

# HYBRID UML COMPONENTS FOR THE DESIGN OF COMPLEX SELF-OPTIMIZING MECHATRONIC SYSTEMS\*

Sven Burmester<sup>†</sup> and Holger Giese  
*Software Engineering Group, University of Paderborn  
Warburger Str. 100, D-33098 Paderborn, Germany*

Oliver Oberschelp  
*Mechatronic Laboratory Paderborn, University of Paderborn  
Pohlweg 98, D-33098 Paderborn, Germany*

**Keywords:** Mechatronic, Self-optimization, Control, Hybrid Systems, Components, Reconfiguration, Unified Modelling Language (UML), Real-Time

**Abstract:** Complex technical systems, such as mechatronic systems, can exploit the computational power available today to achieve an automatic improvement of the technical system performance at run-time by means of self-optimization. To realize this vision appropriate means for the design of such systems are required. To support self-optimization it is not enough just to permit to alter some free parameters of the controllers. Furthermore, support for the modular reconfiguration of the internal structures of the controllers is required. Thereby it makes sense to find a representation for reconfigurable systems which includes classical, non-reconfigurable block diagrams. We therefore propose hybrid components and a related hybrid Statechart extension for the Unified Modeling Language (UML); it is to support the design of self-optimizing mechatronic systems by allowing specification of the necessary flexible reconfiguration of the system as well as of its hybrid subsystems in a modular manner.

## 1 INTRODUCTION

Mechatronic systems are technical systems whose behavior is actively controlled with the help of computer technology. The design of these systems is marked by a combination of technologies used in mechanical and electrical engineering as well as in computer science. The focus of the development is on the technical system whose motion behavior is controlled by software.

The increasing efficiency of microelectronics, particularly in embedded systems, allows the development of mechatronic systems that besides the required control use computational resources to improve their long term performance. These forms of self-optimization allow an automatic improvement of a technical system during operation which increases the operating efficiency of the system and reduces the operating costs.

A generally accepted definition of the term self-optimization is difficult to find. In our view, the core function of self-optimization in technical systems is generally an automatic improvement of the behavior of the technical system at run-time with respect to defined target criteria. In a self-optimizing design, development decisions are being shifted from the design phase to the system run-time.

There are two opportunities of optimization during runtime. The first is to optimize parameters (Li and Horowitz, 1997) the second is to optimize the structure. However, alteration of the free parameters of the system will not lead very far because many characteristics, in particular those of the controller, can be altered only in the internal structures and not just by a modification of parameters (Föllinger et al., 1994; Isermann et al., 1992).

While most approaches to hybrid modeling (Henninger et al., 1995; Bender et al., 2002; Alur et al., 2001) describe how the continuous behavior can be modified when the discrete control state of the system is altered, we need an approach that allows the continuous behavior as well as its topology to be altered in a modular manner to support the design of self-optimizing systems.

\*This work was developed in the course of the Special Research Initiative 614 - Self-optimizing Concepts and Structures in Mechanical Engineering - University of Paderborn, and was published on its behalf and funded by the Deutsche Forschungsgemeinschaft.

<sup>†</sup>Supported by the International Graduate School of Dynamic Intelligent Systems. University of Paderborn

Our suggestion is to integrate methods used in mechanical engineering and software engineering to support the design of mechatronic systems with self-optimization. We therefore combine component diagrams and state machines as proposed in the forthcoming UML 2.0 (UML, 2003) with block diagrams (Föllinger et al., 1994) usually employed by control engineers. The proposed component-based approach thus allows a decoupling of the domains: A control engineer can develop the continuous controllers as well as their reconfiguration and switching in form of hybrid components. A software engineer on the other hand can integrate these components in his design of the real-time coordination. As this paper focusses the modeling aspect we set simulation aside. Simulation results can be found in (Liu-Henke et al., 2000).

In Section 2 we will examine related work. Section 3 discusses problems resulting from reconfiguration by means of an application example. In Section 4, our approach to hybrid modeling with an extension of UML components and Statecharts will be described. Thereafter we describe our model's runtime platform in Section 5 and sum up in Section 6 with a final conclusion and an outlook on future work.

## 2 RELATED WORK

A couple of modeling languages have been proposed to support the design of hybrid systems (Alur et al., 1995; Lamport, 1993; Wieting, 1996). Most of these approaches provide models, like linear hybrid automata (Alur et al., 1995), that enable the use of efficient formal analysis methods, but lack of methods for structured, modular design, that is indispensable in a practical application (Müller and Rake, 1999).

To overcome this limitation, hybrid automata have been extended to hybrid Statecharts in (Kesten and Pnueli, 1992). Hybrid Statecharts reduce the visual complexity of a hybrid automaton through the use of high-level constructs like hierarchy and parallelism, but for more complex systems further concepts for modularization are required.

The hybrid extensions HyROOM (Stauner et al., 2001), HyCharts (Grosu et al., 1998; Stauner, 2001) and Hybrid Sequence Charts (Grosu et al., 1999) of ROOM/UML-RT integrate domain specific modeling techniques. The software's architecture is specified similar to ROOM/UML-RT and the behavior is specified by statecharts whose states are associated with systems of ordinary differential equations and differential constraints (HyCharts) or Matlab/Simulink block diagrams (HyROOM). HyROOM models can be mapped to HyCharts (Stauner et al., 2001). Through adding tolerances to the continuous behavior this interesting specification technique en-

ables automatic implementation, but support for the modular reconfiguration is not given.

In (Conrad et al., 1998) guiding principles for the design of hybrid systems are sketched. It describes how to apply techniques that are used in automotive engineering, like the combination of statecharts, blockdiagrams and commercial tools. Following this approach hybrid systems need to be decoupled into discrete and continuous systems in the early design phases. Therefore a seamless support and a common model are not provided.

Within the Fresco project the description language Masaccio (Henzinger, 2000) which permits hierarchical, parallelized, and serially composed discrete and continuous components has been developed. A Masaccio model can be mapped to a Giotto (Henzinger et al., 2003) model, that contains sufficient information about tasks, frequencies, etc. to provide an implementation. The project provides a seamless support for modeling, verification and implementation, but our needs for advanced modeling techniques that support dynamic reconfiguration are not addressed.

## 3 MODELING RECONFIGURATION

We will use the switching between three controller structures as a running example to outline the resulting modeling problems. The concrete example is an active vehicle suspension system with its controller which stems from the *Neue Bahntechnik Paderborn*<sup>3</sup> research project. The project has been initiated and worked upon by several departments of the University of Paderborn and the Heinz Nixdorf Institute. In the project, a modular rail system will be developed; it is to combine modern chassis technology with the advantages of the Transrapid<sup>4</sup> and the use of existing rail tracks. The interaction between information technology and sensor/actuator technology paves the way for an entirely new type of mechatronic rail system. The vehicles designed apply the linear drive technology used in the Transrapid, but travel on existing rail tracks. The use of existing rail tracks will eliminate an essential barrier to the proliferation of new railbound transport systems (Lückel et al., 1999).

Figure 1 shows a schema of the physical model of the active vehicle suspension system. The suspension system of railway vehicles is based on air springs which can be damped actively by a displacement of their bases. The active spring-based displacement is effected by hydraulic cylinders. Three vertical hydraulic cylinders, arranged on a plane, move the bases

<sup>3</sup><http://www-nbp.upb.de/en>

<sup>4</sup><http://www.transrapid.de/en>

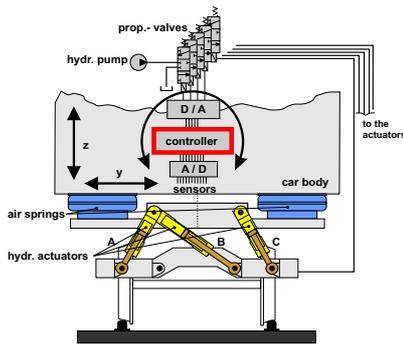


Figure 1: Scheme of the suspension/tilt module

of the air springs via an intermediate frame, the suspension frame. This arrangement allows a damping of forces in lateral and vertical directions. In addition, it is also possible to regulate the level of the coach and add active tilting of the coach body. Three additional hydraulic cylinders allow a simulation of vertical and lateral rail excitation (Hestermeyer et al., 2002). The vital task for the control system is to control the dynamical behavior of the coach body. In our example, we will focus only on the vertical dynamic behavior of the coach body. The overall controller structure comprises different feedback controllers, e.g., for controlling the hydraulic cylinder position and the dynamics of the car body.

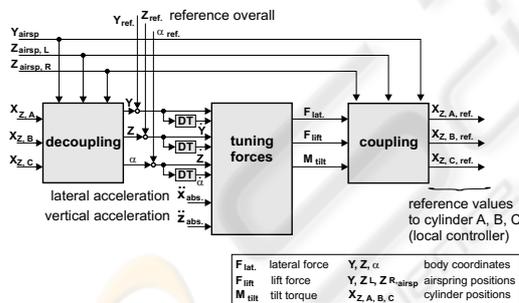


Figure 2: Reference controller

The schema of the reference controller for the overall body controller is depicted in Figure 2. The subordinated controller is not displayed here for reasons of clarity. The controller essentially consists of the blocks decoupling, tuning forces, and coupling. In the block decoupling the kinematics of the cylinders is converted into Cartesian coordinates for the description of the body motion. The actual control algorithm is in the block tuning forces. Here the forces are computed which are to affect the mechanical structure. In the block coupling the forces and torques are converted again into reference values for the subordinated cylinder controller (Liu-Henke et al., 2000).

A common representation for the modeling of mechatronic systems which have been employed here are hierarchical block diagrams. This kind of representation has its seeds in control engineering, where it is used to represent mathematic transfer functions graphically. It is widely-used in different CAE-tools. Block diagrams consist generally of function blocks, specifying function resp. behavior and hierarchy blocks grouping function and hierarchy blocks. This allows a structured draft and reduces the overall complexity of a block diagram. Between single blocks exists connections or couplings, which can have the shape of directed or non-directed links. With directed links data is exchanged whereas non-directed ones often describe functional relations or physical links, such as a link between mass and spring in multi-body system representation. While parameter optimization can be described using simply an additional connection for the parameter, the static structure of block diagrams does not permit to model structural modifications (Isermann et al., 1992).

We can, however, use special switches or fading blocks to describe the required behavior. The controller in our example has two normal modes: Reference and absolute. The controller reference uses a given trajectory that describes the motion of the coach body  $z_{ref} = f(x)$  and the absolute velocity of the coach body  $\dot{z}_{abs}$  (derived from  $\ddot{z}_{abs}$ ). The  $z_{ref}$  trajectory is given for each single track section. A track section's registry communicates this reference trajectory to the vehicle. In case a reference trajectory is not available, another controller which requires only the absolute velocity of the coach body  $\dot{z}_{abs}$  for the damping of the coach-body motion has to be used.

Besides the regular modes another controller named robust is required for an emergency; it must be able to operate even if neither the reference trajectory nor the measurement of the coach-body acceleration are available (see Figure 3).

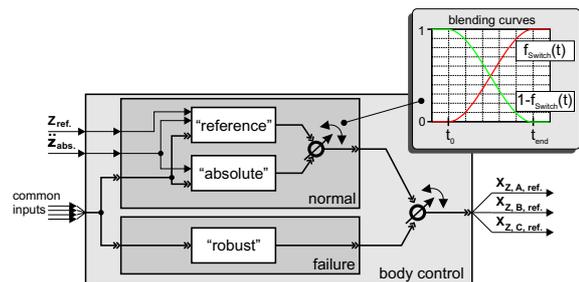


Figure 3: Fading between different control modes

For a switching between two controllers one must distinguish between two different cases: *atomic switching* and *cross fading*.<sup>5</sup> In the case of atomic

<sup>5</sup>The structure and type of cross fading depends on the

switching the change can take place between two computation steps. To start the new controller up, it is often necessary to initialize its states on the basis of the states of the old controller. In our example, the switching from the normal block to the failure block (see Figure 3) can be processed atomically because the robust controller actually has no state. Different theoretical works deal with the verification of stability in systems with any desired switching processes (Lygeros et al., 2003). In the simple case of a switch to a robust controller, stability can be maintained with the help of condition monitoring (Deppe and Oberschelp, 2000).

If, however, the operating points of the controllers are not identical it will be necessary to cross-fade between the two controllers. This handling is required in the normal block depicted in Figure 3, where a transition between the reference and the absolute controller is realized. The cross fading itself is specified by a fading function  $f_{switch}(t)$  and an additional parameter which determines the duration of the cross fading.

While the outlined modeling style allows to describe the required behavior, the result is inadequate because of the following two problems: (1) the different operation modes and the possible transitions between them are more appropriately modeled using a techniques for finite state systems rather than a set of blocks scattered all around in the block diagrams and (2) this style of modeling would require to execute all alternative controllers in parallel even though a straight forward analysis of the state space of the resulting finite state system would permit to run only the blocks which realize the currently active controllers.

To overcome these problems, hybrid modeling approaches such as hybrid automata (Henzinger et al., 1995) can be employed. When modeling the fading and switching between the different controllers in our example according to Figure 3, a hybrid automaton with at least three discrete states – one for each of the controllers– might be used. These are the locations Robust, Absolute and Reference in Figure 4. The locations' continuous behavior is represented by the associated controllers. Contrary to the white-filled arrows, black-filled ones denote the common inputs which are always available and required by all controllers.

When the automaton resides in the start location Robust and for instance the  $\ddot{z}_{abs}$  signal becomes available (as indicated by the discrete zAbsOK signal), the location and the controller change to the absolute mode. As this change requires cross fading an additional location (FadeRobAbs) is introduced in which the cross fading is comprised. To specify a fading duration  $d_1 = [d_{low}^1, d_{up}^1]$  an additional state variable

controller types and could lead to complex structures. In our example we use only output fading.

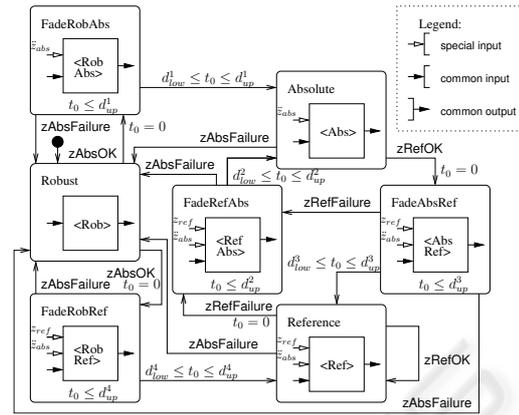


Figure 4: Hybrid body control with fading locations

$t_c$  with  $\dot{t}_c = 1$  is introduced. The reset when entering the fading location FadeRobAbs, its invariant  $t_c \leq d_{up}^1$  and the transition's guard  $d_{low}^1 \leq t_c \leq d_{up}^1$  guarantee that the fading will take  $d_{low}^1$  at a minimum and  $d_{up}^1$  at a maximum. After completing the fading, the target location Absolute is entered. The duration of the other fading transitions is specified similarly. If the  $\ddot{z}_{abs}$  signal is lost during fading or during the use of the absolute-controller, the default location with its robust control is entered immediately (cf. Figure 4).

In an appropriate self-optimizing design the aim must be to use the most comfortable controller while still maintaining the shuttle's safety. If, for instance, the absolute controller is in use and the related sensor fails, the controller may become instable. Thus there is need for a discrete control monitor which ensures that only safe configurations are allowed. The monitor which controls the correct transition between the discrete controller modes must also be coordinated with the overall real-time processing. In our example, the reference controller can only be employed when the data about the track characteristics has been received in time by the real-time shuttle control software. Trying to also integrate these additional discrete state elements into our hybrid description of the body control would obviously result in a overly complex description by means of a single hybrid Statechart which lacks modularity.

## 4 THE APPROACH

To support the design of complex mechatronic systems and to overcome the problems of the available modeling techniques as outlined in the preceding section, we introduce in the following informally our approach for modeling with the UML in Section 4.1, our notion for hybrid Statecharts in Section 4.2, our

notion for hybrid components in Section 4.3, and the modular reconfiguration in Section 4.4. The exact formalization is presented in (Giese et al., 2004).

### 4.1 Hybrid UML Model

In Figure 5c the overall structural view of our example is presented by means of a UML component diagram. The Monitor component, that embeds three components, communicates with the Registry component through ports and a communication channel. The Shuttle-Registration pattern specifies the communication protocol between Monitor and Registry. The behavior of the track section registry which is frequently contacted by the monitor to obtain the required  $z_{ref}$  is depicted in Figure 5b. In Figure 5a the sensor's behavior is described by a Statechart.

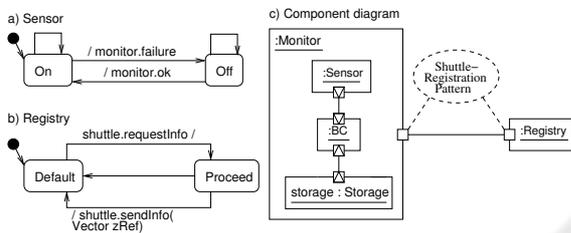


Figure 5: Monitor and its environment (incl. behavior)

The embedded components communicate continuous signals through so called *continuous ports*, depicted by framed triangles whose orientation indicates the direction of the signal flow, and discrete events through discrete, non-filled ports.

### 4.2 Hybrid Statecharts

A look at the hybrid automaton from Figure 4 reveals that the explicit fading locations considerably increase the number of visible locations of the automaton and make it unnecessarily difficult to understand.

Therefore we propose an extension of UML Statecharts towards Hybrid Statecharts that provide a short-form to describe the fading. The short-form contains the following parameters: A source- and a target-location, a guard and an event trigger, information on whether or not it is an atomic switch, and, in the latter case, a fading strategy (here cross fading is used for all *fading transitions*), a fading function ( $f_{fade}$ ) and the required fading duration interval  $d = [d_{low}, d_{up}]$  specifying the minimum and maximum duration of the fading. This short-form is displayed in Figure 6. The fading-transitions are visualized by thick arrows while atomic switches have the shape of regular arrows.

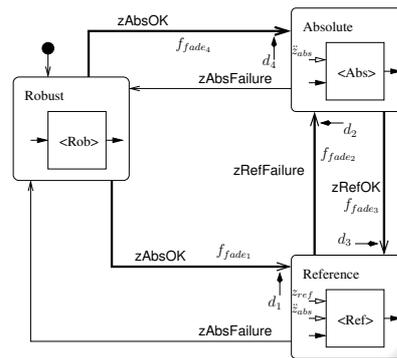


Figure 6: Behavior of the body control component

An atomic transition, leaving the source or the target state of a currently active fading transition, interrupts the execution of the fading transition. In case of conflicting atomic transitions of the source and target state, the source state's transitions have priority by default.

### 4.3 Hybrid Components

One serious limitation of today's approaches for hybrid systems is due to the fact that the continuous part of each location has to have the same set of required input and provided output variables (continuous interface). To foster the reconfiguration we propose to describe the different externally relevant continuous interfaces as well as the transition between them using hybrid interface Statecharts.

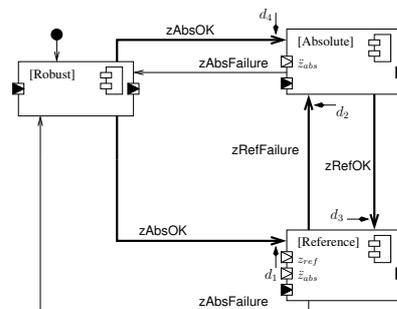


Figure 7: Interface Statechart of the BC component

The related interface Statechart of the body control component of Figure 6 is displayed in Figure 7. It shows that the body control component has three possible different interfaces. The (continuous) ports that are required in each of the three interfaces are filled black, the ones that are only used in a subset of the states are filled white. For all possible state changes, only the externally relevant information, such as durations and the signals to initiate and to break the tran-

sitions, are present.

Interface Statecharts can be employed to abstract from realization details of a component. A component in our approach can thus be described as a UML component – with ports with distinct quasi-continuous and discrete signals and events– by

- a hybrid interface Statechart which is a correct abstraction of the component behavior (cf. (Giese et al., 2004)) which determines what signals are used in what state,
- the dependency relation between the output and input signals of a component per state of the interface Statechart in order to ensure only acyclic dependencies between the components, and
- the behavior of the component usually described by a single Hybrid Statechart and its embedded subcomponents (see Section 4.4).

In our example, the BC component is described by its hybrid interface Statechart presented in Figure 7, the additionally required information on which dependencies between output and input variables exist which is not displayed in Figure 7, and its behavior described by the Hybrid Statechart of Figure 6 where the required quasi-continuous behavior is specified by controllers.

#### 4.4 Modular Reconfiguration

The hybrid statechart from Figure 6, which supports state-dependent continuous interfaces, does still not include the case that the employed controllers show hybrid behavior themselves. Instead, complete separation of discrete and continuous behavior like in (Alur et al., 2001; Bender et al., 2002; Henzinger et al., 1995; Kühl et al., 2002) is still present.

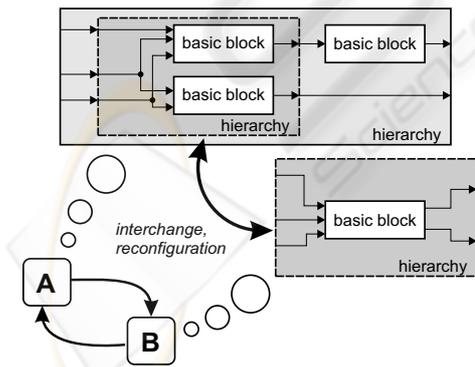


Figure 8: Reconfigurable block diagram

To overcome this limitation, we propose to assign the required configuration of embedded subcomponents (not only quasi-continuous blocks) to each state of a Hybrid Statechart by means of UML instance diagrams. This idea is motivated by the observation

that the topology of hierarchical block diagrams could be seen as a tree. With the leaves of this tree representing the behavior whereas the inner nodes describe the structure of the system. This distinction between structure (hierarchy) and function (block) can be used for the required modular reconfigurable systems. In our context, a reconfiguration can be understood as a change in the structure resp. substructure of a block diagram. It alters the topology of the system; functions are added and/or interlinked anew. Thus to realize modular reconfiguration we only have to provide a solution to alter the hierarchical elements (cf. Fig. 8). In this manner the required coordination of aggregated components can be described using a modular Hybrid statechart which alter the configurations of its hybrid subcomponents (see Figure 9).

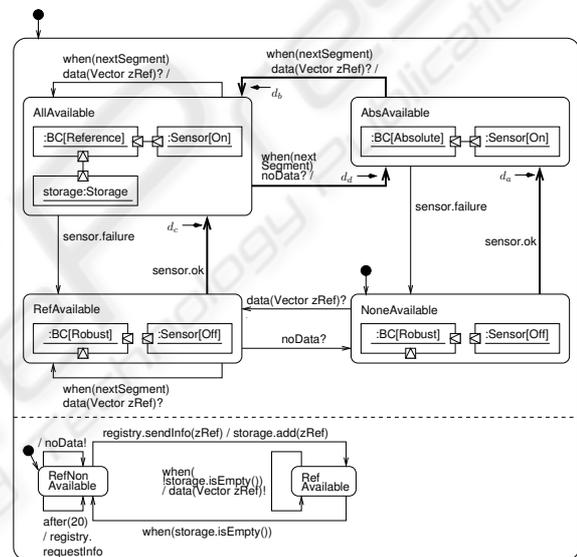


Figure 9: Monitor behavior with modular reconfiguration of the subcomponent BC

Figure 9 specifies the behavior of the control monitor software. The upper AND-state consists of four locations indicating which of the two signals  $\tilde{z}_{abs}$  and  $z_{ref}$  are available. Some transitions can fire immediately, others are associated with a deadline interval  $d = [d_{low}, d_{up}]$  specifying how much time a transition may take minimal and maximal. These transitions are thicker and marked with an arrow and the intervals  $(d_a, \dots, d_d)$ . The lower branch of the modular Hybrid statechart communicates with the track section registry (Figure 5c), frequently requests the  $z_{ref}$  function, and puts it in the storage.

In the Hybrid Statechart, every discrete state has been associated with a configuration of the subcomponents (BC, Sensor, Storage). In the design of these associations, only the interface description of the embedded component BC (see Figure 7) is relevant and

the inner structure can be neglected. Therefore, as shown in Figure 9, we can assign to each location of the upper AND-state of the statechart the BC component in the appropriate state. E.g., the BC component instance in state Reference has been (via a visual embedding) assigned to the location AllAvailable of the monitor where  $z_{ref}$  as well as  $\ddot{z}_{abs}$  are available. The required structure is specified by instances of the Sensor and the Storage component and the communication links.

The Hybrid Statechart of Figure 9 defines a mapping of states to required configurations of the sub-components. The required synchronization between the components is accomplished through explicit raising of the discrete signals defined in Figure 7.

## 5 RUN-TIME ARCHITECTURE

In order to reach interoperability for mechatronic systems, which are only alterable within certain limits, one can use appropriate middleware solutions like IPANEMA (Honekamp, 1998), which allows abstraction from hardware details. IPANEMA is a platform concept for distributed real-time execution and simulation to support rapid prototyping. It allows a modular-hierarchical organization of tasks or processes on distributed hardware.

In order to make interoperability possible also for hybrid components, which contain the kinds of alteration capability described above, the support of alteration capability by the middleware must be considerably extended. First it is necessary to generate the models in accordance to their modular-hierarchical structure. This is the basis for a reconfiguration.

In each discrete location of the system the equations, that implement the currently active controllers, have to be employed to compute the correct continuous behavior. Thus in every location of this kind only the relevant equations have to be evaluated. To reach this aim, the architecture provides means for every component to adjust the set of active equation blocks in such a way that the required reconfiguration of the component system is efficiently managed.

In the modular execution framework outlined, the required execution order results from the local evaluation dependencies within each component as well as from their interconnection. It must thus be determined at deployment-time or run-time.

The continuous nonlinear differential equations are solved by applying suitable numeric solvers. Computation is time-discrete. Incrementation depends on the solver and the dynamics of the continuous system. A time-accurate detection of a continuous condition is not possible if the controller is coupled with a real technical system. Thus, we restrict the urgent reaction

to continuous conditions in the hybrid statecharts to a detection within the desired time slot (cf. (Henzinger et al., 2003; Stauner, 2002)).

## 6 CONCLUSION AND FUTURE WORK

Complex mechatronic systems with self-optimization are hybrid systems which reconfigure themselves at run-time. As outlined in the paper, their modeling can hardly be done by the approaches currently available. Therefore, we propose an extension of UML components and Statecharts towards reconfigurable hybrid systems which supports the modular hierarchical modeling of reconfigurable systems with hybrid components and hybrid Statecharts.

The presented approach permits that the needed discrete coordination can be designed by a software engineer with extended Statecharts. In parallel, a control engineer can construct the advanced controller component which offers the technical feasible reconfiguration steps. These two views can then be integrated using only the component interface of the controller component.

It is planned to support the presented concepts by both the CAE tool CAMEL (Richert, 1996) and the CASE tool Fujaba (Giese et al., 2003). With both tools, the result of every design activity will be a hybrid component. Each of these hybrid components itself can integrate hybrid components. The integration only requires the information given by the component interface. Additional automatic and modular code generation for the hybrid models and run-time support for the reconfiguration will thus result in support for the modular and component-based development of self-optimizing mechatronic systems from the model level down to the final code.

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