Lenstra and Kan Rinnooy, 1975). Many approaches may be used to solve this TSP. In this paper, we attempt to use genetic algorithm.

4 OUR APPROACH

Our approach consists of two phases: obtaining a sequence of sample data records, and classifying records into groups. The first phase uses a genetic algorithm and the second phase adopts a classification criterion for grouping. An illustrative example is used to show the computational process of our approach.

4.1 Genetic Algorithm

The genetic algorithm approach was developed by John Holland (1975). This approach is a subset of evolutionary algorithms that model biological processes to optimize highly complex cost functions. It allows a population composed of many individuals to evolve under specified selection rules to a state that maximizes the “fitness” (i.e., minimizes the cost function).

Clearly, the large population of solutions and simultaneously searching for better solutions give the genetic algorithm its power. Some of the advantages of a genetic algorithm are that it (Haupt and Haupt, 1998):

- Optimizes with continuous or discrete parameters.
- Does not require derivative information.
- Simultaneously searches from a wide sampling of the cost surface.
- Deals with a large number of parameters.
- Is well suited for parallel computers.
- Optimizes parameters with extremely complex cost surfaces; it can jump out of a local minimum.
- Provides a list of optimum parameters, not just a single solution.

- May encode the parameters so that the optimization is done with the encoded parameters, and
- Works with numerically generated data, experimental data, or analytical functions.

4.1.1 Representation

A chromosome represents an individual. For example, $x_1 = (1011001)$ and $x_2 = (0111011)$ are two distinct individuals. Offspring (new individuals) are generated by crossover. A crossover point will be selected randomly. The parent chromosomes will be split at the chosen point and the segments of those chromosomes will be exchanged. Using this basic crossover operator, two fit individuals may combine their good traits and make fitter offspring.

Nevertheless, the simple representation scheme described above is not suitable for TSP. Instead, three vector representations for TSP were proposed (Michalewicz, 1999): adjacency, ordinal, and path. Each representation has its own genetic operators. Among the three representations, the path representation is the most natural representation of a tour. For example, a tour $3 - 4 - 1 - 6 - 5 - 2 - 7$ is simply represented by $(3\ 4\ 1\ 6\ 5\ 2\ 7)$. Our proposed approach uses this representation.

4.1.2 Initialization

Initialization involves generating of possible solutions to the problem. The initial population may be generated randomly or with the use of a heuristic. In our approach, the initial population is generated randomly.

4.1.3 Fitness Function

Fitness function is used to evaluate the value of the individuals within the population. According to the fitness value scored, the individual is selected as a parent to produce offspring in the next generation or is selected to disappear in the next generation.

In TSP, the total distance is calculated as the distance travelled from the starting city to the last city plus the distance from the last city to the starting city. In our data auditing problem, returning to the starting city (i.e., record) does not have any practical meanings. Therefore, the problem is simplified to the associated Hamiltonian Path Problem (HPP). As the first and last records need not be connected, we may calculate the total distance of a path instead of a tour in our fitness functions.

Let ρ be the permutations of records along the row of the initial matrix. For a sequence of cities

(i.e., records): (1 3 7 4 8 5 9 2 6 10), $\rho(2) = 3$ and $\rho(7) = 9$. The proposed approach converts the initial sequence of records (specified by the initial matrix) to a new sequence that minimizes the following fitness function:

$$\sum_{i=1}^{n-1} d_{\rho(i)\rho(i+1)} \quad (6)$$

where n = number of records (i.e., rows or columns).

4.1.4 Parent Selection

Parent selection is a process that allocates reproductive opportunities to individuals. There are several selection schemes: roulette wheel selection, scaling techniques, ranking, etc. (Goldberg, 1989).

As the process continues, the variation in fitness range will be reduced. This often leads to the problem of premature convergence in which a few super-fit individuals receive high reproductive trials and rapidly dominate the population. If such individuals correspond to local optima, the search will be trapped like hill climbing.

In our approach, fitness ranking is used to solve the problem of premature convergence (Whitley, 1989). Individuals are sorted according to their fitness values, the number of reproductive trails are then allocated according to their rank.

4.1.5 Crossover

Several TSP crossover operators are defined: partially-mapped (PMX), order (OX), cycle (CX), and edge recombination (ER) crossover. Whitley et al. (1989) found that ER is the most efficient crossover operator for TSP. Starkweather et al. (1991) proposed an enhancement to ER and find it more efficient than the original operator.

In our approach, we use the EER operator. Since the EER operator incorporates random selection to a break tie, this mechanism creates an effect similar to mutation. In our approach, we do not use any mutation operator.

4.1.6 Mutation

Mutation is applied to each child individually after crossover according to the mutation rate. It provides a small amount of random search and helps ensure that no point in the search space has a zero probability of being examined. Several mutation operations have been suggested by Michalewicz (1999). We do not plan to use mutation operation. This is because the

crossover operator used incorporates a random selection in completing a legal permutation and the effect is similar to a mutation.

4.1.7 Replacement & Termination Criterion

In each generation, only two individuals are replaced. In other words, parents and offspring may co-exist in the population. The genetic process is repeated until a termination criterion is met. In this case, we use a pre-specified maximum number of generations as a termination criterion.

4.2 Classification Criteria

We adopt the classification criteria developed by Stanfel (1983) to classify data records into groups. These classification criteria seek to minimize the average distance within groups and maximize the average distance between groups. Minimizing the average distance within groups will put similar data records into the same groups. At the same time, maximizing the average distance between groups will put dissimilar data records into different groups.

To formulate the chosen selection criterion, we define:

$$Y_{ij} = \begin{cases} 1 & \text{if records } i \text{ and } j \\ & \text{are in the same group} \\ 0 & \text{otherwise} \end{cases} \quad (7)$$

The expression for the average distance *within* groups is given as:

$$\frac{\sum_{i=1}^{M-1} \sum_{j=i}^M d_{ij} (1 - Y_{ij})}{\sum_{i=1}^{M-1} \sum_{j=i}^M (1 - Y_{ij})} \quad (8)$$

While the expression for the average distance *between* groups is given as:

$$\frac{\sum_{i=1}^{M-1} \sum_{j=i}^M d_{ij} Y_{ij}}{\sum_{i=1}^{M-1} \sum_{j=i}^M Y_{ij}} \quad (9)$$

Hence, in order to achieve the objective of maximizing the homogeneity of records within groups as well as the heterogeneity of records between groups, the difference between the average

$$\frac{\sum_{i=1}^{M-1} \sum_{j=i}^M d_{ij} Y_{ij}}{\sum_{i=1}^{M-1} \sum_{j=i}^M Y_{ij}} - \frac{\sum_{i=1}^{M-1} \sum_{j=i}^M d_{ij} (1 - Y_{ij})}{\sum_{i=1}^{M-1} \sum_{j=i}^M (1 - Y_{ij})} \quad (10)$$

distance within groups and the average distance between groups is minimized as shown in criterion (10):

4.3 Illustrative Example

Phase one of the proposed approach takes Matrix (4) as input and rearranges its rows and columns to obtain a sequence of records. As shown in Matrix (5), the sequence produced is (1 3 7 4 8 5 9 2 6 10). Phase two takes the sequence and classifies data records into groups to minimize criterion (10). As a result, we find this grouping result: {1,3,7}, {4,8}, {5,9}, {2}, {6}, {10}.

Unlike most of the existing methods, our approach does not require any parameters. The generation of a sequence of data records and assignment of data records into groups are all automatic.

Based on the grouping result, we conclude that records 2, 6, and 10 are outliers. A survey administrator will examine these outliers and determine whether they contain any error or not. Our example postulates that an employee with a college degree, in a middle management position, and with more years of seniority should have higher current salary. On the other hand, an employee without a college degree, in a supervisor position, and with fewer years of seniority should have a lower current salary. However, record 2 indicates that the employee has a college degree and in a middle management level but has relatively low current salary. Record 10 indicates that the employee has exceptionally high current salary for his/her position (i.e., supervisor) and education (i.e., does not have a degree). Record 6 may (or may not) contain errors.

5 CONCLUSION

In this work, we discussed the use of clustering algorithms for assessing the quality of Internet survey data. Limitations of some existing approaches were identified. To address these limitations, we first examined the nature of the quality problem of Internet surveys and then established that the problem is equivalent to a TSP. Our proposed approach exploits the underlying TSP structure. Although many algorithms may be used for the TSP-quality problem, we adopt genetic algorithm for computational efficiency and advantage. Compared to the existing approaches, our approach provides a better understanding of the nature of the Internet survey problem and seems to

offer improvement potential. However, the quality model must be implemented and tested to verify our claims.

REFERENCES

- Anderberg, M.R., 1993. *Cluster analysis for applications*, Academic Press, New York.
- Cheng, C.H., Goh, C.H., and Lee-Post, A., 2006. Data auditing by hierarchical clustering, *International Journal of Applied Management & Technology*, Vol.4, No.1, pp. 153-163.
- Felligi, I.P. and Holt, D., 1976. A systematic approach to automatic editing and imputation, *Journal of American Statistics Association*, Vol. 71, pp. 17-35.
- Freund, R.J. and Hartley, H.O., 1967. A procedure for automatic data editing, *J Journal of American Statistics Association*, Vol. 62, pp. 341-352.
- Goldberg, D.E., 1989. *Genetic Algorithms in Search, Optimization and Machine Learning*. Massachusetts: Addison Wesley.
- Hartigan, J.A., 1975. *Clustering Algorithms*, McGraw-Hill, New York.
- Haupt, R.L. and Haupt, S.E., 1998. *Practical Genetic Algorithms*. New York: John Wiley & Sons.
- Holland, J.H., 1975. *Adaptation in Natural and Artificial Systems*. Michigan: Michigan Press.
- Lee, R.C., Slagle, J.R., and Mong, C.T., 1978. Towards automatic auditing of records, *IEEE Transactions on Software Engineering*, Vol. SE-4, pp. 441-448.
- Lenstra, J.K. and Kan Rinnooy, A.H.G., 1975. Some Simple Applications of the Traveling Salesman Problem, *Operations Research Quarterly*, Vol. 26, pp. 717-733.
- Michalewicz, Z., 1999. *Genetic Algorithms + Data Structures = Evolution Programs*. Third, Revised and Extended Edition, Hong Kong: Springer.
- Naus, J.I., Johnson, T.G., and Montalvo, R., 1972. A probabilistic model for identifying errors and data editing, *Journal of American Statistics Association*, Vol. 67, pp. 943-950.
- Slagle, J.R., Chang, C.L., and Heller, S.R., 1975. A clustering and data-reorganizing algorithm, *IEEE Transactions on Systems, Man, and Cybernetics*, Vol. SMC-5, pp. 125-128.
- Stanfel, L.E., 1983. Applications of clustering to information system design, *Information Processing & Management*, Vol. 19, pp. 37-50.
- Starkweather, T., McDaniel, S., Mathias, K., Whitley, D., and Whitley, C., 1991. A Comparison of Genetic Sequencing Operators, *Proceedings of the fourth International Conference on Genetic Algorithms and their Applications*, pp.69-76
- Storer, W.F. and Eastman, C.M., 1990. Some Experiments in the use of clustering for data validation, *Information Systems*, Vol. 15, pp. 537-542.
- Whitley, D., 1989. The Genitor Algorithm and Selection Pressure: Why Rank-based Allocation of Reproductive Trials Is Best, *Proceedings of the Third International Conference on Genetic Algorithms*, pp.116-121.