MODEL SHARING IN THE SIMULATION AND CONTROL OF DISTRIBUTED DISCRETE-EVENT SYSTEMS

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Abstract: Today, sophisticated discrete-event systems are being designed whose complexity necessitates the employment of distributed planning and control. While using a distributed control architecture results in the overall system model consisting of a collection of independent models, today's commercially available simulation languages can only accommodate a single model. As a result, in order to use these simulation languages one must create a new system model that consists of a single model but yet models a collection of models. Typically the communication among the distributed models is ignored causing inaccurate results.

In this paper we use our simulation concept, also presented in this paper, to create a simulation tool that enables the simulation of distributed systems by using a collection of models rather than a single model. With our concept we create a methodology that accomplishes this by simulating the communications among the distributed models. Besides the benefit of not having to create a new model for simulation, this methodology produces an increase in accuracy since the communication among the models is taken into consideration. Furthermore this tool has the capability to control the system using the same collection of models.

1 INTRODUCTION

Flexible manufacturing systems represent an interesting control problem because they generally have its control architecture distributed among many different computers. While creating a controller for each subsystem is not difficult, coordinating the individual controller to perform a common task may be more challenging. Manufacturing systems are typically modeled using discrete-events. And discrete-event simulation is commonly used to evaluate possible alternative actions used in decision making. This leads to a modeling problem. Since the controller of the plant is distributed, that is it consists of many controllers running independently, and the simulation model must not be distributed how is one to model a system using a single model when the actual controller consists of a collection of models? Currently the solution is to create a new model that captures the control logic of the complete controller system. This of course requires writing a new model that hopefully captures this true control logic correctly including all of the communication among the controllers. In addition to representing a significant modeling effort, it leads to a source of error. This is especially true when you consider modeling the communication. In this paper we present a modeling approach that has the capability of accepting a collection of models. We can use the models that are used for control directly by simply including them into the simulation. Not only can our simulation handle the collection of models but it also automatically models the communication among the controllers. This paper is based on the simulation and control of distributed large complex systems whose behavior is characterized by discrete events such as a flexible manufacturing system, FMS. FMS produce many part types concurrently in very small quantities. These systems are difficult to control because they generally operate in a highly transient state.

The following research efforts are related to our work. (Peters et al. 1996), (Smith et al. 1994) and (Smith and Peters 1998) have adapted Arena (Kelton, Sadowski and Sadowski 2001) to control their experimental FMS. However, unlike the modeling approach to be discussed in this paper, their control architecture uses one hierarchical level where a single supervisor, the cell controller,
manages a set of subordinate processes. In their modification of Arena, they have included special events in order to facilitate message passing among the controllers. Furthermore, to model a complex FMS or other multi-level, distributed system, the set of modeling elements provided by ARENA, as well as most simulation languages, severely constrains the modeling process. This is particularly true when one attempts to assess the impact of the control architecture upon the system. (Davis et al. 1993) and (Davis 1998) details the numerous restrictions that current simulation languages impose upon the modeling of hierarchically distributed systems. (Mize et al. 1992) further discusses the inaccuracies that ensue from using current simulation languages in order to model an FMS. (Narayanan et al. 1992), (Govindaraj et al. 1993) and (Bodner and Reveliotis 1997) have also developed simulation tools that are capable of controlling a system. They claim that the typical approach to address the complexity of FMS control systems is a hierarchical decomposition. A widely used decomposition features tree layers, Strategic Decisions, Tactical Decisions and Operational Decisions, (Arango and Prieto-Diaz 1991) and (Stecke 1985). In (Bodner and Reveliotis 1997), (Govindaraj et al. 1993), and (Narayanan et al. 1992), they claim that to handle lower level issues involving the real-time control, two additional layers are needed. The Structural and Equipment Control layers address issues such as deadlock avoidance. In these papers, hierarchical decomposition refers to the logical decomposition of the decision making process while we refer to it as the actual decomposition of the controller that yields a distributed controller. Incidentally they also claim that FMS exist but there are no controllers for them.

The notions of multi-resolutional architectures and task decomposition have been discussed by many. The interested reader is referred to (Albus and Meystel 1995 and 1997) for a discussion of these terms. In the same reference, these authors discuss their reference architecture for distributed intelligent control and the associated real-time control system for its implementation. The major distinction between their architecture and ours, discussed below, is that they have separated the planning and control functions within a given controller. Our approach does not separate these functions. Secondly, we make heaver use of on-line simulation.

2 THE DISTRIBUTED MODELING METHODOLOGY

This methodology presents a very novel approach to simulating distributed systems using a single thread. It is assumed that the control architecture is hierarchical. Figure 1 shows a typical control architecture for an example system.

![Control Hierarchy](image)

Figure 1: The hierarchical architecture used in our physical model

This methodology is based upon the belief that the interactions among the controllers must be considered by the simulation model in order to accurately model a system with a distributed control architecture. The single most important characteristic of the methodology and what separates it from other object-oriented simulation approaches is its attention to modeling the flow of messages among the controllers included within the architecture. By modeling the flow of messages, the methodology allows the simulation to accommodate any number of models. Recall that the controller is distributed. That means that there are many independent controllers each running with its own model. This methodology allows us to simulate the complete system using the models from the controllers (Gonzalez 2004). This has many advantages:

1. The simulation is more accurate since it includes all of the communication among the distributed controllers. This was the original motivation for developing our methodology.

2. The simulation produces maximum fidelity since the same models are used for both simulation and control. Model verification is a major effort in modeling usually requiring more effort then to actually build the model. Furthermore there is always the risk that the model does not represent the control logic correctly. Using our method reduces the chances of discrepancies between the models used for simulation and those used for control. This simplifies the verification phase of
modeling and produces a model that represents the true logic.

3. The modeling effort is shared between control and simulation. They both share the same models so only one set is made. In fact without our methodology many more models will have to be created. Since each controller has a different set of subordinate controllers in its control domain, each one will need a different model. Only the model in the very bottom row will be able to use there exiting models. Each model will need to include its own logic as well as the logic of all of its subordinate controllers.

4. The control and simulation models necessarily employ the same state definition since they are in fact the same code. This simplifies the task of initializing the simulation model to the current system-state. If the simulation and control models employ different state definitions, which is the case when one employs conventional simulation approaches, then one must translate the measured system variables into values for the state variable employed within the simulation model. This is a relatively large effort (Gonzalez and Davis 1998a).

5. This methodology offers multiresolutional modeling for simulation. Since the system models have a hierarchical organization, the simulation resolution can be dynamically selected by selecting the levels in the hierarchy to include models for. For example, in Figure 1, if the top most controller wants to make a fast but rough decision it may only include the next level down. This increases the simulation speed at a cost of accuracy. Furthermore if a controller on the 2nd level wants to make a decision it will only include the 3 controllers below it since the others have no relevance to its decision. Since the models are independent, to vary the resolution we simply include the desired models with no extra modeling effort required.

3 THE SIMULATION AND CONTROL TOOL

In order to use the control models in the simulation directly a tool needs to be used that can run in control and simulation mode, in this way the same model, written for the same tool, can be used for both control and simulation. The state definition can then be simply the contents of the model’s variables without any need for interpreting the information in these variables. We have developed a tool based on this methodology. This tool is described in more detail in (Gonzalez and Davis 2003). The goal for developing this software tool is to provide a simulation approach using C++ that provides the modeling convenience of a conventional simulation language while providing the additional modeling capabilities that are needed to control a real-world system, particularly hierarchically distributed systems. A set of C++ objects was developed in order to provide the basic necessary modeling elements that are employed in nearly all simulation and control scenarios, irrespective of which modeling methodology is being employed.

In this tool each included modeling element is represented as an object in C++. As with most simulation approaches, the adopted simulation approach is event-driven, and an executive object manages the sequential processing of events. Two event lists (discussed below) are maintained where the events are stored as they wait to be executed at the proper simulated time. The simulation object, Figure 2, is a C++ object that contains the scheduled event list, pending event list, list of resources, the executive function, a pointer to the system model, and the communication object that interacts with the hardware. The executive function manages the chronologically ordered scheduled event list and the pending event list. It manages the processing of the events as they are removed from the event lists, schedules new events that are to be placed into either the scheduled or pending event list, and manages the allocation of resources. After each event is processed, the simulation object removes the next event with the smallest event time from the event list and then invokes the proper object to manage the event’s execution.

Using a streamlined approach, there are two basic types of events that can occur. The QUESEIZE event occurs when resources become available for assignment to a requesting entity. The
DELAY event occurs whenever a delay is completed. There are also other types of events that provide capability for control but which are not considered in most simulation tools, including the PENDING, TIMEOUT, and ERROR events.

In order to run, the system model must be attached to the simulation object. A communication function must also be provided to the simulation object in order to permit this object to communicate with the hardware when the model is running in a control mode. This function handles all of the communications among the distributed controllers and the various machines.

The simulation object can operate in the following two modes:

Simulation mode: In this mode, the system clock advances in discrete increments. Every time an event is pulled off the event list, the time is incremented to the event’s time of occurrence. No hardware is present. The model executes purely as a software program.

Control mode: In control mode, the system clock runs in real time. That is, the system clock is a conventional clock, like the one on the wall, where time advances continuously. The events are pulled off the event list when their event time occurs. Hardware may be present. The pending event list is used to allow the executive function to determine when events occur.

4 DISTRIBUTED MODELING

In order to build the complete simulation model, one simply includes all of the simulation objects into a single program along with the coordination object. The coordination function within the coordination object coordinates the execution of all of the simulation objects. It also has the global event list and the message relay discussed later. When the distributed model is executed upon a single computer, only one of the simulation objects can be executed at a given time because there is only one computational thread. In order to emulate all of the simulation objects operating concurrently on a single processor, the coordination function executes one simulation object for a short time and then switches to another object. The coordination function uses its global event list and its message relay to determine which simulation object to execute next. The individual simulation objects return control back to the coordination function when it is done executing all of the events that occur during that instance of time. Note that in simulation, the time it takes the computer to execute the segment of event code that handles is usually neglected, and an event can be assumed to occur in an instance of time. Furthermore, if an event triggers the immediate execution of another event, then this second event is assumed to occur at the same instant of time. This is why the coordinating function has the liberty to give control to an individual simulation object for the duration that it takes that object to execute an event while meeting the stated objective of emulating concurrent operation.

In a nondistributed simulation of the system, the executive function constantly cycles in a loop. This loop starts by checking the event list for the next event. It then executes this event. After executing the event, it checks to see whether there are any new events resulting from the execution of this event. An event may cause a second event to execute by the releasing of resources. Thus, the execution of an event may cause a chain reaction. All the events in this chain are executed within this single executive function cycle. In control mode, at the end of the current cycle, the executive function checks to see whether any message has arrived from the communication pipelines. If a message has arrived, the pending event list is then checked to find the proper event to be executed. In either mode, the executive function then initiates a new event-processing cycle.

In order to permit several executive objects to operate concurrently while in the control mode, the executive function is modified so that it cycles through the loop only once every time it is called. Furthermore, it does not check the communication pipelines at the end of the cycle. All the communication is handled by a message relay contained within the primary executive function. The executive function is responsible for calling the executive function of the model that will execute the next event. At this time, the submodel’s executive function cycles through the event-processing loop once and then returns program control back to the coordinating function. This method of sharing the computing processor operates on a similar principle as many multi-tasking operating systems.

Another special feature of our simulation tool is the manner by which messages enter the model. Instead of checking the communication lines every time the executive function reaches the end of an event processing cycle, the lines are not checked at all. Rather, the coordination function manages the flow of messages. Whenever a message is present, the coordination function calls the executive function of the appropriate submodel’s object and provides it with the message. The executive function of that object then handles the message by executing the proper events in the same way that it previously
handled a message when it was responsible for detecting the message arrival. We adopted this approach to ensure that a single entity, namely the coordination function, was responsible for monitoring the arrival of all messages.

The last feature in our simulation software is the DELAY modeling element. When a DELAY element is executed, in addition to scheduling a DELAY event onto the submodel’s scheduled event list, it also automatically schedules a DELAY event into the coordination object’s global event list as well. Using this approach, the coordination function maintains a chronologically ordered, global event list with all of the events associated with every simulation object. The coordination function uses this global event list to determine which simulation object will be invoked next. Since the scheduling of events into this global event list occurs automatically, the user does not need to consider this global event list when modeling a local delay. From the user’s point of view, the delay is modeled in the same way as if the model were not a member of a collection of distributed simulation objects. As a matter of fact, no part of the submodel has to be modified to accommodate the distributed simulation. Each submodel executes as though it is running on its own dedicated computational thread. This is important because it allows the modeling engineer to implement a distributed model the same manner as when the model was not distributed. Again this feature is essential because the model will be used both to project the response and to control the system.

5 THE COORDINATION FUNCTION

The coordination function is a supervisory function that manages the execution of all of the submodels’ simulation objects. The coordination object contains a global event list and a message relay. This function runs in a cycle much the same way the simulation objects do when running independently (not as part of a distributed simulation). In the typical simulation mode, the event-processing cycle starts by removing the next event from the global event list and passing it to the appropriate submodel for execution. While a submodel is being executed, the message relay, which acts as the network, receives all of the messages that are generated by the submodel that is currently executing. These messages are stored in the rear of the message queue within the message relay, which is used to model the network. Once the submodel finishes executing its current event-processing cycle, the coordination function removes the first message from the front of the message queue and passes it to the recipient submodel. The submodel then receives the message and executes the appropriate functions that are needed to handle the message. Additional messages may again be generated and are inserted at the rear of the message queue. Once control is returned to the coordination function, it removes the next message from the front of the message queue and recycles it though the message-processing loop. This procedure is repeated until no messages remain in the message queue. At this time the coordination executive function finishes its cycle and begins the next cycle by removing the next event in the global event list and passing it to the appropriate simulation object.

The events in the global event list tell which submodel will address the event and the time at which the event occurs. When an event is pulled off the list, the simulated time is then advanced to the next event’s event time. The event type is not recorded on the global event list, as this information is contained within the submodel’s local event list and need not be duplicated. The coordination function then calls the appropriate submodel and passes it the current simulated time. The submodel, knowing that it is responsible for executing the next event, pulls the next event off its local event list and executes it. As stated above, it only executes its executive function’s event processing cycle up to the point where the communication lines are to be checked. At this point, program control is returned to the coordination function where the function then checks the message relay and continues its cycle. See Figure 3 above.
6 AN ILLUSTRATIVE EXAMPLE

The following section illustrates an example of a real physical system that is controlled using a distributed model. A simulation using the same set of distributed models for the considered system is also performed. Each of the sub models for the included controllers was programmed using the simulation tool discussed above. In Gonzalez and Davis (1998b), the physical system is discussed in greater detail while in Gonzalez and Davis (1997) the employed simulation model is discussed.

The physical system is a model of a flexible manufacturing system. The schematic for the constructed FMS emulator is shown in Figure 4 and a photograph is depicted in Figure 1. This model provides a testbed for the development of the simulation and control tool that is needed to manage the system. Note that since we are assuming the system is a real FMS, we thus have modeled it using a collection of independent models.

![Figure 4: Schematic for the constructed FMS physical model.](image)

The emulator has four Processing Centers, numbered 1 through 4. Each Processing Center (PC) contains one primary processing resource and a dedicated Material Handling System (MHS). Within the emulated FMS (see Figure 3), another process is the Fixturing Center (FC). The FC has a dedicated MHS consisting of a primary carousel capable of holding sixteen jobs and two smaller carousels for loading and unloading jobs from the Automatic Guided vehicle (AGV). The movement of these carousels is controlled by a dedicated controller. The FC has two fixturing positions that represent the subordinate unit processes. The final subordinate process is the cell’s MHS. AGVs are employed as the primary material handlers at the cell level and are modeled with an HO-scale electric train. In this layout, there are over forty track segments that can be individually powered. Sensory switches are provided on each track segment to detect the presence of an AGV.

![Figure 5: Photograph of the hardware model](image)

In constructing this control architecture, 25 independent copies of the simulation tool was employed, each with its own model of the subsystem that it is addressing and all running concurrently, see Figure 1. The only thing that ties them together is the communication among them. The communication is performed across a local area network (LAN) connecting seven computers where the cell controller and each of the cell’s six subordinate controllers are situated on their own computer. Additional communication links are provided via RS-232 links between the FC and PC controllers and their dedicated hardware controller boards.

In this example, the developed model controlled the physical system. The controller was given a total of three jobs, each with two processing steps. Each processing step required the part to be moved to a different machine. In addition to the processing of the part, before and after each processing step the part had to be moved to the fixturing center for fixturing. There were four Sun Work stations and six controller boards. The software controllers were distributed among the four workstations. Each controller board controlled the hardware that was attached to it. The same model was employed to run the simulation in order to project the future performance of the system given its current state. The state of the hardware is used to initialize the state of the simulation models. However since they are in fact the very same software code the state transfer simply represents a straightforward initialization of each variable in the simulation model from the corresponding variable in the control model. No interpretation as to the physical meaning of these variables data is necessary. The message transcript for both the actual run and the simulation run were similar. The only difference was due to the
discrepancy between the estimated and actual duration times. No changes had to be made to any of the models in order to switch from using them for control or simulation.

Since there is no tool that can accept a collection of models, simulating this system with commercial simulation software like ARENA will be impossible. The method that is commonly used to model distributed systems is to create a new equivalent simulation model that is not distributed. This however represents a considerable modeling effort. For example, the complete modeling effort for our system, including the models for control as well as the simulation models, is cut by 72% over using a conventional modeling tool. Using a conventional tool would have resulted in a modeling effort equivalent to writing models for 88 units where as we only wrote 25 models, one for each of the 25 units. And this does not include the extra modeling effort that goes into modeling the communications among the controllers, which, for our system is integrated into the tool and therefore represent no additional modeling. If one wants to incorporate simulation at different resolutions then this requires additional modeling using a commercial tool. Our tool provides all levels of resolution using the basic set of models. One simply includes those models whose logic we want represented.

7 CONCLUSIONS

The distributed simulation and control system was successfully tested on our physical simulator of an FMS consisting of 4 single processor machines, 1 double processor machine and an automated guided vehicle system. For the sake of demonstration the system used 6 independent computers to control the hardware. Real-time simulations of the distributed controller using our methodology were concurrently executed using the same control models. The simulations were initialized to the current state of the hardware before starting. This allows the simulations to produce future predictions.

We have shown an implementation of a tool that along with a coordination object can simulate a distributed system using the individual models. The key issue for this implementation to work is our distributed modeling methodology where each individual simulation object is included into the simulation program and the coordination function coordinated the execution of each simulation object to model the concurrent processing. The coordination function uses its global event list to organize the execution of the simulation object by chronological order. The simulation objects in turn only execute the immediate event and do not cycle.

One of the most important advantages of using the same models for both simulation and control is that the state definition is simply the data in the variables of the models. This allows for the simulation to be initialized with the current state of the controller with only a trivial copying of the data to a 2nd copy of the model. This solves a major problem in control. That is, the state does not need to be interpreted to meaningful information to initialize a different model.

REFERENCES


