

TRANSFORMING SA/RT GRAPHICAL SPECIFICATIONS INTO CSP+T FORMALISM

Obtaining a formal specification from semi-formal SA/RT essential models

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Abstract: A correct system specification is systematically obtained from the essential user requirements model by applying a set of rules, which give a formal semantics to the graphical analysis entities of SA/RT. The aim of the systematic procedure is to establish the methodological infrastructure necessary for deriving a complete system specification of a given real-time system in terms of CSP+T processes. A detailed complete solution to the Production Cell problem is discussed, showing how the method can be applied to solve a real-world industrial problem.

1 INTRODUCTION

We present a complete bottom-up systematic method to derive a correct system specification from a semi-formal system requirements specification in SA/RT (Hatley 1988, Ward 1985, Svensson 1991) by systematically applying a set of transformation rules. The method integrates two complementary approaches to describe a real-time system: (1) SA/RT based notations, and (2) CSP+T process terms (Žic, 1994) to model real-time processes, including the specification of their timing requirements.

The approach takes advantage of the long tradition of SA/RT graphical notations and development methodologies in the industry and, at the same time, aims to foster the use of Process Algebras as an adequate way to overcome the intrinsic imprecision that SA models present in describing real-time systems.

Many proposals have tried to overcome the lack of formal semantics of SA notations. Noteworthy among these is the formalization of SA through Z and Larch (Semmens, 1990), the *translation* from SA to *Communication Processes* formalism (Fencott, 1994) and the set of rules to give SA an interpretation by using High-Level Petri nets (Elmstrom, 1993). However, as Baresi and Pezzè stated (Baresi, 1998), all of these proposals irremediably damage the flexibility of SA by assuming a given interpretation to the ambiguity of analysis entities,

some of which should only have a weak semantics in order to remain useful as constructs of a process description language.

The objective of our approach follows the guideline proposed in the latter reference, namely to overcome the imprecision and ambiguities that the different *families* of SA notations present in describing real-time systems. However, the proposed method does not determine a particular semantics when there are several possible ways to solve a given ambiguity in an SA analysis entity. It is left up to the analyst to select the most appropriate notation semantics, depending on the system that needs to be specified. This feature of the method is a result of the flexibility provided by the CSP+T design notation for real-time systems.

The proposed systematic derivation technique and the transformational rules can be easily integrated into state-of-the-art SA/RT software tools and the complete derivation process can be fully implemented in Java with the support of CTJ (Hilderink, 2000) or JCSP (Welch, 2001) libraries. The rest of this paper is structured as follows: first, we give some background on SA models, which is necessary to understand the transformation rules. In section 3, we describe the system specification method. In section 4, using the example of the Production Cell, we present a complete system specification. Finally, the conclusions and ongoing work lines are presented.

2 REAL-TIME STRUCTURED ANALYSIS

The methodologies and notations referred to under the generic denomination of Structured Analysis (SA) are mainly directed towards specifying the system behaviour as a set of *data transformations or processes*, which describe the basic functions of the target system. The first was proposed by Ward and Mellor (Ward, 1985) and the second by Hatley and Pirbhai (Hatley, 1988). A third variation, called the *Extended System Modeling Language (ESML)* (Svensson, 1991) is an attempt to combine these two approaches to the structured analysis of real-time systems.

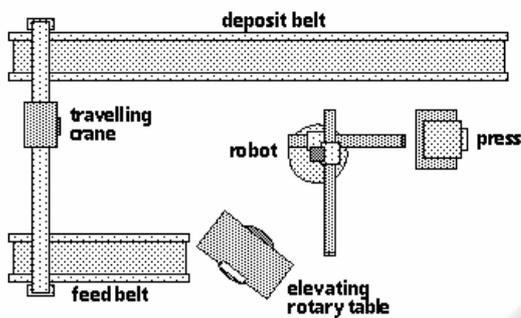


Figure 1: The Production Cell.

2.1 Production Cell Example

This well-known case study (Lewerentz, 1995), Fig.1, presents a realistic industry-oriented problem, where safety requirements play a significant role and can be met by the application of formal methods. In the fundamental configuration, the production cell processes metal blanks that are conveyed to a press by a feed belt. The first robot arm takes each blank from the feed belt and places it in the press. Because the belt and the robot are at different heights, there is an elevating rotating table which is designed to give blanks to the robot. The press forges a new metal blank and opens again. Forged metal plates are taken out of the press and put on a deposit belt by a second robot arm. Since the robot is fitted with two arms, the utilization of the press is enhanced, thus making it possible for the first arm to pick up the next blank while the press is forging another plate with the previous blank.

2.2 System Requirements Model

The model consists of a hierarchy of transformation schemes rooted on the *System Context Diagram (SCD)*. Each scheme “explodes” into a *State Transition Diagram (STD)* or into a *Data Flow Diagram (DFD)*. The scheme denoted as SCD defines the border between the system and the environment, comprising the external entities (or *terminators*) to the system. Fig. 2 shows an example of SCD. DFDs may explode into new, more detailed DFDs.

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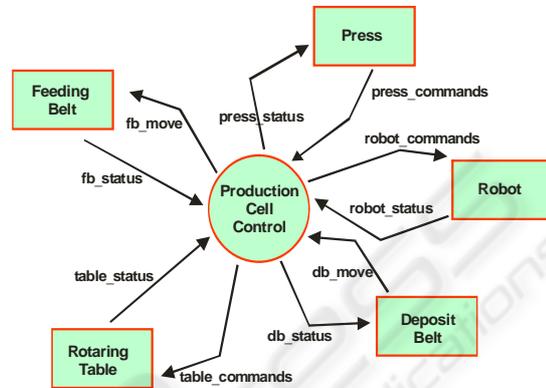


Figure 2: Production Cell SCD.

The DFD transformation scheme in the SA/RT notations must include at least one *Data Transformation Process (DTP)*, whose role is to change the input data or *event flows (control)* into output flows with no relation between the number of inputs and outputs. The same output can be sent to several analysis entities and a DTP must have at least one output flow.

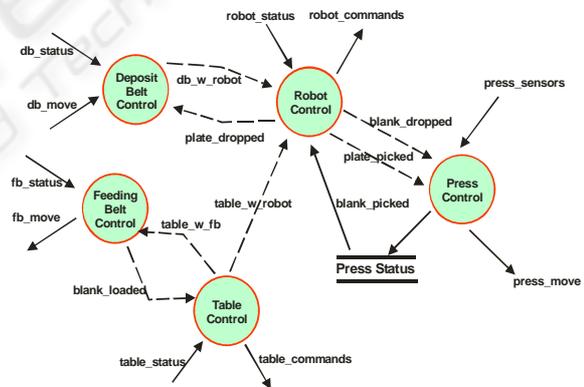


Figure 3: First level Production Cell DFD.

Control Transformation Processes (CTPs) serve to transform inputs into output event flows. They cannot accept or generate any type of data flow and are formally specified by means of a *State Transition Diagram (STD)*. STDs should be deterministic Moore or Mealy automata, and they describe a sequence of *state transitions* of the system that cause the execution of DTPs to be triggered. Different output events can be specified in an STD to activate the DTPs it controls. Enable/disable events indicate that the DTP will execute between the enable and disable signals. Trigger events indicate that the time needed

by the DTP to perform its action is unknown. These types of events cannot be used to model terminator activation, since event flows towards them must only be signals that switch on the device. Data stores (DS) loosely represent data of a certain type that cannot be considered structured. Destructive readings over a DS cannot be assumed and DS cannot be directly inter-connected, since the movement of data is only performed by the DTP in diagrams. The DFD for the complete Production Cell Control System is shown in Fig. 3.

2.3 Flaws of SA/RT as a Specification Notation for RTS

Some WM and HP analysis entities do not have a concrete semantics, thereby causing imprecision in the specification, which may confer non-predictability on a final real-time system at a later development stage:

- a) *Lack of any rule for defining primitive process specifications (PSPECs) in DTP schemes.* PSPECs are purely functional descriptions. However, in realistic applications, DTPs not only describe the purely functional behaviour of processes but often also include control and timing information.
- b) *The enabling conditions of processes are not fixed.* The SA rationale is that processes are enabled whenever “sufficient data” appear in any of their input flows. Nevertheless, the rules do not clearly indicate the expected behaviour of a process when more than one of its input flows have values.
- c) *Execution time requirements for processes are excluded.* These requirements, when applied to practical cases, are used to specify a maximum or minimum time to be associated with the execution of a process.
- d) *The number and the type of the input flows entering a process are described in vague terms.* To obtain predictability in real-time applications, when there are multiple input flows entering a process, it is necessary to define whether all the inputs (synchronous) or only a subset (asynchronous case) are needed.
- e) *Simultaneous events awakening more than one transition.* This possibility is excluded from SA/RT notations. However, there should be no objection to allowing nondeterministic selections of transitions in notations for soft real-time systems. Such imprecisions can be solved by giving a semantic interpretation to the SA entities that exclude any of these ambiguities. These interpretations can be easily programmed in CSP+T by using a set of rules which translate each SA entity into a pattern that defines a CSP+T process.

2.4 Real-time Systems Specification with CSP+T

Many proposals have tried to overcome the problem of SA imprecision by complementing it with formal methods. The use of extensions of algebraic process description languages, such as CSP (Hoare, 1985), CSP+T (Žic, 1994), the standard specification language LOTOS (Eijk, 1989), can give a precise and flexible interpretation to SA entities.

In the group of CSP derivatives to describe time intervals, we could mention *Timed CSP* (Hoare, 1985) and CSP+T, the latter being a simpler approach, although still powerful enough to formally describe a set of deterministic processes with time constrained behaviour. The syntax of CSP+T has been adapted to our method so as not to include nondeterministic operators for the moment, since real-time controllers are deterministic pieces of software, due to predictability is preferred to programmability in these systems. The adapted notation is described as follows:

-Every process P defines its own set of communication symbols, termed the *communication alphabet* $\alpha(P)$. These communications represent the events that the process P receives from its *environment* or that occur internally, such as the event τ that is not visible in the environment. Any type of event causes a change of state of the process.

-The communication interface $comm_act(P)$ of a given process P contains all the CSP-like (Hoare, 1978) communications ($\{?, !\}$) in which P can engage and the alphabet $\alpha(P)$.

-A new operator \star (star) denotes process instantiation. Given P' , the timed version of P , which is instantiated at time l , s is a time stamp associated to a , and the specification of P' is,

$$P' = l. \star \rightarrow s.a \rightarrow STOP, \text{ where } s \in [l, \infty)$$

The instantiation event is unique in the system, since it represents the origin of time at which the processes can start their execution.

-A new event operator $\triangleright \triangleleft$ is introduced, to be used jointly with a *marker* variable to record the time instant at which the event occurs. $ev \triangleright \triangleleft v$ means that the time at which ev is observed is in the *variable* v . The value of time stamps is taken from the set of positive real numbers, so that successive events form a non-decreasing monotonic sequence.

$$P = l. \star \rightarrow a \triangleright \triangleleft var \rightarrow STOP$$

For each execution of P , the time stored in the variable will always satisfy $var \geq l$. The scope of marker variables is limited to one sequential process.

-Each event is associated with a time interval, which is called the *event-enabling interval*.

$$P = 0. \star \rightarrow [1, 2] a \triangleright \triangleleft v \rightarrow STOP$$

The value of the marker variable v satisfies the inequality $1 \leq v \leq 2$. Only during this continuous time interval is the event available to the process and its environment. A process is considered to be the *STOP* process if it cannot engage in any communication or synchronize in any event within the interval that precedes the event.

-The enabling intervals can also be defined in terms of functions over a set of marker variables,

$$P = \dots E.P' \quad E = \{s \mid s \in \text{rel}(x, v)\}$$

The bound variable x sets the upper limit of the interval. If the preceding event occurs at time t_0 , then $\text{rel}(x, v) = [v-t_0, x+v-t_0,]$, since the times for events are absolute and the times for intervals are relative to the preceding event. When there are no marker variables referenced, the enabling interval is defined relative to the immediate preceding event.

-Finally, it should be noted that only deterministic processes can be described in *CSP+T* formal description language.

In order to obtain a *CSP+T* model of the system, it is necessary to represent every analysis entity of the System Requirements Model (SRM) by a class of *CSP+T* processes. Following this approach, we intend to write a process *CSP+T* prototype for every DTP, CTP, DS, CS, continuous data flow, etc.

3 A FORMAL SPECIFICATION FROM THE SRM

A series of transformation rules will allow us to create a *CSP+T* model for every transformation scheme that appears in any diagram of the SRM.

Definition 1. Given the set of SA/RT analysis entities E , *proc* an injective application, such that $P = \text{proc}(E) \in \text{CSP} + T$, we define,

$$\text{Interface}(P) \subseteq \text{comm_act}(P) - \{\tau\},$$

as a set of actions that model the data or control flows on which the analysis entity interacts with its environment. P is a syntactically correct process term of *CSP+T* that models the entity E .

Modelling process interface (rule 1). *interface(P)* is made up of an input communication symbol for every entity O , which is the origin of a communication towards P , and, vice-versa, of an output communication for every destination entity D , where O and D are analysis entities with the only limitation being that both of them cannot be of type DS.

Renaming is obviously necessary when several entities D_1, D_2, \dots, D_n on a DFD accept the same input flow and, vice-versa, when several entities, O_1, O_2, \dots, O_n accept the same output flow, as otherwise the *CSP* communications could deadlock. The control transformation process (CTP) interface is modelled in the same way by including events with a

special meaning in *comm_act*. These are called e, d, t , after the SA/RT synchronization events *enable, disable, trigger*, that a CTP uses to control its DTPs.

Modelling continuous data flows (rule 2). Continuous data flows cannot be directly modelled by means of communication events in *CSP+T*, since in the latter the communication is understood to be a synchronous message passing between 2 processes and a continuous flow of data denotes an uninterrupted communication between different processes. It is therefore necessary to write an extra process (termed S in the rule) for each continuous data flow.

Modelling State Transition Diagrams (rule 3). Every CTP, called P , of the lower level in the SRM hierarchy is represented by a unique STD from the point of view of control specification. An STD can be defined as a tuple (Q, C, A, T, q) in which:

- Q is a set of states.

- C is a set of conditions, i.e., every condition denotes the occurrence of an external event, which corresponds to an input flow of control in P , or to the occurrence of an internal event which is different from any internal control flow in P , such as the internal action τ .

- A is a set of actions. An action causes the execution of an activity in the system. It can be easily identified since it corresponds to an output control flow in a DTP, or to the occurrence of an internal event of an STD.

- T is a set of transitions. A transition is a tuple of the form (q_1, c, a, q_2) in which $q_1, q_2 \in Q$, $c \in C$ or is null, $a \in A$ or is null, and its interpretation is: if in state q_1 , condition c is satisfied, then action a will be performed and also a change to state q_2 will occur. Either c or a can be *nul*.

- q is the initial state of the STD and $q \in Q$.

The transition concept can be extended to specify timing constraints in the system by describing enabling intervals and marker events.

Timing constraints. These constraints can be described as a set R of tuples (e_1, I, e_2) in which $e_1 \in (C \cup A)$, and e_1 receives the name of the *marker event*, I is a real number interval of the form $[\alpha, \beta]$, where $\alpha, \beta \in \mathbb{R}^+$, and $\alpha \leq \beta$ or I is an interval relative to the preceding event or to the event e_1 . $I(e_1)$ denotes the interval I in the following text and $e_2 \in C$ or $e_2 \in A$ receives the name of a *restricted event*. The interpretation of a timing constraint R is as follows: *event e_2 can only occur within the interval of time I from the occurrence of event e_1 , where both events can represent the satisfaction of a condition c or the execution of an action a .*

If the restricted event coincides with condition c , this means that the condition is satisfied during the time interval I to which it is restricted, the satisfaction of the condition outside the interval not being considered. In the case of the restricted event being

action a , the system is forced to carry out this action within the interval I to which it is restricted, or otherwise the process in which the restricted event is programmed fails.

Modelling timed Control Transformation Processes. A *timed* STD is therefore considered as the tuple (Q, C, A, T, q, R) , i.e. an initial STD plus the timing constraints imposed on the system. The process that models the STD is the process associated to the initial state of the system, i.e. it is *activated* when the system starts. According to the SA/RT rationale, transitions exiting the same state are associated with different events. The deterministic choice operator $|$ is used to represent different outgoing transitions from a given process state.

Modelling data and control storages (rule 4).

A Data Store (DS) in the SRM is simply a class of entities capable of storing *pieces of* information for which we cannot assume any structure or formal definition. Therefore, no mechanism to specify data or to retrieve/insert data from/into a DS has been anticipated in the SRM of a system.

In our system specification model, a DS or a CS is modelled as a CSP+T process capable of accepting information by communicating with other processes, or capable of offering its stored data through another communication.

Modelling of Transformation Specifications (rules 5 and 6). There is no agreement on how to perform the correct specification of a PSPEC (primitive DTP) in SA/RT. The specification of PSPECs is usually carried out in pseudocode, *structured English*, pre/post-conditions, etc. In this respect, we assume that the functionality of primitive DTPs is simple enough to allow us to obtain a model for each DTP as a single CSP+T process. There is, therefore, room for the analyst to set the concrete semantics that any primitive DTP should have, according to the system being modelled.

A primitive CTP in a DFD is specified by means of an STD, in such a way that any flow of events that occur in a CTP results in a condition or action in its associated STD.

Hierarchic integration of the entities of a diagram (rule 7). Since we follow a bottom-up design method, we begin by applying the above rules to the lower level schemes of the SRM of the system. When all the entities in which a diagram *explodes* have been modelled from its component processes, we are able to obtain the complete diagram represented by a complex CSP+T process term. The definition of the *interface* is not recursive, since the term *interface* (CTP) on the right side of the equation is previously calculated by rule 1. The functioning of the method is based on an iterative composition of the constituent processes and on the abstraction of their internal communications. The iterative process

finishes when the context diagram of the system is obtained.

Systematic derivation process

The above set of transformation rules are the basis of our complete top-down systematic specification technique, and are listed in Table I. In general, the following steps are taken:

- 1) Prepare the analysis schemes for carrying out the transformation. It may be necessary to rename some analysis entities to avoid conflicts (i.e., unwanted synchronizations) when constructing their model in CSP+T.
- 2) Transform the control transformation (CTP) and data transformation (DTP) schemes of the lower level, i.e. those that do not explode into other schemes, into CSP+T processes.
- 3) Select the other schemes in ascending order, i.e. a CSP+T process for each Data Storage (DS), Control Storage (CS), Continuous Flow of Data, DTP or CTP that appears in the scheme, and build a CSP+T process for each entity within the scheme.
- 4) Once the CSP+T model has been obtained for all the entities in a scheme, one CSP+T process is defined to model the complete scheme. If this scheme is already included in a CTP or a DTP of a higher level, repeat from step (3), thus progressively integrating the CSP+T model of the system in an ascending way.

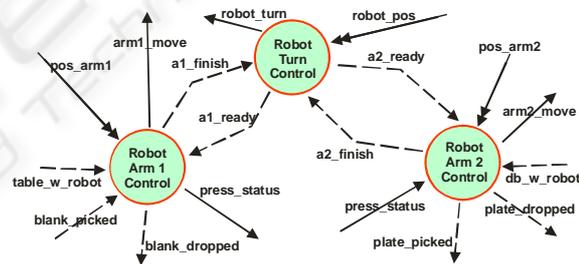


Figure 4: 2nd level DFD for a generic Robot Control P.

- 5) The process of hierarchic integration of transformation schemes finishes when the model of the *System Context Diagram* is obtained.

4 SPECIFICATION EXAMPLE: THE PRODUCTION CELL

Let us first present a detailed modelling of the Robot Control Process (*RC*) of the Production Cell Control System, since it is the process with the richest functionality among those conforming its design. We assume that *table*, *press* and *belt control* processes are already modelled, since they do not contribute additional design strategies to the general compre-

hension of the method. Finally, the integration of all the derived schemes is obtained to show that the interface of the unique process coincides with the data+control flows shown in the SCD of Fig. 2.

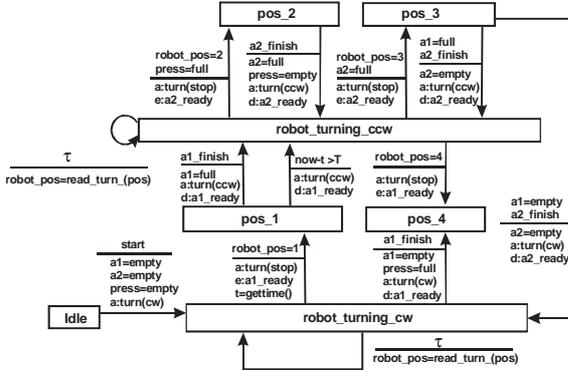


Figure 5: STD for Robot Turn Control.

Robot Turn Control (RTC)

The flows accepted by this process, Fig. 4, ensure that the robot arms are able to reach the positions established as safety positions in the production cell specification (Lewerentz, 1995), i.e., there are no collisions between *robot arms* and *belts*, blanks cannot fall from the *table* or the *press*, etc. The robot's safe positions are named in the sequel, *robot_pos_0* (initial position with arms retracted), *robot_pos_1* (arm 1 is placed in front of the rotating table and is prepared to extend and pick up a blank), *robot_pos_2* (arm 2 points towards the press and is prepared to initiate picking up a plate), *robot_pos_3* (arm 2 points to the deposit belt to place a plate on it), finally *robot_pos_4* (arm 1 is prepared to extend and deposit the blank on the press). The latter robot arms positions are specified by the STD shown in figure 5.

Firstly, we need to identify the *marker* event, i.e. *robot_pos_1*, of the syntactic CSP+T term named *RTurnCW*. The latter represents the subprocess of RTC controlling the clockwise (CW) movement of robot arm 1; then the event *a1_finish* in the term *POS_1* reports that arm 1 is positioned and has reached the point *POS_1*. On reaching this point, there are two alternative cases to consider: either a blank is on the table and the arm 1 gripper picks it up within the deadline *T*, or no new blanks have prompted before the deadline expired and therefore there needs to be a timeout. Since we must model a process with 2 alternatives, rule 3 must be applied in order to construct it,

```

Robot_Turn_Control → RTC
RTC = start → actions(start) → RTurnCW
RTurnCW = (robot_pos_1 ▷ <t → action(RturnCW) → POS_1
| robot_pos' ? robot_pos → RTurnCW)
POS_1 = (I1(robot_pos_1, a1_finish) →

```

```

actions(POS_1_CCW) → RTurnCCW
| I2(robot_pos_1, a1_finish) →
action(blank_timeout) → RTurnCCW

```

Its associated enabling intervals are defined as follows,

```

I1(robot_pos_1, a1_finish) = [t, t+T]
// the robot arm 1 has picked up a
blank from the table.
I2(robot_pos_1, a1_finish) = (t+T, ∞)
// there was no blank to pick up
within time T.

```

action(event) is a notation used to summarize all the actions associated with a transition, *event* represents the condition or the process term in which it is defined, for instance, *action(RturnCW) = activate_stop_turn; t := gettime()*.

After arm 1 picks up a blank from the table, we use the syntactical term *RTurnCCW* to indicate that the rotation of robot arms then becomes counter clockwise (CCW),

```

RTurnCCW =
(robot_pos_2 & press_full →
action(POS_2) → POS_2
| robot_pos_3 & a2_full →
action(POS_3) → POS_3
| robot_pos_4 → action(POS_4) → POS_4
| robot_pos' ? robot_pos → RTurnCCW)
POS_2 = a2_finish → action(POS_2_CCW) →
RTurnCCW
POS_3 = (a1_full & a2_finish →
action(POS_3_CCW) → RTurnCCW
| a1_empty & a2_finish →
action(POS_3_CW) → RTurnCW)
POS_4 = a1_finish → action(POS_4_CW) →
RTurnCW

```

Having reached the control position described by the *RTurnCCW* process term, the robot enters into *POS_2* only if there is a forged plate on the press, otherwise it goes directly into position 4. The communication *robot_pos'...* representing the flow of the same name in Fig. 4 continuously informs the control process of the current position of both arms. After picking up a forged plate from the press, arm 2 turns towards the belt and gets the state given by *POS_3*. The first alternative of *POS_3* represents the case in which there is also a blank in arm 1. In this case, it continues to turn CCW so that arm 2 can put the plate on the deposit belt and arm 1 can then drop the new blank on the press; the second alternative addresses the case in which there is no blank in arm 1, in which case the plate in arm 2 is dropped onto the deposit belt and the robot turning state changes to clockwise (CW) in order to return both robot arms to the control state given by *RTurnCW*, i.e. the state previous to entering into *POS_1*.

In position 4, only arm 1 has a blank to drop onto the press, turning afterwards to position 1 when it finishes the dropping action, thereby preventing an unnecessary turn (CCW) towards the deposit belt of

arm 2 that would have been compulsory if arm 2 had held a forged plate in this position.

Robot Arms Control (RA1, RA2)

As in the specification of the above control processes, we also need to apply rules 3.1 and 3.2 to derive these processes, which model the robot arms' extension, contraction and actions on the electromagnet to pick up blanks from the belts.

Robot Control (RC)

Following the ascending order of the hierarchy of schemes within the SR model, we must now model the higher abstraction scheme, the SA/RT Robot Control entity in Fig. 3, from its integrating processes in Fig. 4. RC must offer its users a simpler communication interface than the union of the interfaces of RTC, RA1C and RA2C. Rule 7 addresses this case in order to obtain the parallel compounded CSP+T term named RC and its interface from its subprocesses. This rule takes advantage of the compositionality of process terms and their algebraic properties in order to obtain new processes seamlessly from their parts. All the internal events (*start*, *aX_finish*, *aX_ready*, *robot_pos'*, *pos_armX'*) of the RC process term must be hidden, so that the interface of RC coincides with the flows in Fig. 3.

```
Robot_Control = RC
RC = (RTC\{start, a1_finish, a1_ready,
a2_finish, a2_ready, robot_pos'} ||
RA1\{start, a1_finish, a1_ready} ||
RA2\{start, a2_finish, a2_ready} ||
RTP\{robot_pos'} || RAP1\{pos_arm1'}
|| RAP2\{pos_arm2'})
```

As fig.4 shows, the above control processes (RAC1, RAC2, RTC) receive three continuous data flows (*robot_pos*, *pos_arm1* and *pos_arm2*) reporting the current robot position and the horizontal positions of the arms, respectively, it is therefore necessary to write 3 synchronization processes (RTP, RAP1, RAP2), in addition to the above ones, to model the continuous data flows, according to rule 2.

Table, Press, Feeding Belt, Deposit Belt Control (TC, PRC, FBC, DBC)

The same transformational approach can be applied to the Table Control Process (TC) and to its components, i.e. the Press Control (PRC), the Feeding Belt Control (FBC) and the Deposit Belt Control (DBC) process.

The Production Cell

The complete system is finally obtained by integrating all the elements, together with a PS (PressStatus) process. The PS process is used to keep updated the press position that is received as a continuous data flow by the robot arm control processes. Finally, we can integrate the entire Production Cell Control System by parallel composition of the derived process terms,

```
Production_Cell = PC
```

```
PC = (FBC\{failure, start} ||
TC\{start, table_turn, turn_stop, table_move, move_stop} || RC\{start}
|| PRC\{start} || DBC\{start} ||
PS\{start})
```

The bottom-up design process is completed by defining the following instantiated CSP+T process that models the whole system and is derived from the previous *PC* term,

```
Production_Cell_Context = PCC
PCC = 0. * -> PC\{start, db_w_robot,
plate_dropped blank_dropped,
plate_picked, table_w_robot,
blank_picked, table_w_fb,
blank_loaded}
```

in which the hidden events correspond to internal communication flows between the processes appearing on the DFD of Fig. 3. By applying rule 7, these events must appear hidden in the final process describing the whole system, so that only the flows connecting the system and its environment remain in *PCC*, according to the system context diagram, as shown in Fig. 2. As this latter condition is the finalization condition of the proposed method, we can conclude that a consistent and detailed design of the Production Cell Example has been derived.

5 CONCLUSIONS

We have presented a systematic transformation procedure to obtain a correct system specification of a real-time system (The Production Cell) from a semi-formal SA/RT system user requirements essential model. Our methodological scheme is based on a set of rules that are able to transform SA/RT entities into a formal specification made up of CSP+T process terms. This system specification is used to overcome the intrinsic imprecision that SA models present in describing real-time systems, by giving a formal modelling framework that permits the analyst to define the expected behaviour and functionality of all the primitive processes in a system design, and also to define the time requirements for processes by making use of the enabling interval and marker event. Our method complements the SA/RT methods modelling facilities by using ad hoc CSP+T constructs, so that hard timing constraints on the execution of a target system under development can be reflected in its formal specification. The method has been defined in such a way that it can be easily integrated within current *Automated Software Engineering* (ASE) environments and/or formal tools based on SA/RT.

Table I: Transformation rules for RT/SA entities

rule	SA/RT entities	CSP+T model
1.1	Discrete data flow f of x , with origin O and target D .	$(\text{interface}(\text{proc}(O)) \cup \{f!x\}) \wedge (\text{interface}(\text{proc}(D)) \cup \{f?x\})$
1.2	Or discrete event flow e .	$(\text{interface}(\text{proc}(O)) \cup \{e\}) \wedge (\text{interface}(\text{proc}(D)) \cup \{e\})$
2	Continuous flow f of x , with origin O and destination D	$P = \text{proc}(f)$ $P = f?x \rightarrow S ; S = (f?x \rightarrow S \mid f!x \rightarrow S)$
3.1	STD defined as (Q, C, A, T, q) , $\forall (q_i, c, a, q_j) \in T, c \in C \cup \{\lambda\}, a \in A \cup \{\lambda\}$	$P = \text{proc}(q_i), Q = \text{proc}(q_j)$ $P = c \rightarrow (a \rightarrow Q)$ or $P = a \rightarrow Q, c = \lambda$ or $P = c \rightarrow Q, a = \lambda$ $P = c \triangleright \triangleleft m_c \rightarrow (a \triangleright \triangleleft m_a \rightarrow Q)$ or $P = a \triangleright \triangleleft m_a \rightarrow Q, c = \lambda$ or $P = c \triangleright \triangleleft m_c \rightarrow Q, a = \lambda$.
3.2	a and/or c are marker events with marker variables m_a, m_c .	
4	Data storage DS with input flows $\{f_{i1}, \dots, f_{in}\}$ and output flows $\{f_{o1}, \dots, f_{om}\}$	$P = \text{proc}(DS)$ $\text{interface}(P) = \{f_{i1}, \dots, f_{in}, f_{o1}, \dots, f_{om}\}$
5	Data Transformation Process DTP	$P = \text{proc}(DTP)$ (<i>a DTP can explode in additional entities</i>). $\text{interface}(P) = \text{interface}(DTP)$
6	Control Transformation Process CTP	$Q = \text{proc}(CTP) \vee Q = \text{proc}(STD)$ (<i>Q can model the STD associated to a primitive CTP or its explosion</i>) $\text{interface}(CTP) = Q \setminus \{\text{alphabet}(Q)\} - \text{interface}(CTP)$
7	E_1, E_2, \dots, E_p , SA/RT entities in the same scheme S	$\text{interface}(S) = E_1 \setminus \{\text{alphabet}(E_1)\} - \text{interface}(S) \parallel E_2 \setminus \{\text{alphabet}(E_2)\} - \text{interface}(S) \parallel \dots \parallel E_p \setminus \{\text{alphabet}(E_p)\} - \text{interface}(S)$

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